CS 225

Data Structures

October 11 – AVL Analysis
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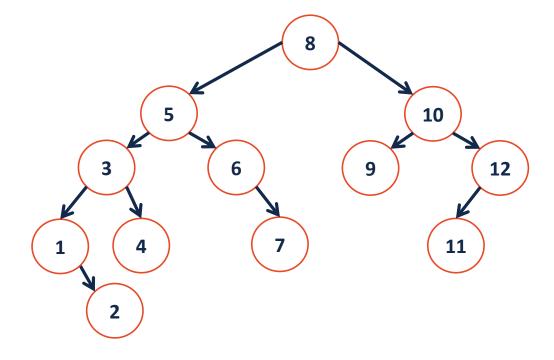
_insert(6.5)

Insertion into an AVL Tree

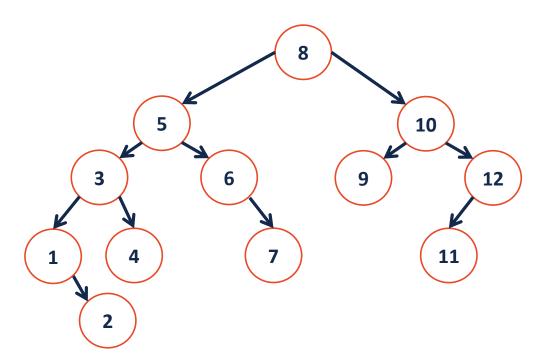
Insert (pseudo code):

- 1: Insert at proper place
- 2: Check for imbalance
- 3: Rotate, if necessary
- 4: Update height

```
1 struct TreeNode {
2   T key;
3   unsigned height;
4   TreeNode *left;
5   TreeNode *right;
6 };
```



```
template <class T> void AVLTree<T>:: insert(const T & x, treeNode<T> * & t ) {
 2
    if( t == NULL ) {
 3
    t = new TreeNode<T>( x, 0, NULL, NULL);
 4
 5
     else if (x < t->key) {
 7
      insert( x, t->left );
      int balance = height(t->right) - height(t->left);
 9
      int leftBalance = height(t->left->right) - height(t->left->left);
     if (balance == -2) {
10
       if ( leftBalance == -1 ) { rotate _____(t); }
11
                               { rotate (t); }
12
        else
13
14
15
16
     else if (x > t->key) {
      insert( x, t->right );
17
      int balance = height(t->right) - height(t->left);
18
      int rightBalance = height(t->right->right) - height(t->right->left);
19
20
      if( balance == 2 ) {
        if( rightBalance == 1 ) { rotate_____(t); }
21
                             { rotate ( t ); }
        else
23
24
25
26
     t->height = 1 + max(height(t->left), height(t->right));
27
```

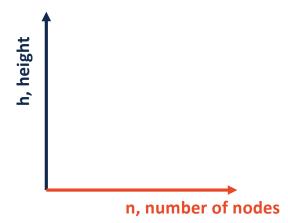


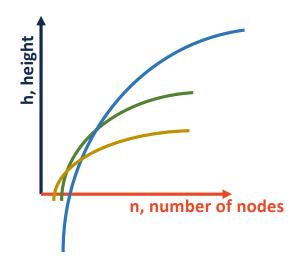
We know: insert, remove and find runs in: _____.

We will argue that: h is _____.

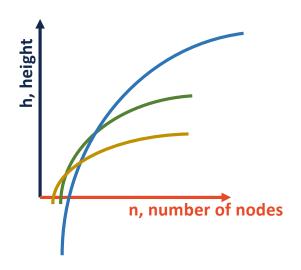
Definition of big-O:

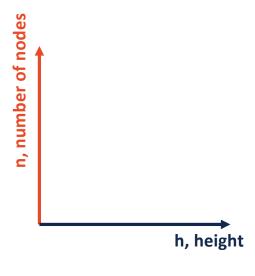
...or, with pictures:

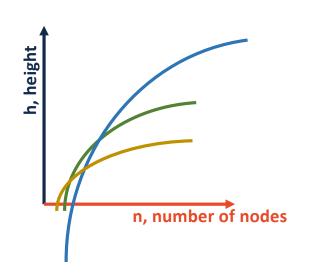


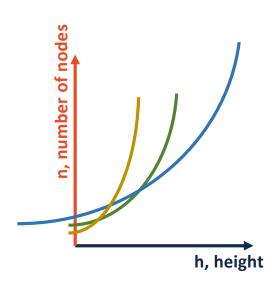


The height of the tree, f(n), will always be less than
c × g(n) for all values where n > k.









• The number of nodes in the tree, $f^{-1}(h)$, will always be greater than $c \times g^{-1}(h)$ for all values where n > k.

Plan of Action

Since our goal is to find the lower bound on **n** given **h**, we can begin by defining a function given **h** which describes the smallest number of nodes in an AVL tree of height **h**:

Simplify the Recurrence

$$N(h) = 1 + N(h - 1) + N(h - 2)$$

State a Theorem

Theorem: An AVL tree of height h has at least _____.

Proof:

- I. Consider an AVL tree and let **h** denote its height.
- II. Case: _____

An AVL tree of height ____ has at least ____ nodes.

Prove a Theorem

III. Case: _____

An AVL tree of height ____ has at least ____ nodes.

Prove a Theorem

IV. Case: _____

By an Inductive Hypothesis (IH):

We will show that:

An AVL tree of height ____ has at least ____ nodes.

Prove a Theorem

V. Using a proof by induction, we have shown that:

...and inverting:

Summary of Balanced BST

Red-Black Trees

- Max height: 2 * lg(n)
- Constant number of rotations on insert, remove, and find

AVL Trees

- Max height: 1.44 * lg(n)
- Rotations:

Summary of Balanced BST

Pros:

- Running Time:
 - Improvement Over:

- Great for specific applications:

Summary of Balanced BST

Cons:

- Running Time:

- In-memory Requirement: