Name:												
NetID:			_	Le	ectur	e:	\mathbf{A}	В				
Discussion:	Thursday	Friday	9	10	11	12	1	2	3	4	5	6
(15 points) Cl	heck the (single)	box that bes	st cha	aracteri	izes ead	ch iten	1.					
one input, A	takes $\log_2 n$ time takes x time. How double the input	w long x	+1		2x		2^x			x^2		
T(1) = c $T(n) = 4T(n)$	$\Theta(n)$ $\Theta(n^{\text{le}})$	——	$\Theta(n)$ $\Theta(n^{1})$	$\log n$) $\log_2 3$)		$\Theta(n^2)$ $\Theta(2^n)$		$\Theta(i)$				
The running $O(n \log n)$.	time of binary	search is	t	rue [false]				
-	n to satisfy the omust have a succ				true [fals	se				
expression be	ther an input lo made true by app	_	nomia	al] 6	expone	ntial			in NI	o [

Name:												
NetID:			_	${ m L}\epsilon$	Lecture:		\mathbf{A}	В				
Discussion:	Thursday	Friday	9	10	11	12	1	2	3	4	5	6
(15 points) Cl	heck the (single)	box that bes	st cha	racteri	zes eac	ch iten	n.					
Karatsuba's in multiplication algorithm	nteger	$\log n$) $\phantom{aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa$	$\Theta(n)$ $\Theta(n^{l}$		=	$\Theta(n \log \Theta(n^{\log n}))$	- /		$\Theta(n)$	· F		
T(1) = d $T(n) = 2T(n)$	$\Theta(\log 4) + n$ $\Theta(n^2)$	· · ·	$\Theta(\sqrt{n})$)				$O(n\log n)$	g(n)			
0	time of merge is efined by $T(1)$		7 27	$\Gamma(n-1)$	1) + c $1) + c$			$\Gamma(n-1)$				
Circuit satist solved in poly	fiability can be nomial time.	true [fals	e] :	not kı	nown]		
_	n to satisfy the d				rue [fal	se				

Name:												
NetID:			_	Lecture: A B								
Discussion:	Thursday	Friday	9	10	11	12	1	2	3	4	5	6
(15 points) Ch	neck the (single)	box that be	est cha	racteri	zes ea	ch item	1.					
The running to of merge	ime $\Theta(1)$ $\Theta(r)$	$\log n$) $\phantom{aaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaaa$	$\Theta(n)$ $\Theta(n^{\mathrm{l}}$			$\Theta(n \log n)$	_ /		$\Theta(n)$ $\Theta(2^n)$			
T(1) = d $T(n) = T(n - 1)$	$1) + n \qquad \Theta(n)$ $\Theta(n^{l_0})$	===	$\Theta(n \mathbf{l})$ $\Theta(n^{\mathrm{lo}})$	_ /		$\Theta(n^2)$ $\Theta(2^n)$			n^3) 3^n)			
T(1) = d $T(n) = 2T(n/2)$	$\Theta(n)$ $\Theta(n^{lo}$		$\Theta(n \operatorname{l}$ $\Theta(n^{\operatorname{lo}}$	_ /		$\Theta(n^2)$ $\Theta(2^n)$		`	n^3) 3^n)			
The solution than in the Hanoi puzzle requires $\Theta(2^n)$		true		false		not	know	n [
	c number of a gr found in polynon	-		rue [false			not k	nown	ı	

Name:												
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Discussion:	Thursday	Friday	9	10	11	12	1	2	3	4	5	6
(6 points) First of integers sorted element), and and		der. Use the										
$Merge(L_1, L_2)$	2: sorted lists of	real number	·s)									
,	is empty and L_2 (L_2 is empty or f	,			ist							
else												
(9 points) Che	eck the (single) b	ox that best	char	acteriz	es each	ı item						
T(1) = d $T(n) = 3T(n/2)$	$\Theta(n)$ $\Theta(n^{\log n})$		$\Theta(n)$ $\Theta(n^{\log n})$	$\log n$) $\log_2 3$)		$\Theta(n^2)$ $\Theta(2^n)$		Ť	(n^3) (3^n)			
The Towers of requires expor	•	true [fals	e]	not kr	nown]		
_	hromatic number requires $\Theta(2^n)$ ti		t	rue [false			not k	knowr	1	

Name:											
NetID:			_	Le	ectur	e:	\mathbf{A}	\mathbf{B}			
Discussion:	Thursday	Friday	9	10	11	12	1	2	3	4	5
(15 points) Cl	heck the (single)	box that be	st cha	racteri	izes eac	ch item	1.				
~	ime of Karatsuba defined by $T(1)$	_		$\Gamma(n/2)$ $\Gamma(n/2)$				4T(n/2)			
T(1) = d $T(n) = 2T(n$	$\Theta(n) - 1 + c \qquad \Theta(n)$	Pg ₃ ²)	$\Theta(n)$ $\Theta(n^{\operatorname{ld}})$	$\log n$)		$\Theta(n^2)$ $\Theta(2^n)$,		
_	time of the Towersively defined by			$\Gamma(n-1)$ $2T(n/2)$,		2	$T(n-2T(n_{j}))$			
_	n to satisfy the der must have a suc				true [fa	alse [
	f Hanoi puzzle can polynomial time.	true [fals	se] r	not kı	nown			

Name:											
NetID:			. L	ecture	e:	\mathbf{A}	\mathbf{B}				
Discussion:	Thursday	Friday	9 10	11	12	1	2	3	4	5	6
(15 points) Cl	heck the (single)	box that bes	st character	rizes eac	ch item	1.					
Adding eleme start of array gets longer)	O /		$\frac{1}{2}(\log n)$				$\Theta(n \log \Theta(3^n))$	g(n)			
T(1) = d $T(n) = 3T(n)$	$\Theta(n)$ $\Theta(n^{\log n})$		$\Theta(n \log n)$ $\Theta(n^{\log_2 3})$		$\Theta(n^2)$ $\Theta(2^n)$		$\Theta(i)$				
T(1) = c $T(n) = 2T(n)$	$\Theta(n)$ $\Theta(n)$ $\Theta(n^{\text{loc}}$		$\Theta(n \log n)$ $\Theta(n^{\log_2 3})$		$\Theta(n^2)$ $\Theta(2^n)$		$\Theta(i)$				
Problems in certain exponential ti	elass NP require	tru	е	false		no	ot kno	wn			
The Marker M	faking problem ca nomial time.	ın be tru	e 🗍	false		no	ot kno	wn			

Name:												
NetID:				Lecture:			\mathbf{A}	В				
Discussion:	Thursday	Friday	9	10	11	12	1	2	3	4	5	6
(15 points) Cl	heck the (single)	box that best	chara	acteriz	es eac	h item						
T(1) = d $T(n) = T(n/3)$	$\Theta(\log B) + c$ $\Theta(n^2$		$O(\sqrt{n})$ $O(n^3)$		$\Theta(n)$ $\Theta(2^n)$	´ =		$(n \log (3^n))$	(n)			
Dividing a linhalf	nked list in $\frac{\Theta(1)}{\Theta(n)}$		$(\log n)$		$\Theta(n)$ $\Theta(2)$	_		$O(n \log n)$	g(n)			
_	time of the Towersively defined by			T(n/2) $(n-1)$				$2T(n_{j})$ $T(n-1)$				
Producing all for a sentence	-	polynom	ial		expo	onentia	al _		in	NP		
The Travelling	g Salesman	polynom	ial [expo	onentia	al		in	NP		

Name:											
NetID:			$\mathbf{L}_{\mathbf{c}}$	ecture	e:	\mathbf{A}	В				
Discussion:	Thursday	Friday 9	10	11	12	1	2	3	4	5	6
(15 points) C	heck the (single)	box that best c	haracter	izes eac	ch item	1.					
T(1) = d $T(n) = T(n/2)$	$\Theta(\log 2) + n$ $\Theta(n^2)$			$\Theta(n)$ $\Theta(2)$		_	$O(n\log n)$	(n)			
input, A takes	takes 2^n time. Constant x time. How lored one to the input	x+2 ng will		2x		2^x			x^2		
T(1) = d $T(n) = 3T(n)$	$\Theta(n)$ $\Theta(n^{\log n})$		$n \log n$) $n^{\log_2 3}$)		$\Theta(n^2)$ $\Theta(2^n)$		$\Theta(r)$,			
Problems in crequire expon	class P (as in P valential time	s. NP)	true [false		1	not k	nown	ı [
The Travelling problem can be polynomial time	be solved in	true	fals	se] 1	not kr	nown]		