



CS 225

Data Structures

April 9 – Graphs Intro

Wade Fagen-Ulmschneider

Disjoint Sets Analysis

The **iterated log** function:

The number of times you can take a log of a number.

$\log^*(n) =$

0, $n \leq 1$

$1 + \log^*(\log(n))$, $n > 1$

What is $\lg^*(2^{65536})$?

Disjoint Sets Analysis

In an Disjoint Sets implemented with smart **unions** and path compression on **find**:

Any sequence of **m union** and **find** operations result in the worse case running time of $O(\text{_____})$,
where **n** is the number of items in the Disjoint Sets.

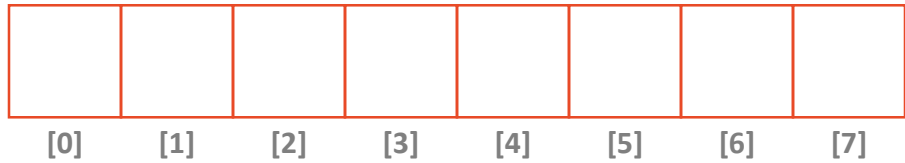
In Review: Data Structures

Array

- Sorted Array
- Unsorted Array
- Stacks
- Queues
- Hashing
- Heaps
 - Priority Queues
- UpTrees
 - Disjoint Sets

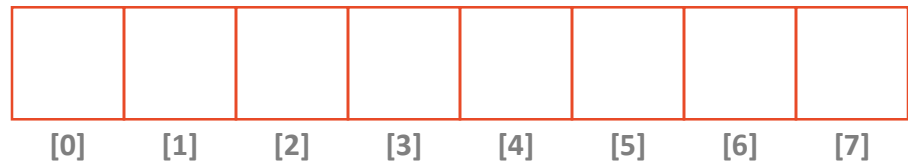
List

- Doubly Linked List
- Skip List
- Trees
 - BTree
 - Binary Tree
 - Huffman Encoding
 - kd-Tree
 - AVL Tree

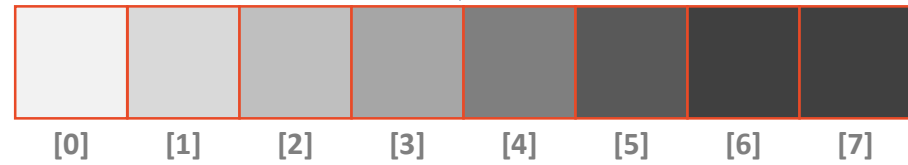


Array

- **Constant time access to any element, given an index $a[k]$ is accessed in $O(1)$ time, no matter how large the array grows**
- **Cache-optimized**
Many modern systems cache or pre-fetch nearby memory values due the “Principle of Locality”. Therefore, arrays often perform faster than lists in identical operations.

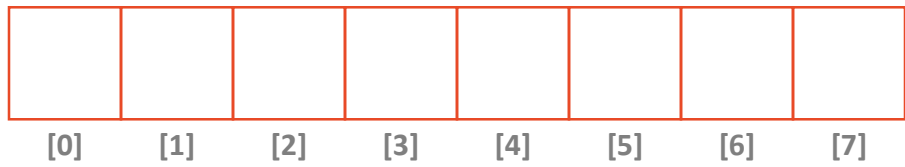


Array

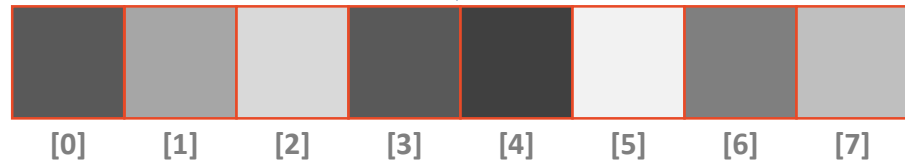


Sorted Array

- Efficient general search structure
Searches on the sort property run in $O(\lg(n))$ with Binary Search
- Inefficient insert/remove
Elements must be inserted and removed at the location dictated by the sort property, resulting shifting the array in memory – an $O(n)$ operation

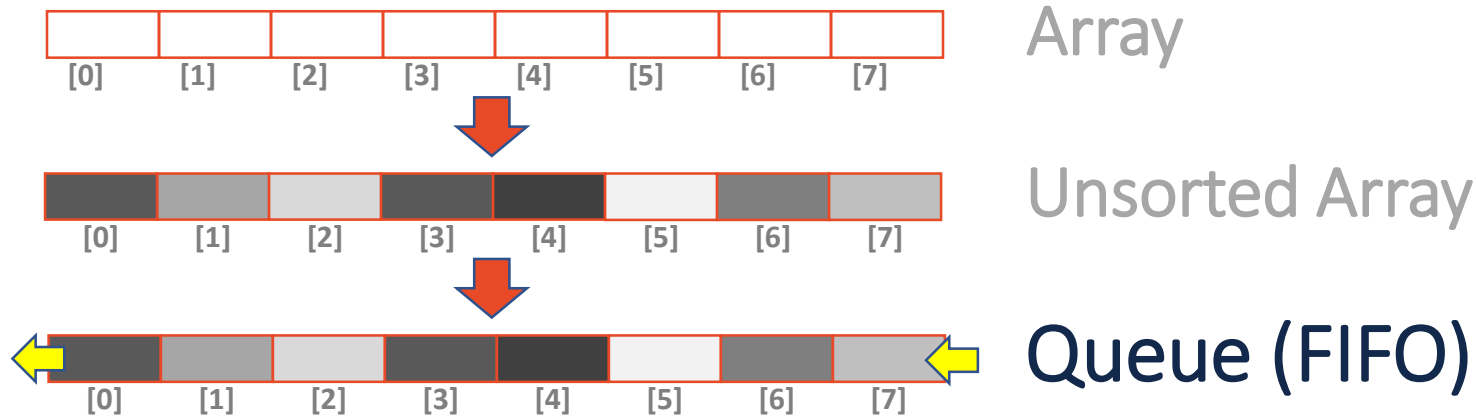


Array

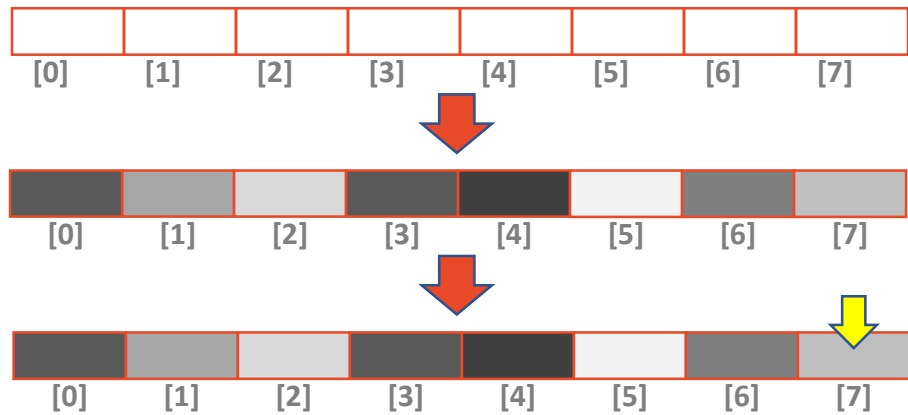


Unsorted Array

- Constant time add/remove at the beginning/end
Amortized $O(1)$ insert and remove from the front and of the array
Idea: Double on resize
- Inefficient global search structure
With no sort property, all searches must iterate the entire array; $O(1)$ time



- **First In First Out (FIFO) ordering of data**
Maintains an arrival ordering of tasks, jobs, or data
- **All ADT operations are constant time operations**
enqueue() and dequeue() both run in $O(1)$ time



Array

Unsorted Array

Stack (LIFO)

- Last In First Out (LIFO) ordering of data
Maintains a “most recently added” list of data
- All ADT operations are constant time operations
`push()` and `pop()` both run in $O(1)$ time

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In Review: Data Structures

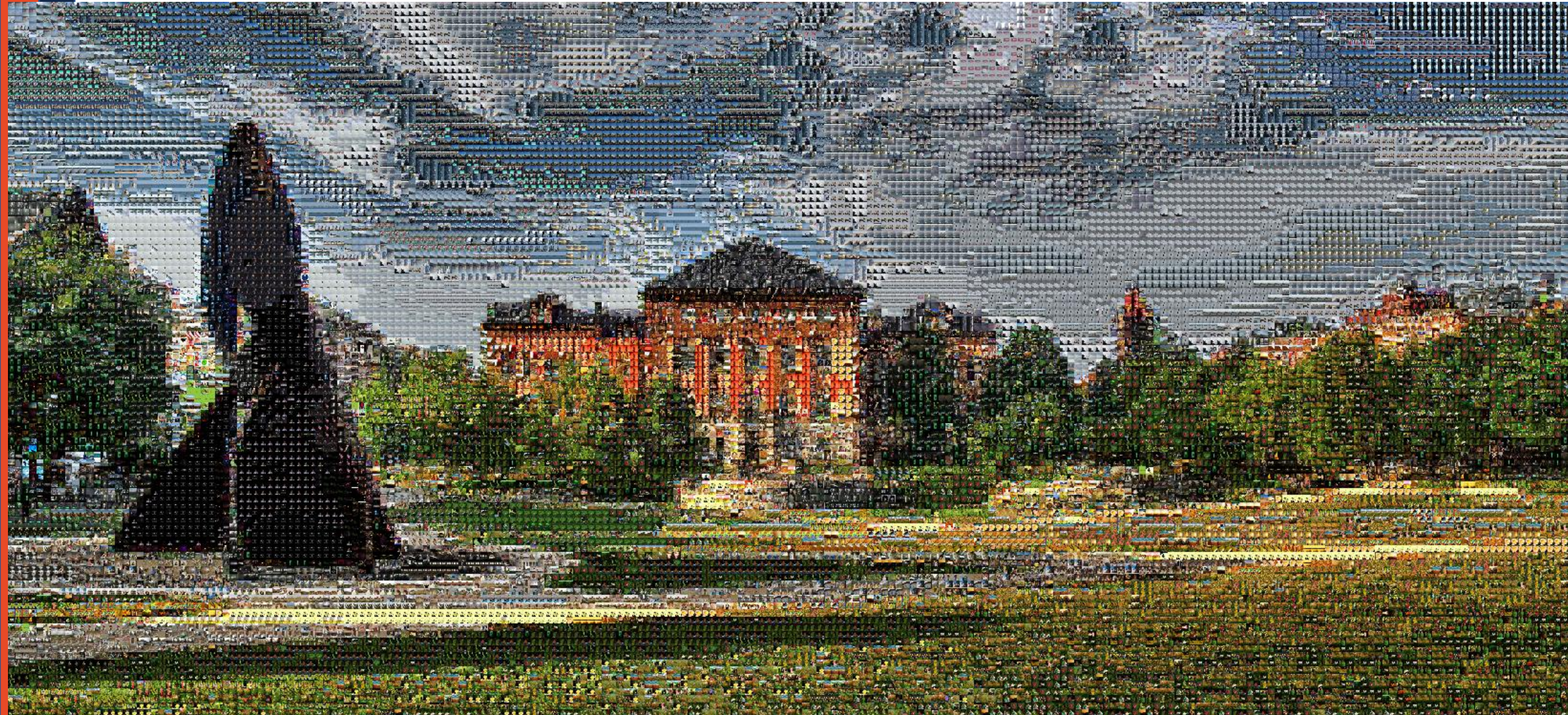
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Graphs

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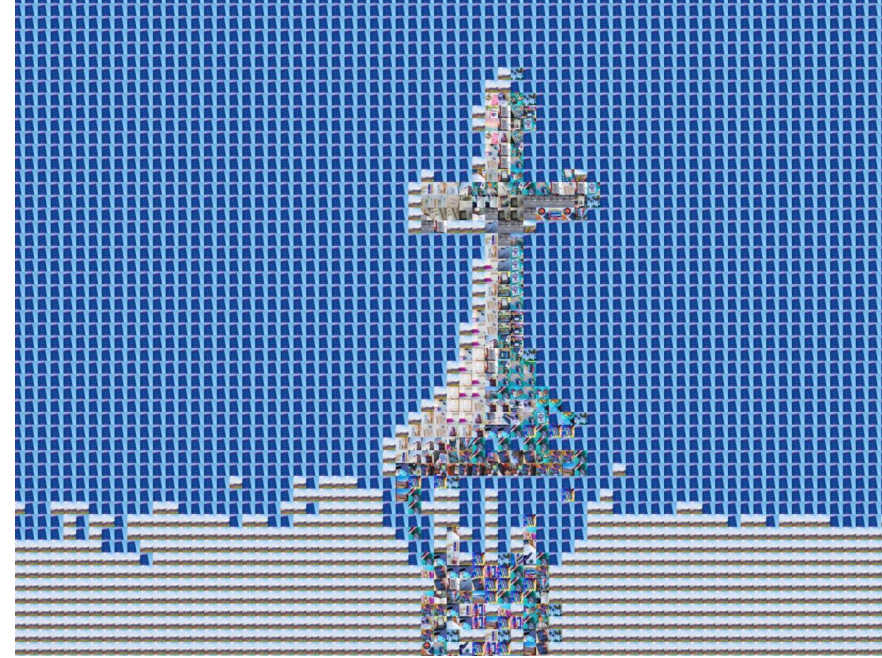
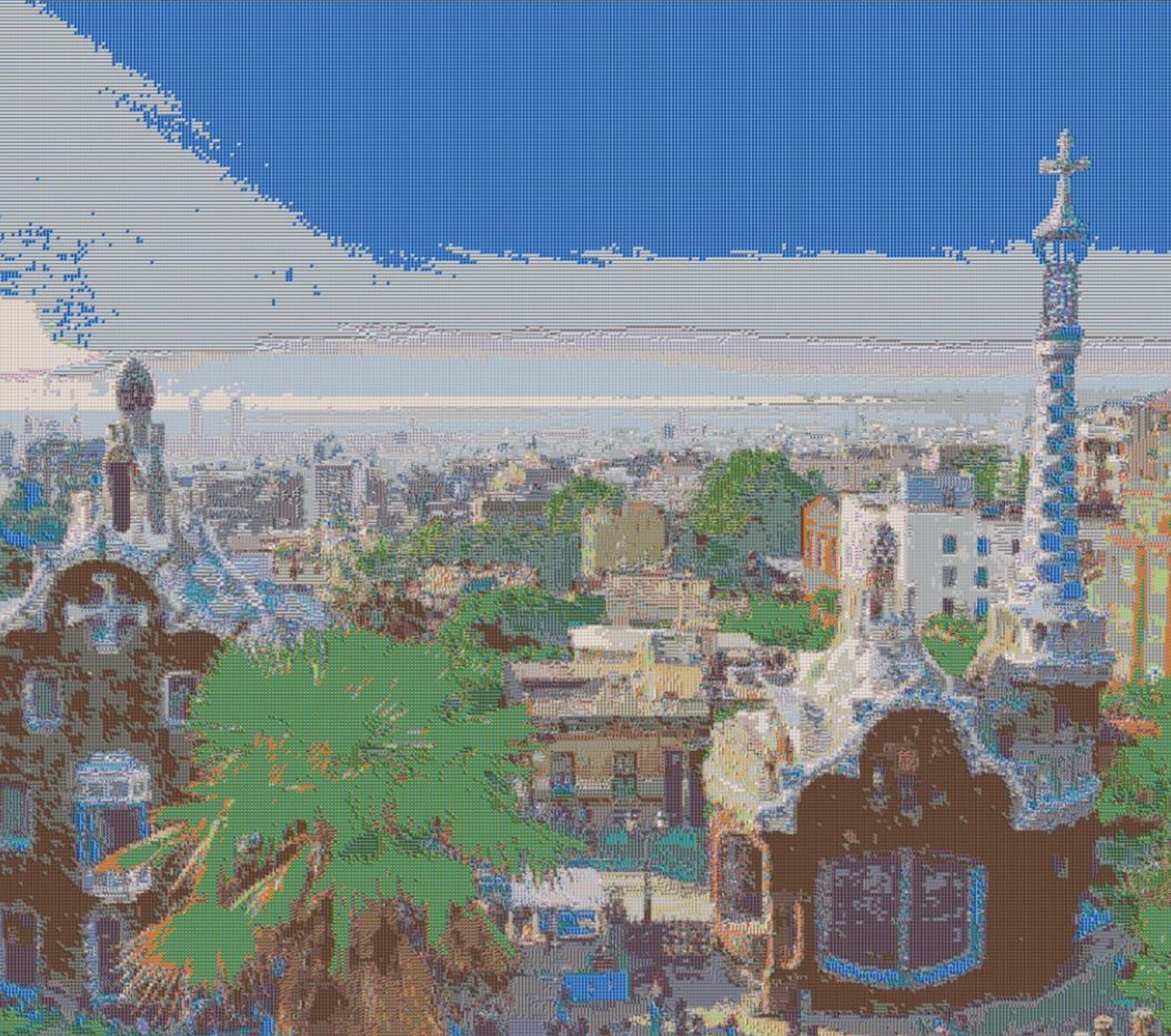


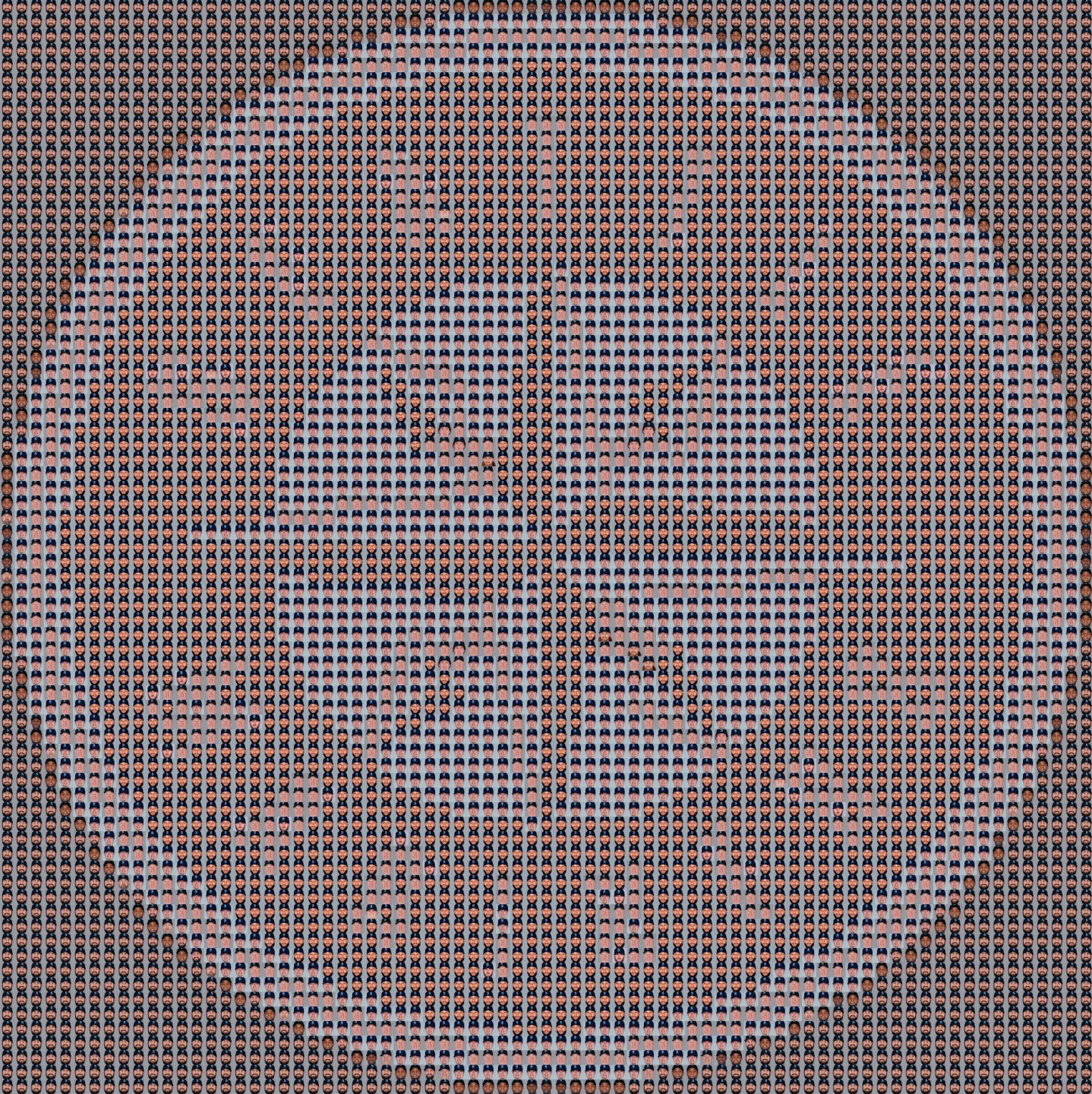


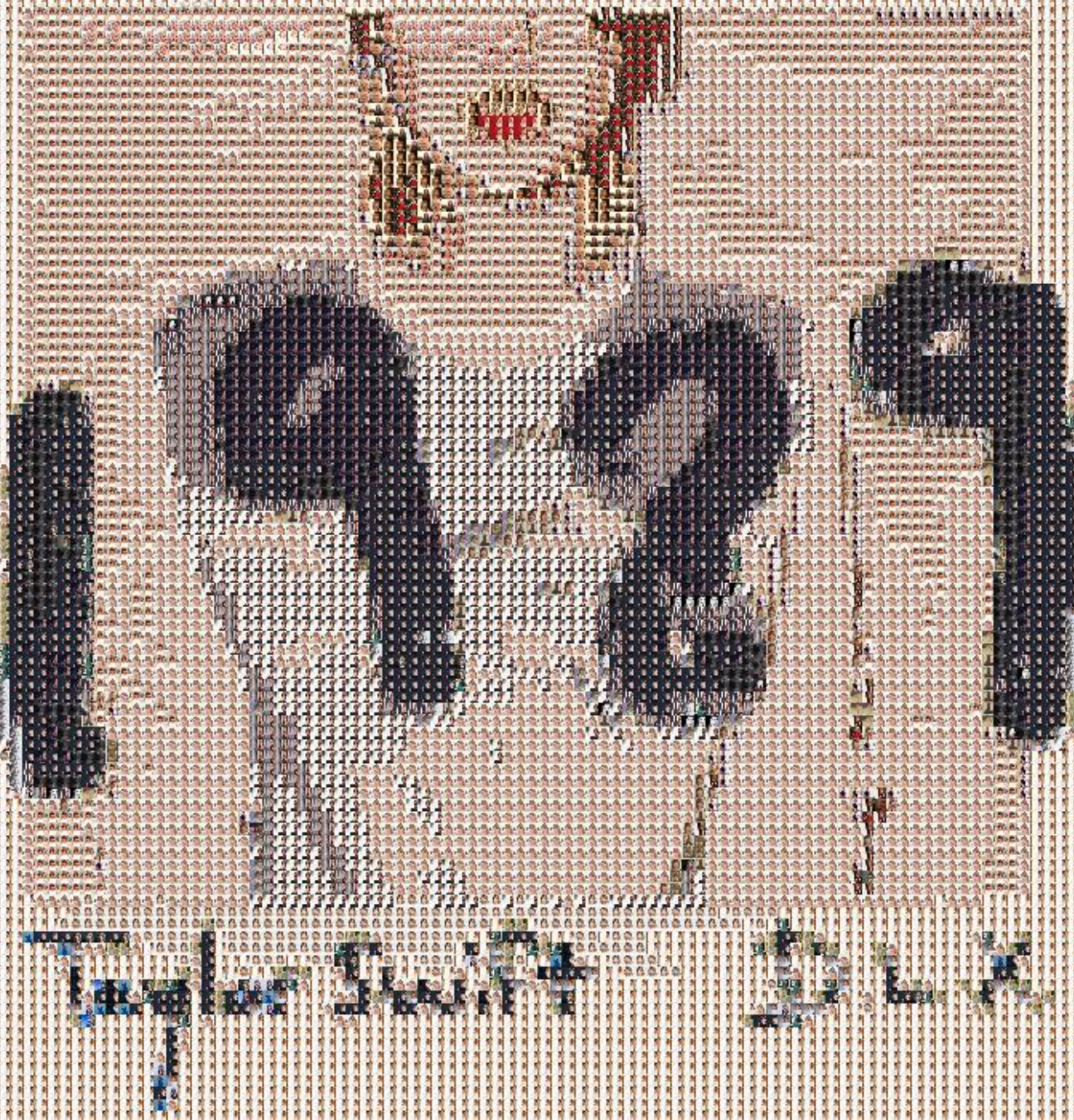














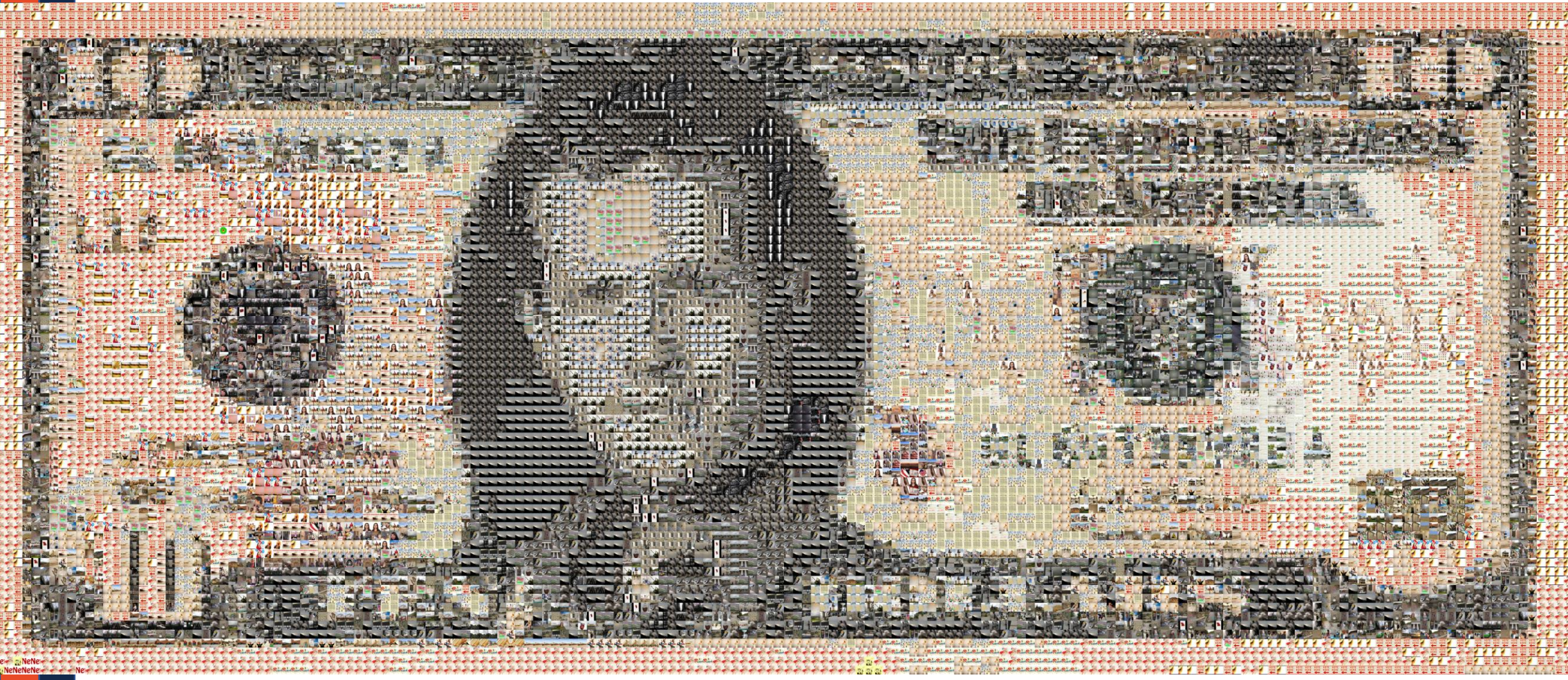
Who
Who

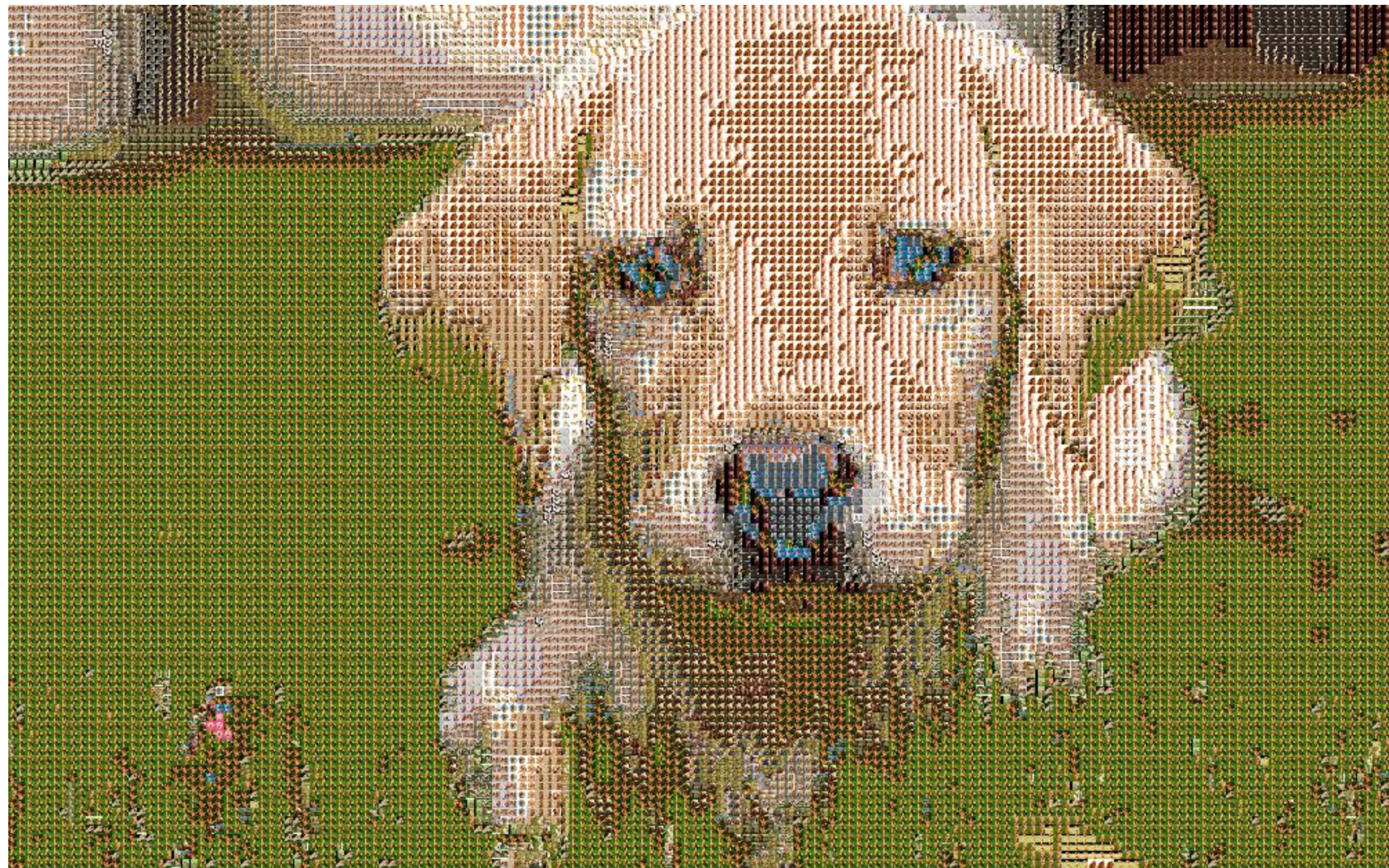
WHO IS IT?

IS

Two white rectangular boxes, likely for contestants or answers.

“When you're asked about kd-trees in an interview and Wade comes to mind:”





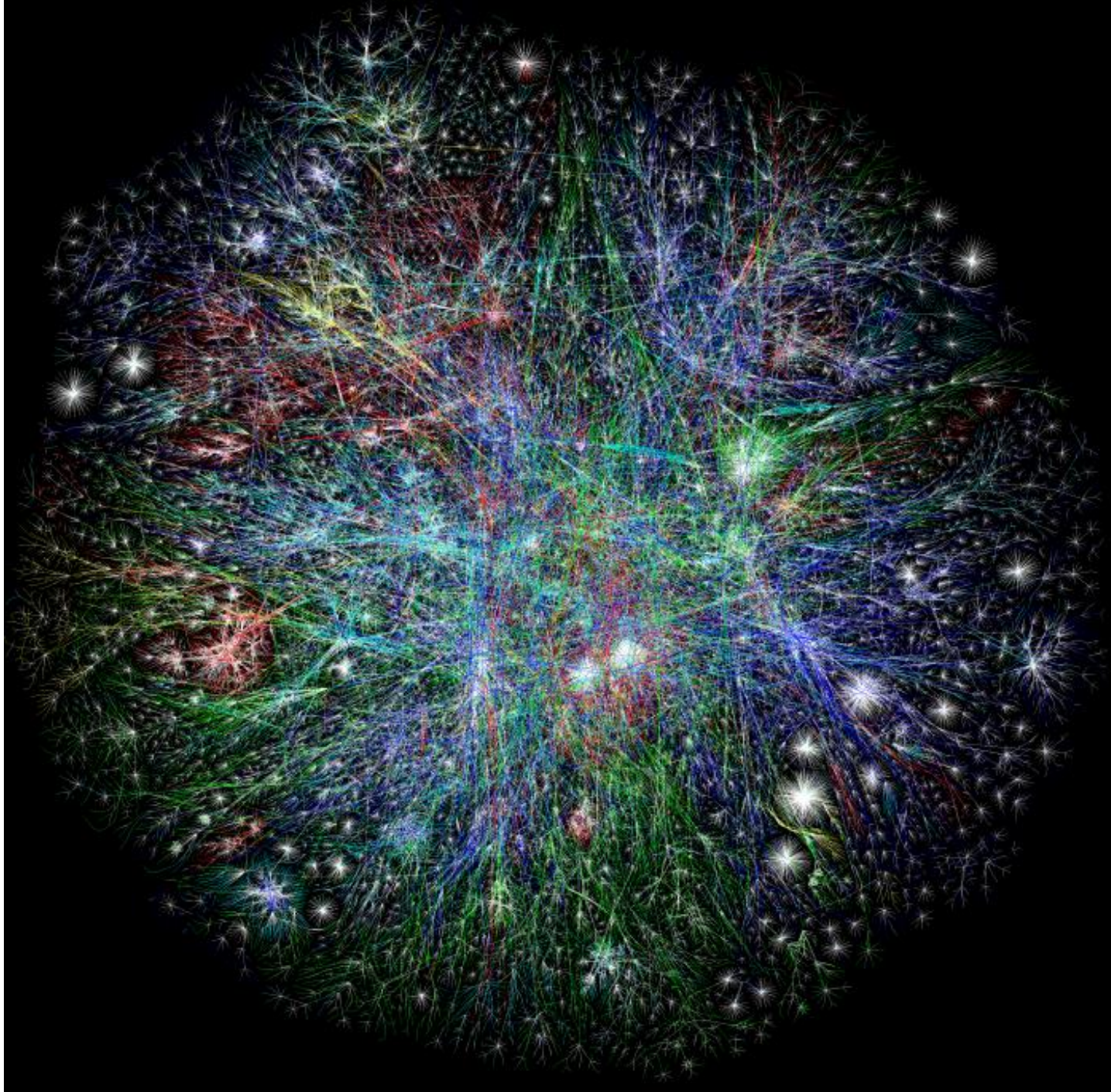


Aw, Snap!

Something went wrong while displaying this webpage.

[Learn more](#)

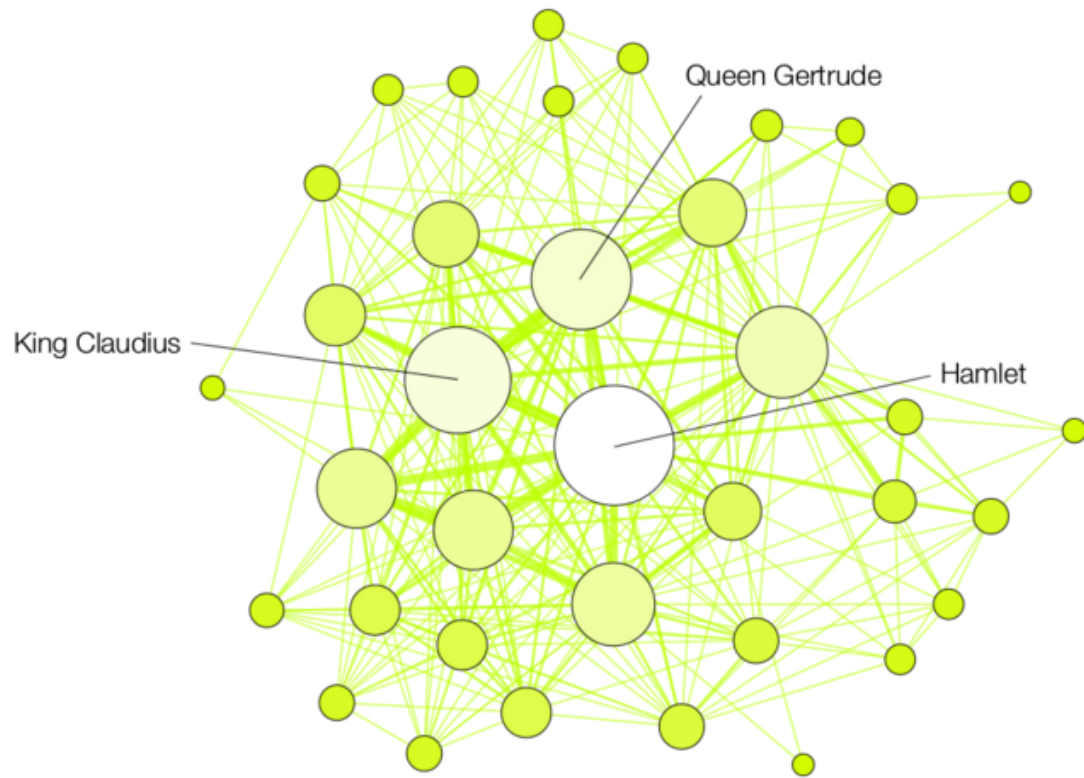
Reload



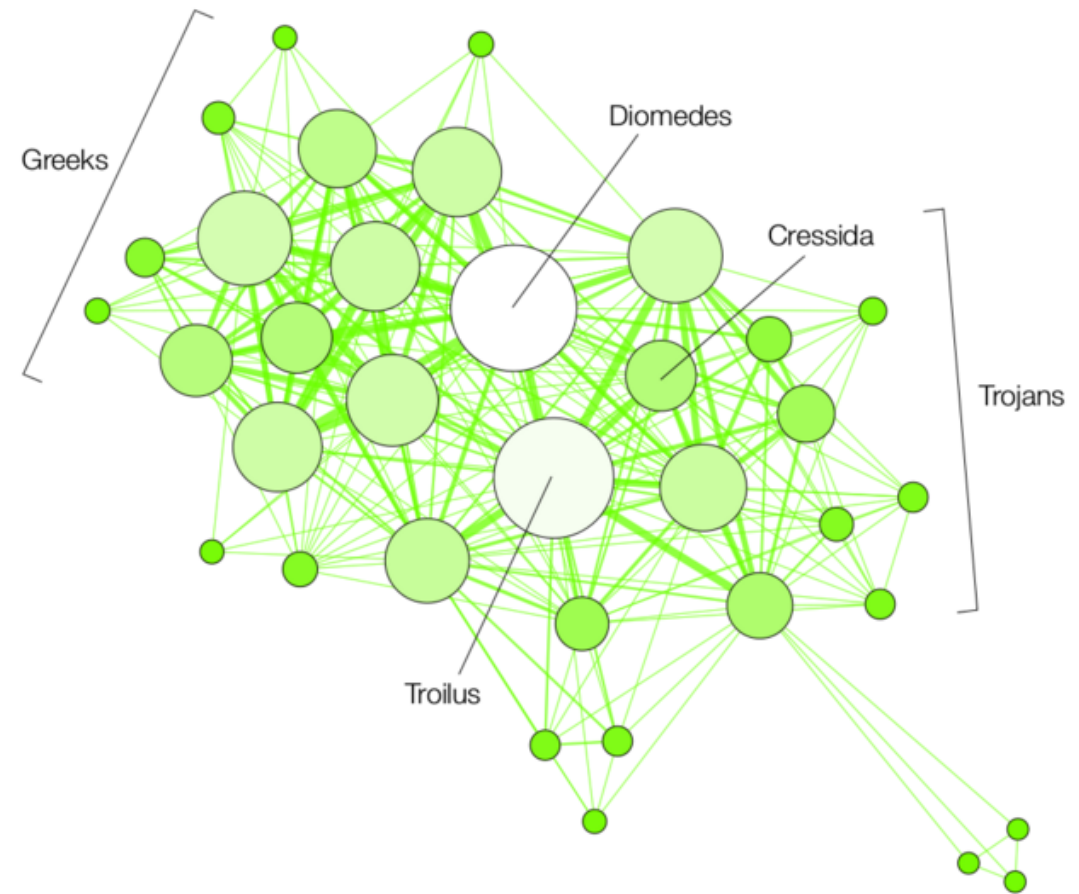
The Internet, 2003

The OPTE Project (2003)

Map of the entire internet; nodes are routers; edges are connections.



HAMLET

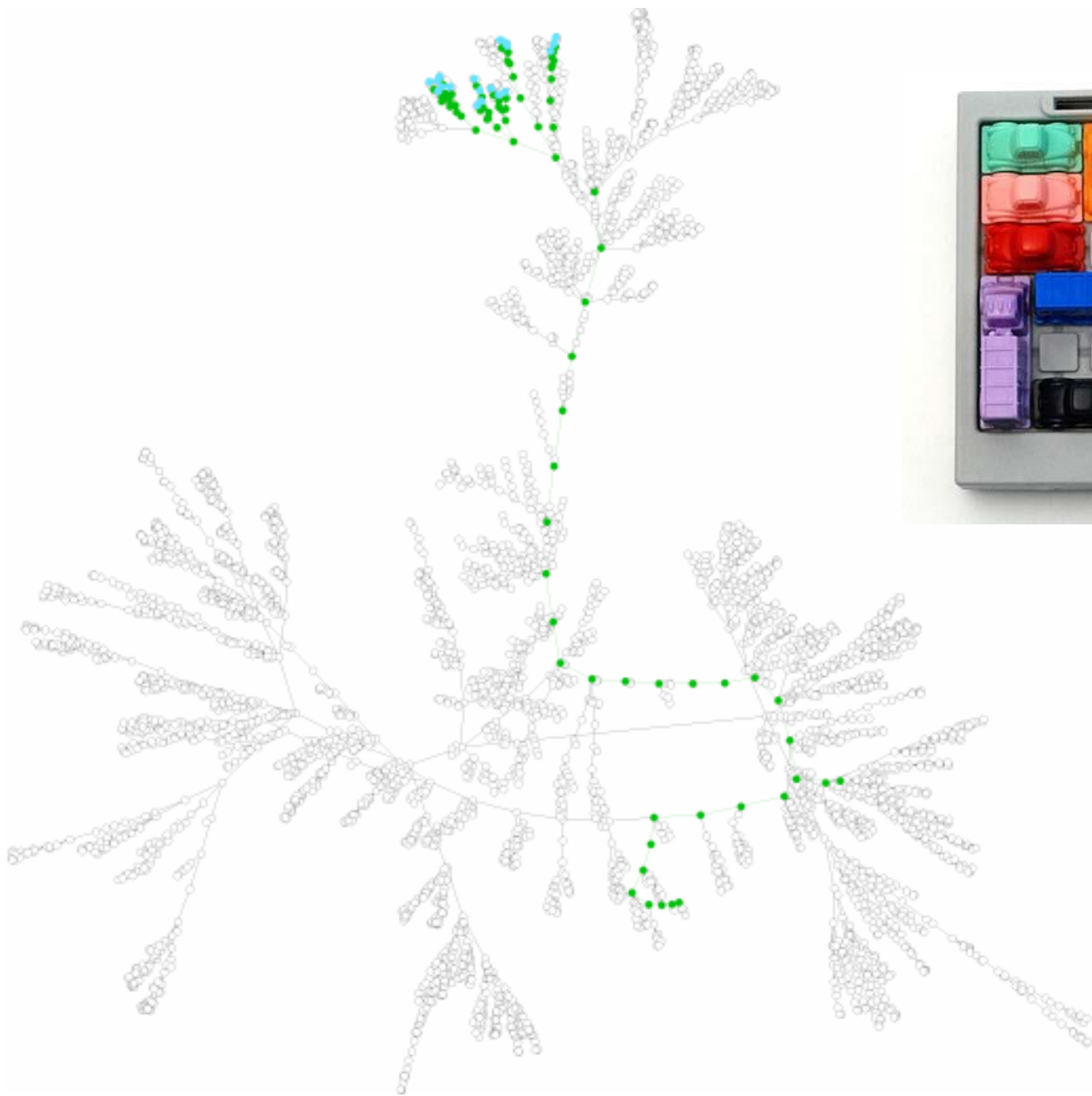


TROILUS AND CRESSIDA

Who's the real main character in Shakespearean tragedies?

Martin Grandjean (2016)

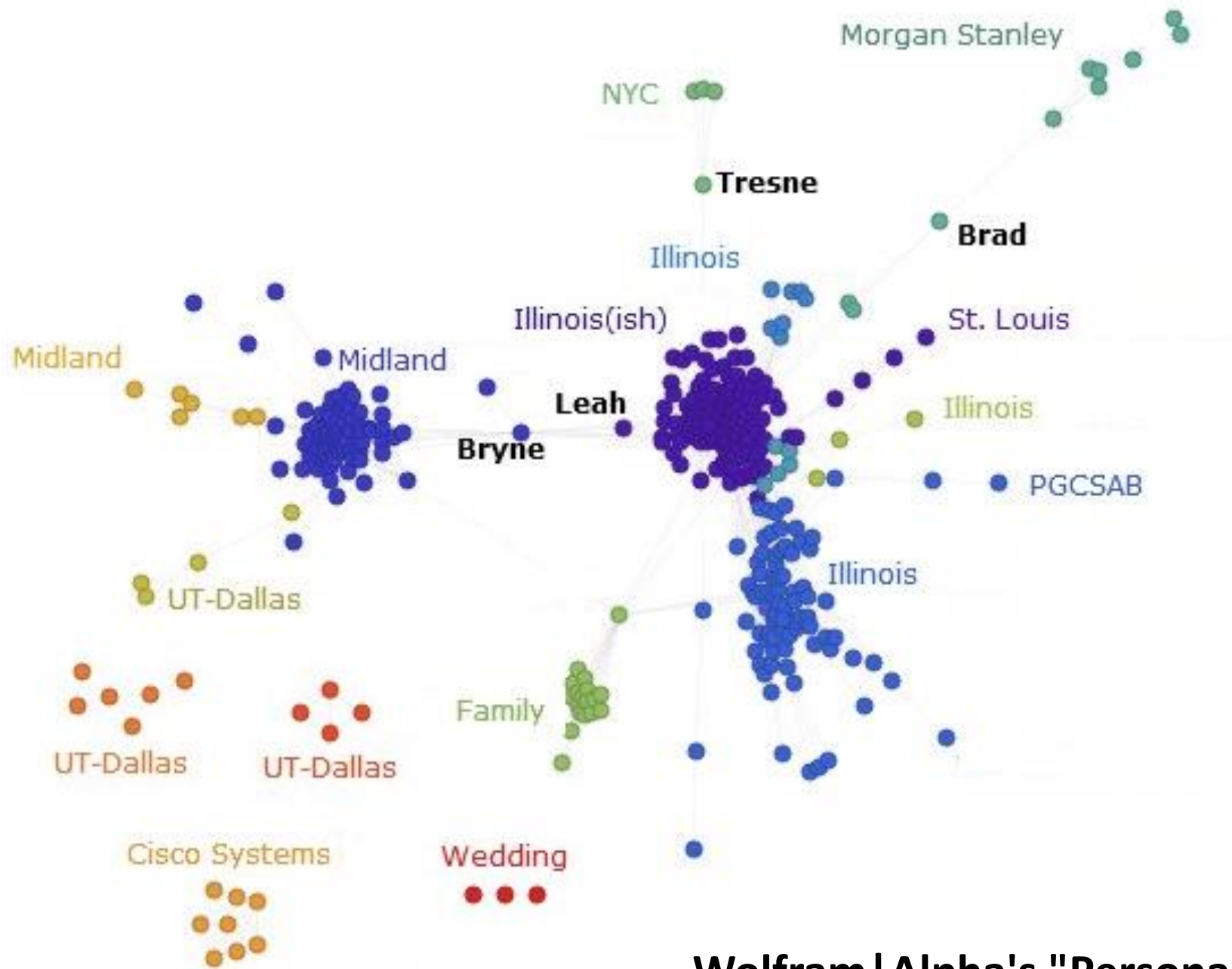
<https://www.pbs.org/newshour/arts/whos-the-real-main-character-in-shakespearean-tragedies-heres-what-the-data-say>



“Rush Hour” Solution

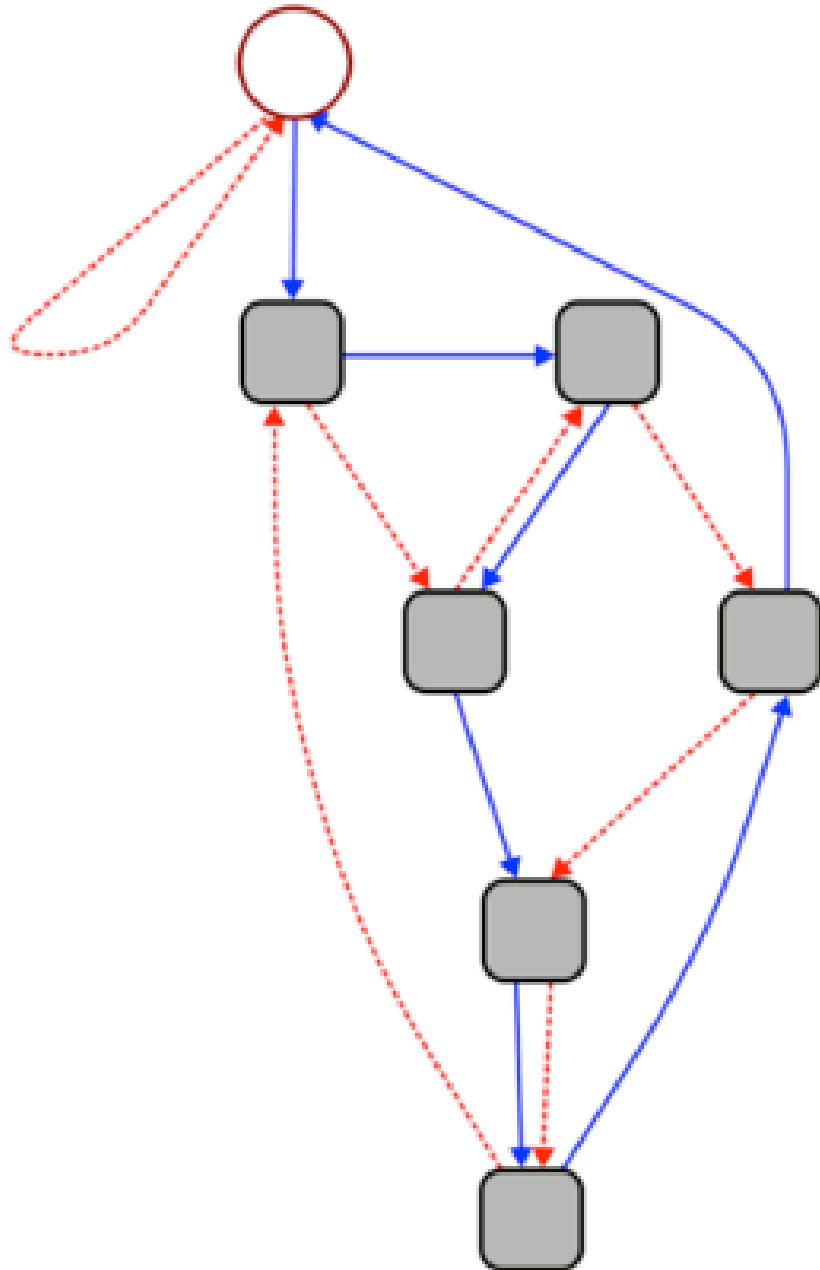
Unknown Source

Presented by Cinda Heeren, 2016



Wolfram|Alpha's "Personal Analytics" for Facebook

Generated: April 2013 using Wade Fagen-Ulmschneider's Profile Data



This graph can be used to quickly calculate whether a given number is divisible by 7.

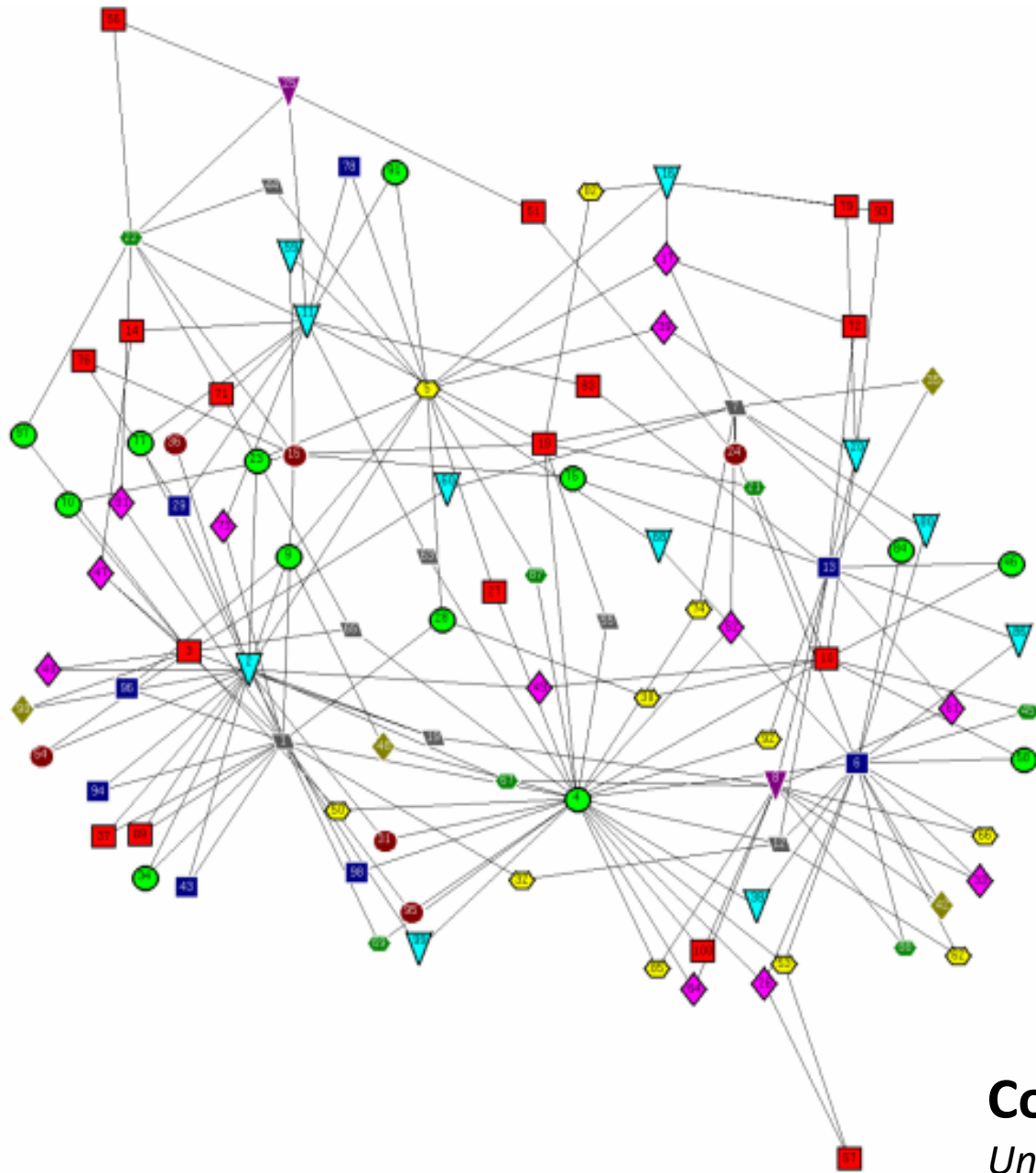
1. Start at the circle node at the top.
2. For each digit **d** in the given number, follow **d** blue (solid) edges in succession. As you move from one digit to the next, follow **1** red (dashed) edge.
3. If you end up back at the circle node, your number is divisible by 7.

3703

“Rule of 7”

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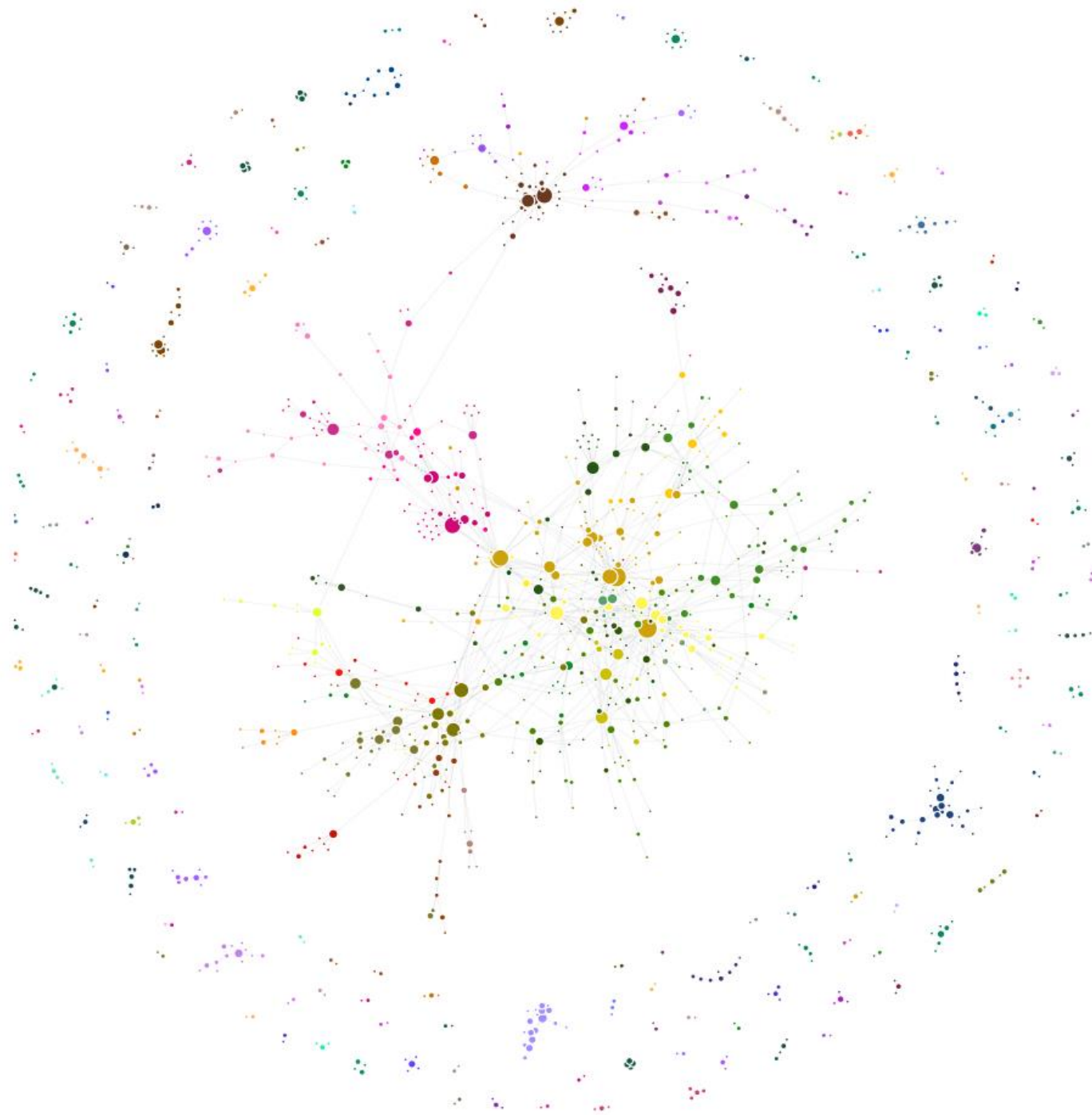
Presented by Cinda Heeren, 2016



Conflict-Free Final Exam Scheduling Graph

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Presented by Cinda Heeren, 2016

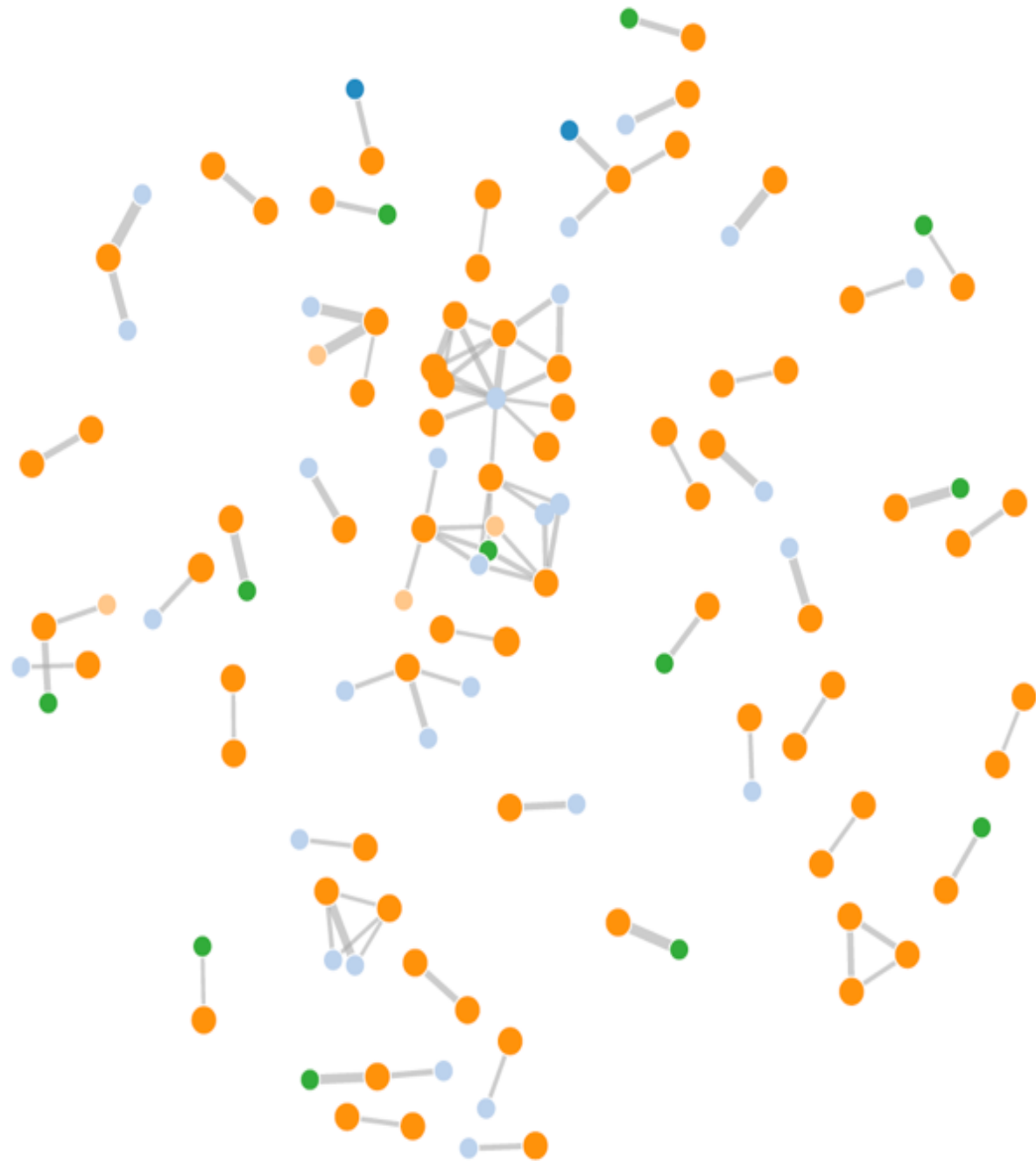


Class Hierarchy At University of Illinois Urbana-Champaign

A. Mori, W. Fagen-Ulmschneider, C. Heeren

Graph of every course at UIUC; nodes are courses, edges are prerequisites

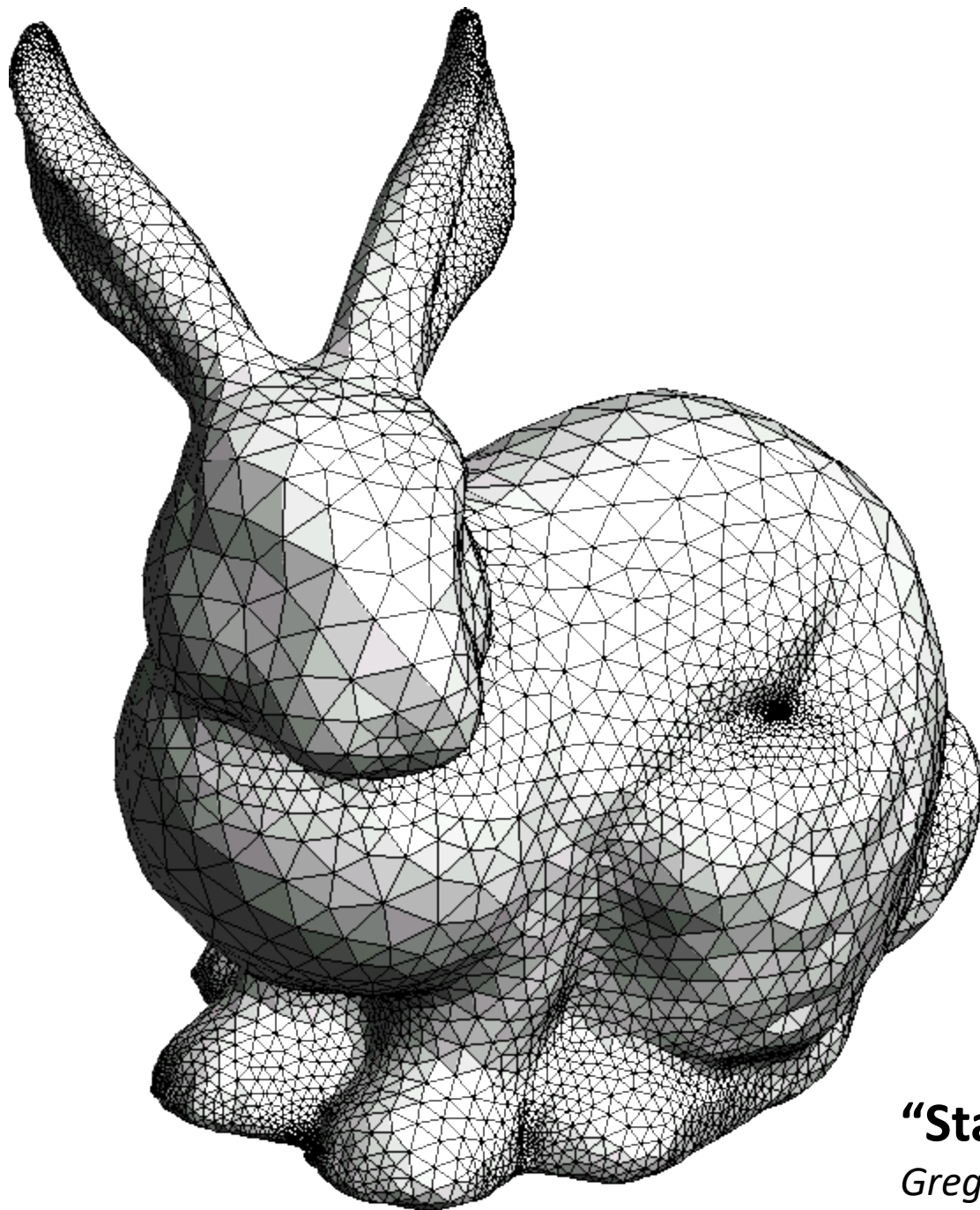
http://waf.cs.illinois.edu/discovery/class_hierarchy_at_illinois/



MP Collaborations in CS 225

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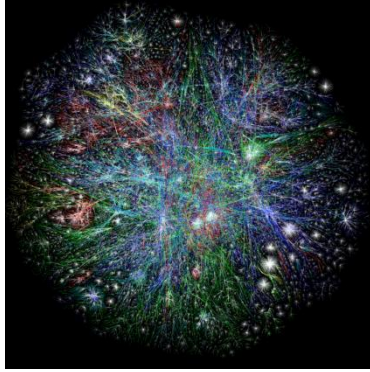
Presented by Cinda Heeren, 2016



“Stanford Bunny”

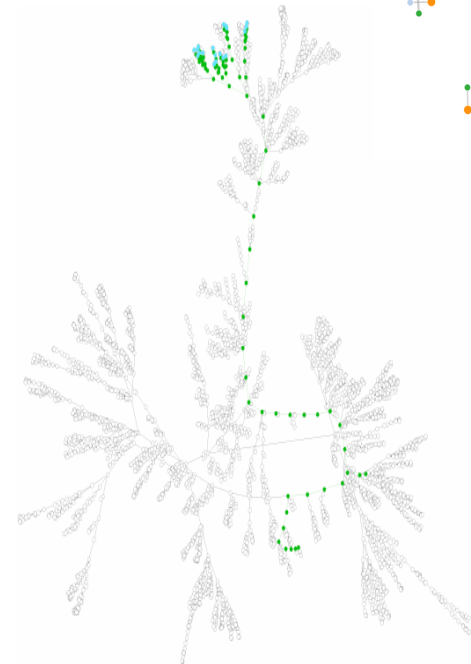
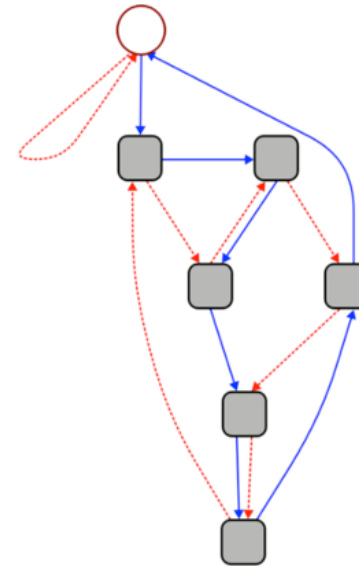
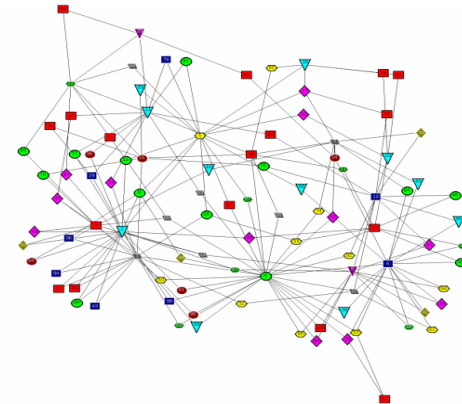
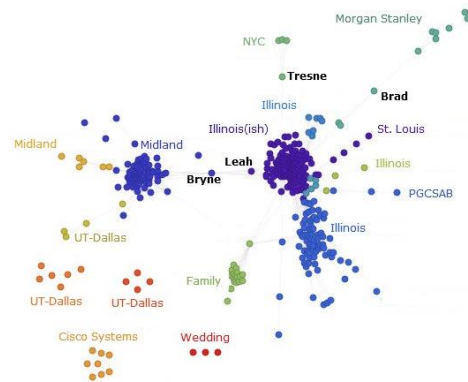
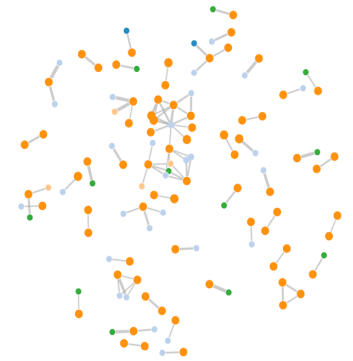
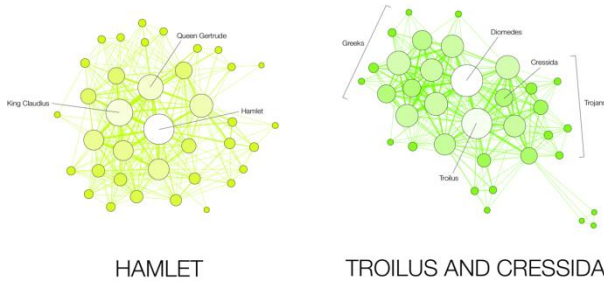
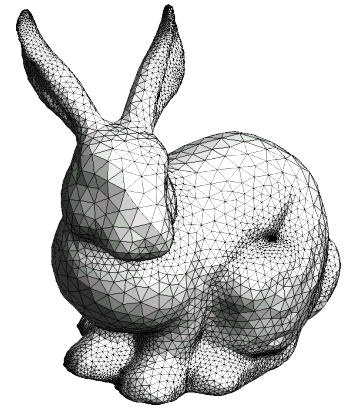
Greg Turk and Mark Levoy (1994)

Graphs



To study all of these structures:

1. A common vocabulary
2. Graph implementations
3. Graph traversals
4. Graph algorithms

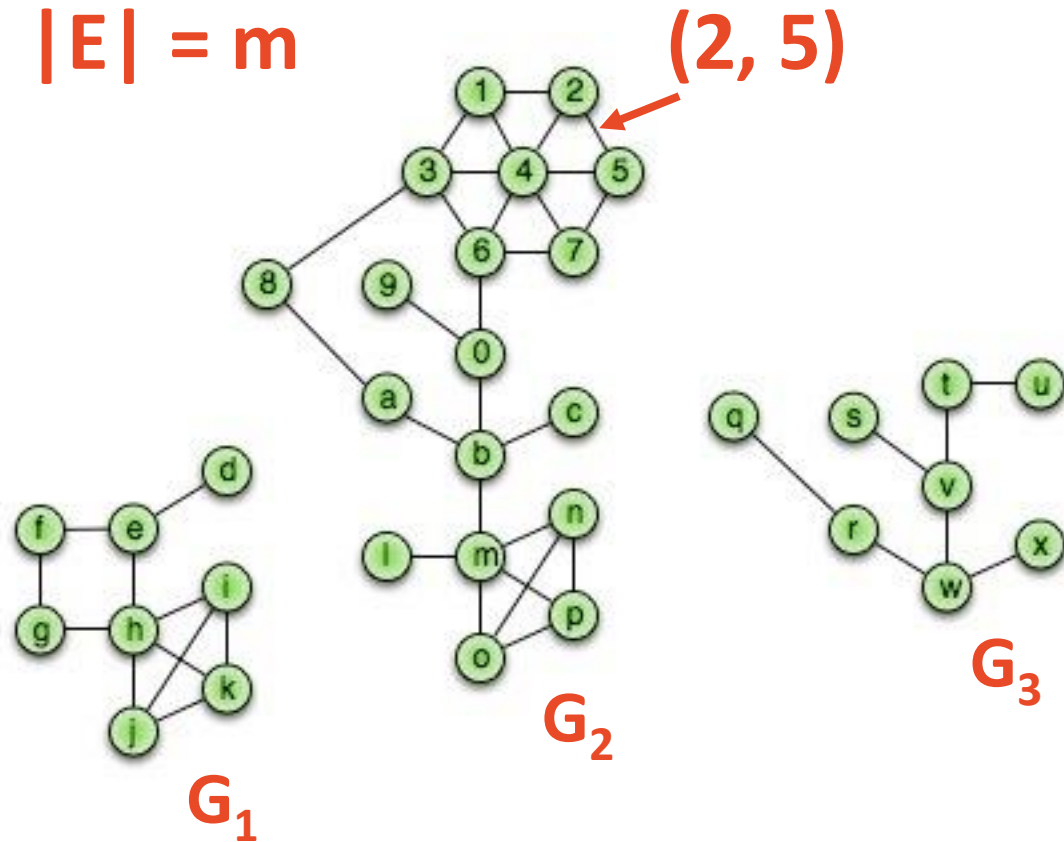


Graph Vocabulary

$$G = (V, E)$$

$$|V| = n$$

$$|E| = m$$



Incident Edges:
 $I(v) = \{ (x, v) \text{ in } E \}$

Degree(v): $|I|$

Adjacent Vertices:
 $A(v) = \{ x : (x, v) \text{ in } E \}$

Path(G_2): Sequence of vertices connected by edges

Cycle(G_1): Path with a common begin and end vertex.

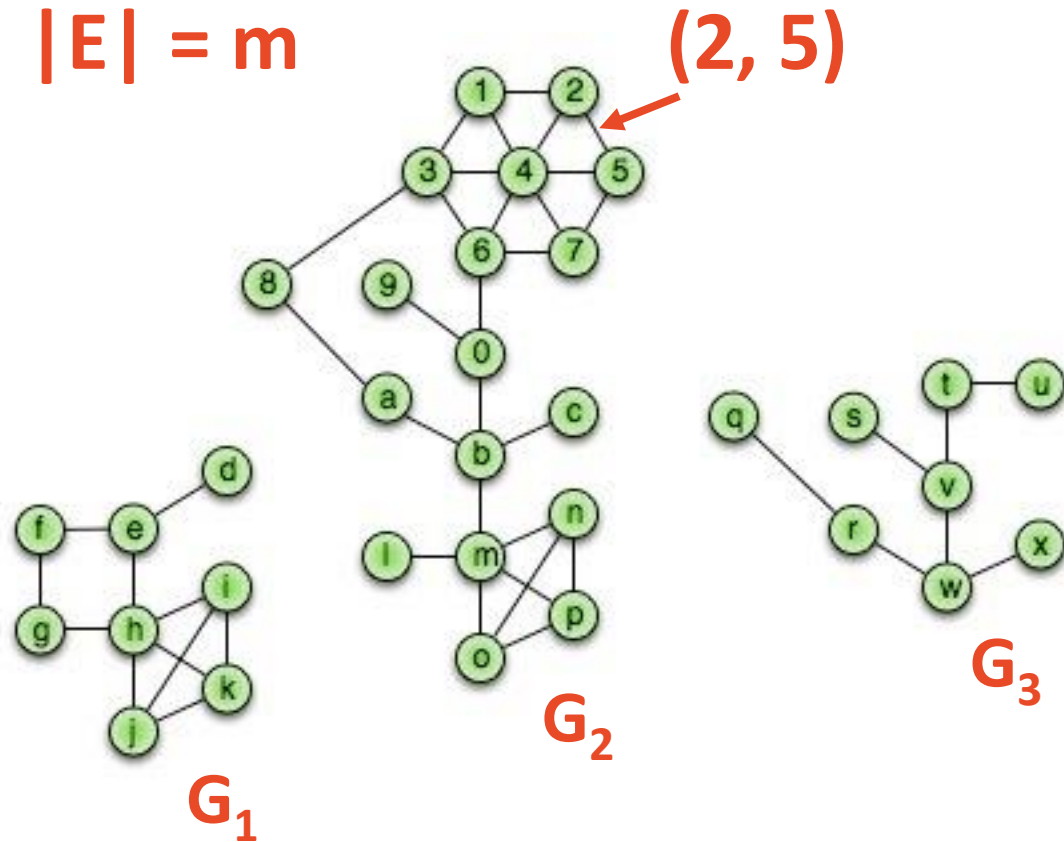
Simple Graph(G): A graph with no self loops or multi-edges.

Graph Vocabulary

$$G = (V, E)$$

$$|V| = n$$

$$|E| = m$$



Subgraph(G):

$$G' = (V', E')$$

$V' \subseteq V, E' \subseteq E$, and

$$(u, v) \in E \rightarrow u \in V', v \in V'$$

Complete subgraph(G)

Connected subgraph(G)

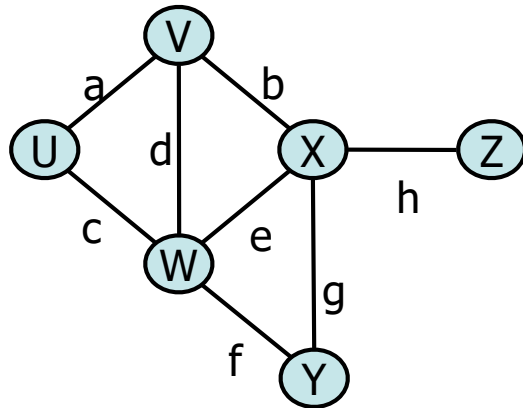
Connected component(G)

Acyclic subgraph(G)

Spanning tree(G)

Running times are often reported by n , the number of vertices, but often depend on m , the number of edges.

How many edges? **Minimum edges:**
Not Connected:



Connected*:

Maximum edges:
Simple:

Not simple:

$$\sum_{v \in V} \deg(v) =$$