

# #22: kd-Trees and BTrees Intro

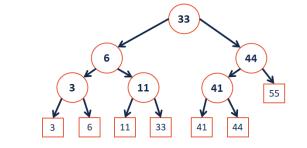
October 17, 2018 · Wade Fagen-Ulmschneider

#### **Range-based Searches:**

Q: Consider points in 1D:  $p = \{p_1, p_2, ..., p_n\}$ . ...what points fall in [11, 42]?



## **Tree Construction:**

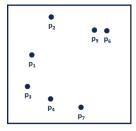


**Range-based Searches:** 

## **Running Time:**

# Extending to k-dimensions:

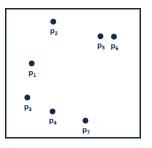
Consider points in 2D:  $\mathbf{p} = {\mathbf{p}_1, \mathbf{p}_2, ..., \mathbf{p}_n}$ :



...what points are inside a range (rectangle)? ...what is the nearest point to a query point **q**?

# kd-Tree Motivation:

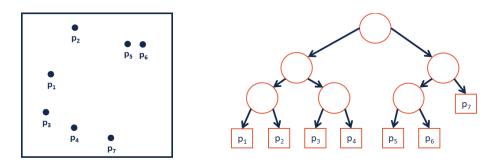
First, let's try and divide our space up:

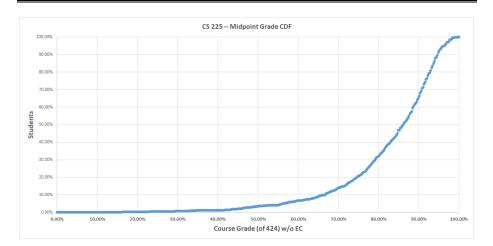


#### kd-Tree Construction:

How many dimensions exist in our input space?

How do we want to "order" our dimensions?





#### Motivation

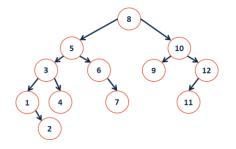
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Can we always fit our data in main memory?

#### Where else do we keep our data?

**vs.** CPU: 3 GHz == 3m ops / \_\_\_\_\_\* \_\_\_ cores

**AVL Operations on Disk:** 



How deep do AVL trees get?

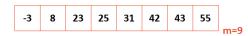
#### **BTree Motivations**

Knowing that we have long seek times for data, we want to build a data structure with two (related) properties:

1.

2.

**BTree**<sub>m</sub>



**Goal:** Build a tree that uses \_\_\_\_\_/node! \_\_\_\_\_/node! \_\_\_\_\_/node!

A **BTree of order m** is an m-way tree where:

**1.** All keys within a node are ordered.

2. All leaves contain no more than **m-1** nodes.

BTree Insert, using m=5

...when a BTree node reaches **m** keys:

## CS 225 – Things To Be Doing:

- **1.** Programming Exam B starts next Thursday (Oct 25<sup>th</sup>)
- 2. MP4 extra credit ongoing (final deadline Monday, Oct. 17th)
- 3. lab\_avl released this week; course feedback in lab this week!
- **4.** Daily POTDs are ongoing!