

# Worksheet on Simple Graphs and Trees

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## TAKE-AWAYS

- A *simple graph* is a set of vertices along with a set of (undirected) edges with no self-loops.
- Most concepts from directed graphs, like walks and paths, can be used almost unchanged for simple graphs. A cycle in a simple graph must be of length at least 3.
- The *degree* of a vertex is the number of edges *incident* to it. Two vertices connected by an edge are *adjacent*.
- A graph is *connected* if there is a path from each vertex to each other vertex.
- A connected acyclic graph is called a *tree*; its degree-1 vertices are *leaves*.
- A tree can also be directed; edges point from *parent* to *child* away from the *root*, and the *height* is the maximum distance from the root to a leaf.
- Two graphs are *isomorphic* if they have the same number of vertices, all connected in the same way.
- Some common graphs are the  $n$ -vertex line graph  $L_n$ , the  $n$ -vertex cycle graph  $C_n$ , the  $(n+1)$ -vertex wheel graph  $W_n$ , and the  $n$ -vertex complete graph  $K_n$ .
- A  $k$ -*coloring* in a graph is an assignment of  $k$  colors to vertices so that adjacent vertices always have different colors.
- A graph's *chromatic number*  $\chi$  is the smallest number of colors needed to color it.
- A *bipartite* graph is one whose vertices can be split into two non-empty groups and all edges go from one group to the other.
- A *matching* in a bipartite graph is a way of pairing up all vertices in one group with adjacent vertices in the other.
- Colorings and matchings are useful for solving different kinds of allocation problems.

**Problem 1.** What is the average degree of the vertices in a graph  $G$ ? (answer in terms of  $|V(G)|$  and  $|E(G)|$ )<sup>1</sup>

<sup>1</sup> Hint: First compute the *sum* of the vertex degrees - how does each edge contribute to that sum? Hint 2: Handshaking Lemma (MCS Lemma 12.2.1)

**Problem 2.** Find the chromatic number for each graph and briefly justify your answers:

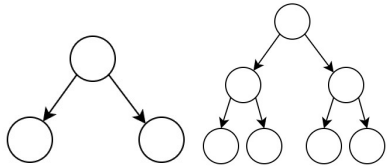
- a)  $K_n$
- b)  $L_n (n \geq 2)$
- c)  $C_n (n \geq 3)$
- d)  $W_n (n \geq 3)$

**Problem 3.** Prove that a graph  $G$  with at least two vertices is bipartite iff  $\chi(G) \leq 2$ .

**Problem 4.** Let the *signature* of a graph be the list of all vertex degrees in descending order. For example, the signature of  $L_4$  is 2, 2, 1, 1.

- a) Briefly argue that any two isomorphic graphs have the same signature.
- b) Prove that the converse is false, i.e. prove that there are some graphs that have the same signature but are not isomorphic.

**Problem 5.** For  $n \geq 1$ , define  $T_n$  as the "full binary tree" of height  $n$ , where each vertex, other than the leaves which are all distance  $n$  from the root, has exactly 2 children. For example, here are  $T_1$  and  $T_2$ :



- a) Fill in the blank, and justify your answer: For any  $k > 1$ , if  $x$  is the number of leaves in  $T_{k-1}$  then  $T_k$  has \_\_\_\_\_ leaves.
- b) How many leaves does  $T_k$  have? (Briefly justify your answer.)<sup>2</sup>
- c) Let  $k$  be some fixed natural number greater than 1, and assume that  $T_{k-1}$  has  $2^k - 1$  vertices. How many vertices does  $T_k$  have?<sup>3</sup>
- d) Let  $k$  be some fixed natural number greater than 1, and assume that  $T_{k-1}$  has  $2^k + 1$  vertices. How many vertices does  $T_k$  have?
- e) How many vertices does  $T_k$  have? (Briefly justify your answer.)<sup>4</sup>

<sup>2</sup> Hint: Start with the number of leaves in  $T_1$  and apply the lemma from (a) as many times as needed.

<sup>3</sup> Hint: Use part b. In the most straightforward solution to this part and the next, you don't have to care at all if the assumption is true or false.

<sup>4</sup> As with part (b), start with a small tree and then repeatedly apply insights from part c or d. In the next two weeks you will see a more rigorous way to prove this kind of thing using induction.