

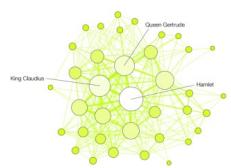
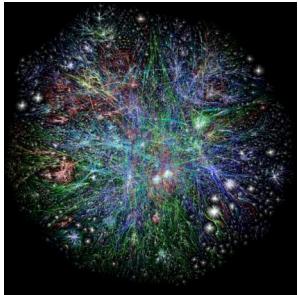


# CS 225

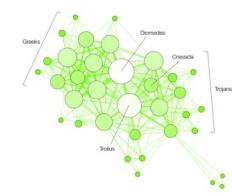
## Data Structures

*April 25 – Dijkstra’s Algorithm  
G Carl Evans*

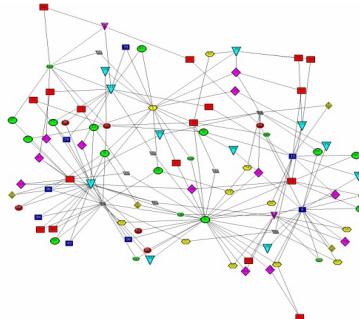
# Graphs



HAMLET

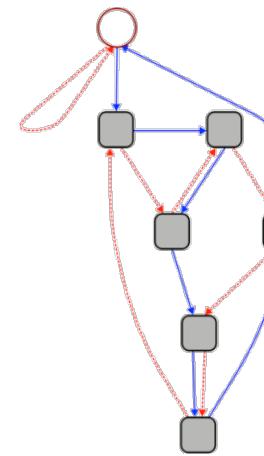
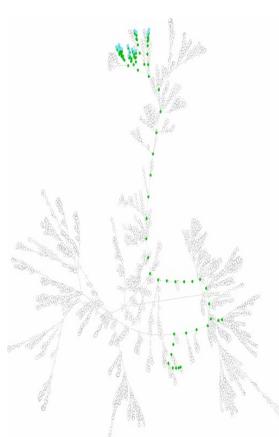


TROILUS AND CRESSIDA



To study all of these structures:

1. A common vocabulary
2. Graph implementations
3. Graph traversals
4. Graph algorithms

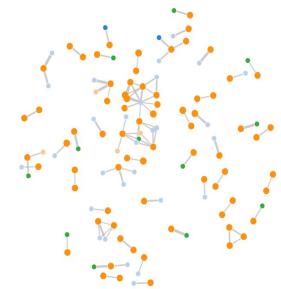
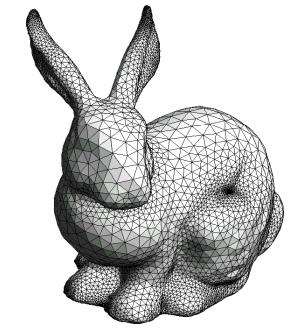


```
heapsyUp(list*, unsigned int):  
    push  
    rbp  
    mov  
    rbp, rsp  
    sub  
    rbp, 16  
    mov  
    dword ptr [rbp - 8], rdi  
    mov  
    dword ptr [rbp - 12], rsi  
    mov  
    rbp  
    jne .LBB0_4
```

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    rbp, rsp  
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    mov  
    dword ptr [rbp - 12], rsi  
    mov  
    rbp  
    jne .LBB0_4
```

```
.LBB0_4:  
    add  
    rbp, rbp  
    16  
    pop  
    rbp  
    ret
```



## MST Algorithm Runtime:

We know that MSTs are always run on a minimally connected graph:

$$n-1 \leq m \leq n(n-1) / 2$$

$$O(n) \leq O(m) \leq O(n^2)$$

## MST Algorithm Runtime:

- Kruskal's Algorithm:  
 $O(n + m \lg(n))$

Sparse Graph:

Dense Graph:

- Prim's Algorithm:  
 $O(n \lg(n) + m \lg(n))$

Sparse Graph:

Dense Graph:

# Suppose I have a new heap:

	Binary Heap	Fibonacci Heap
Remove Min	$O(\lg(n))$	$O(\lg(n))$
Decrease Key	$O(\lg(n))$	$O(1)^*$

## What's the updated running time?

```
 6  PrimMST(G, s):
 7      foreach (Vertex v : G):
 8          d[v] = +inf
 9          p[v] = NULL
10         d[s] = 0
11
12         PriorityQueue Q // min distance, defined by d[v]
13         Q.buildHeap(G.vertices())
14         Graph T           // "labeled set"
15
16         repeat n times:
17             Vertex m = Q.removeMin()
18             T.add(m)
19             foreach (Vertex v : neighbors of m not in T):
20                 if cost(v, m) < d[v]:
21                     d[v] = cost(v, m)
22                     p[v] = m
```

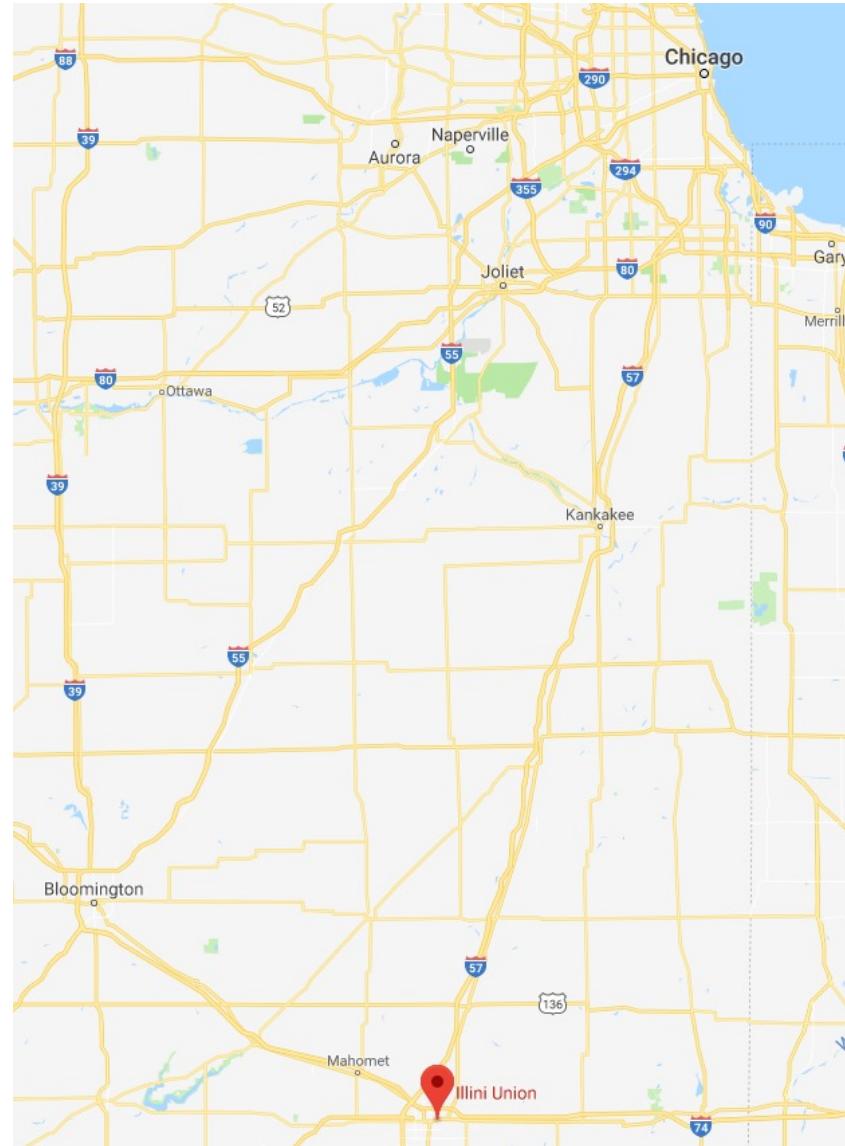
## MST Algorithm Runtimes:

- Kruskal's Algorithm:  
 $O(m \lg(n))$
- Prim's Algorithm:  
 $O(n \lg(n) + m \lg(n))$

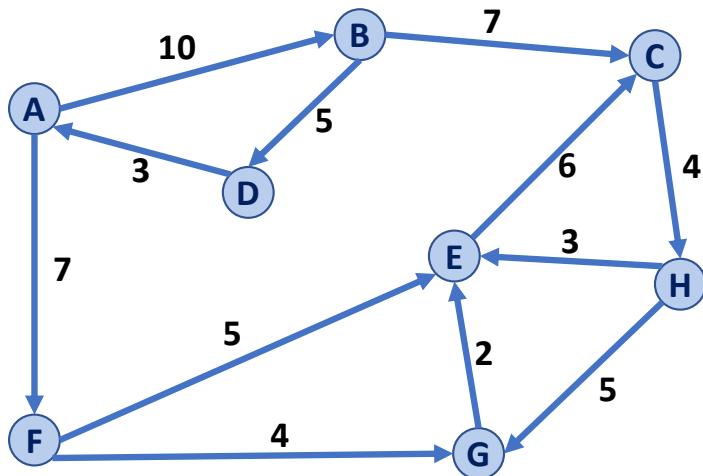
## Final Big-O MST Algorithm Runtimes:

- Kruskal's Algorithm:  
 $O(m \lg(n))$
- Prim's Algorithm:  
 $O(n \lg(n) + m)$

# Shortest Path



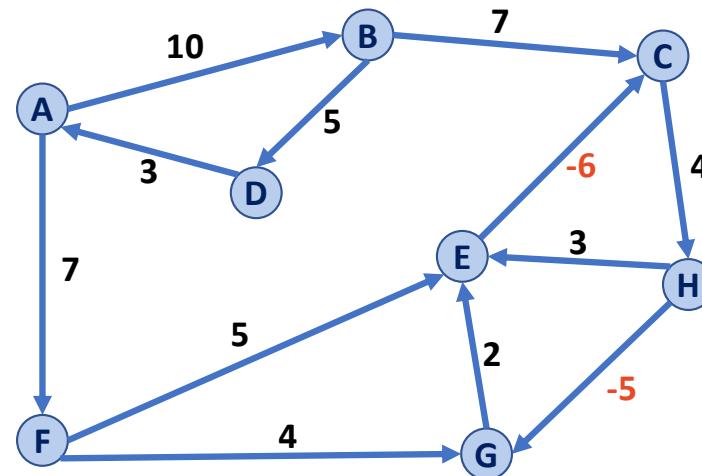
# Dijkstra's Algorithm (SSSP)



```
6  DijkstraSSSP(G, s) :  
7      foreach (Vertex v : G) :  
8          d[v] = +inf  
9          p[v] = NULL  
10         d[s] = 0  
11  
12         PriorityQueue Q // min distance, defined by d[v]  
13         Q.buildHeap(G.vertices())  
14         Graph T           // "labeled set"  
15  
16         repeat n times:  
17             Vertex u = Q.removeMin()  
18             T.add(u)  
19             foreach (Vertex v : neighbors of u not in T) :  
20                 if _____ < d[v] :  
21                     d[v] = _____  
                         p[v] = m
```

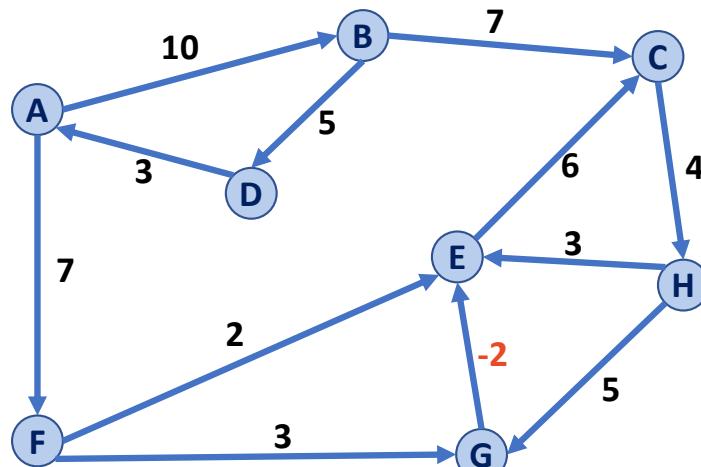
# Dijkstra's Algorithm (SSSP)

What about negative weight cycles?



# Dijkstra's Algorithm (SSSP)

What about negative weight edges, without negative weight cycles?



# Dijkstra's Algorithm (SSSP)

What is the running time?

```
    DijkstraSSSP(G, s):
6      foreach (Vertex v : G):
7        d[v] = +inf
8        p[v] = NULL
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11       PriorityQueue Q // min distance, defined by d[v]
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