String Algorithms and Data Structures
Burrows-Wheeler Transform

CS 199-225
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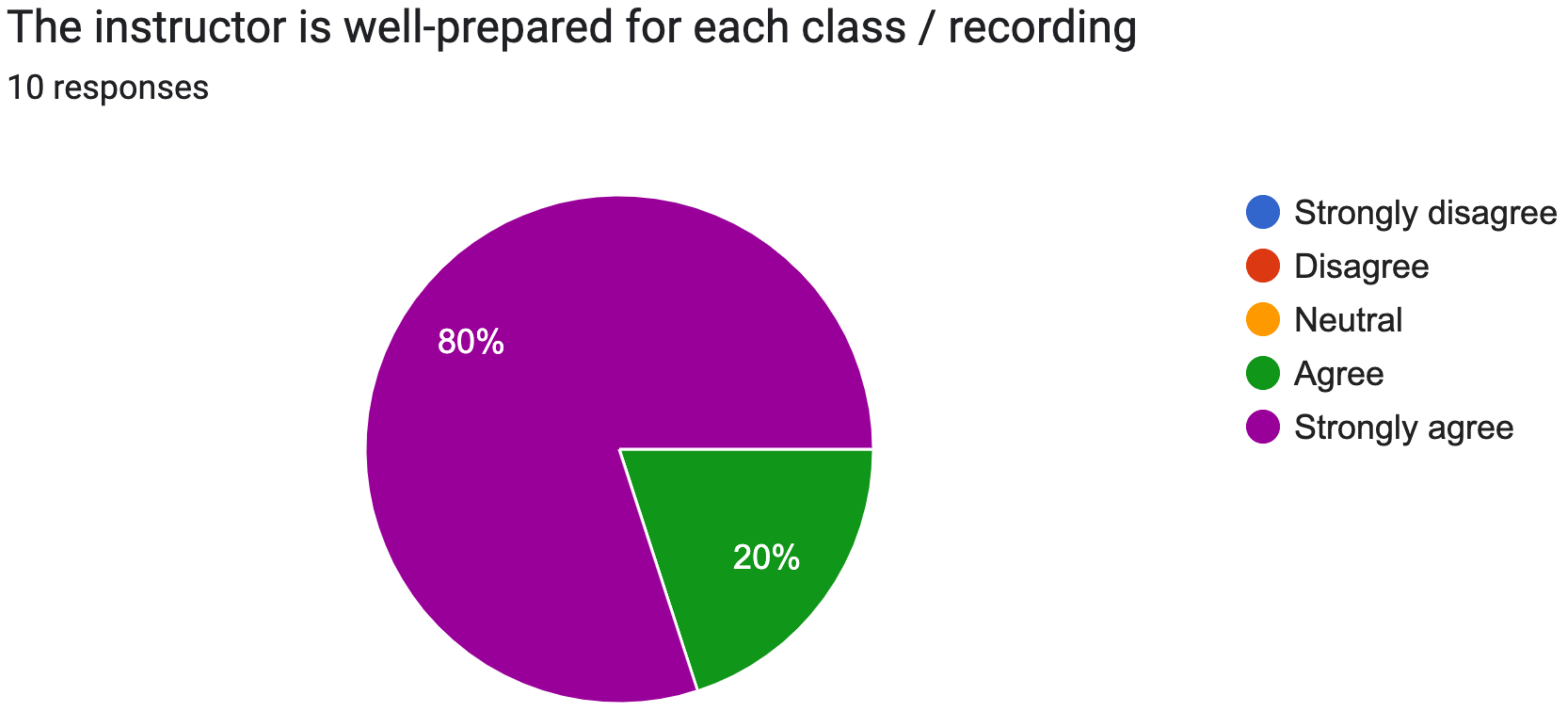
March 27, 2023

UNIVERSITY OF ILLINOIS URBANA-CHAMPAIGN
Department of Computer Science

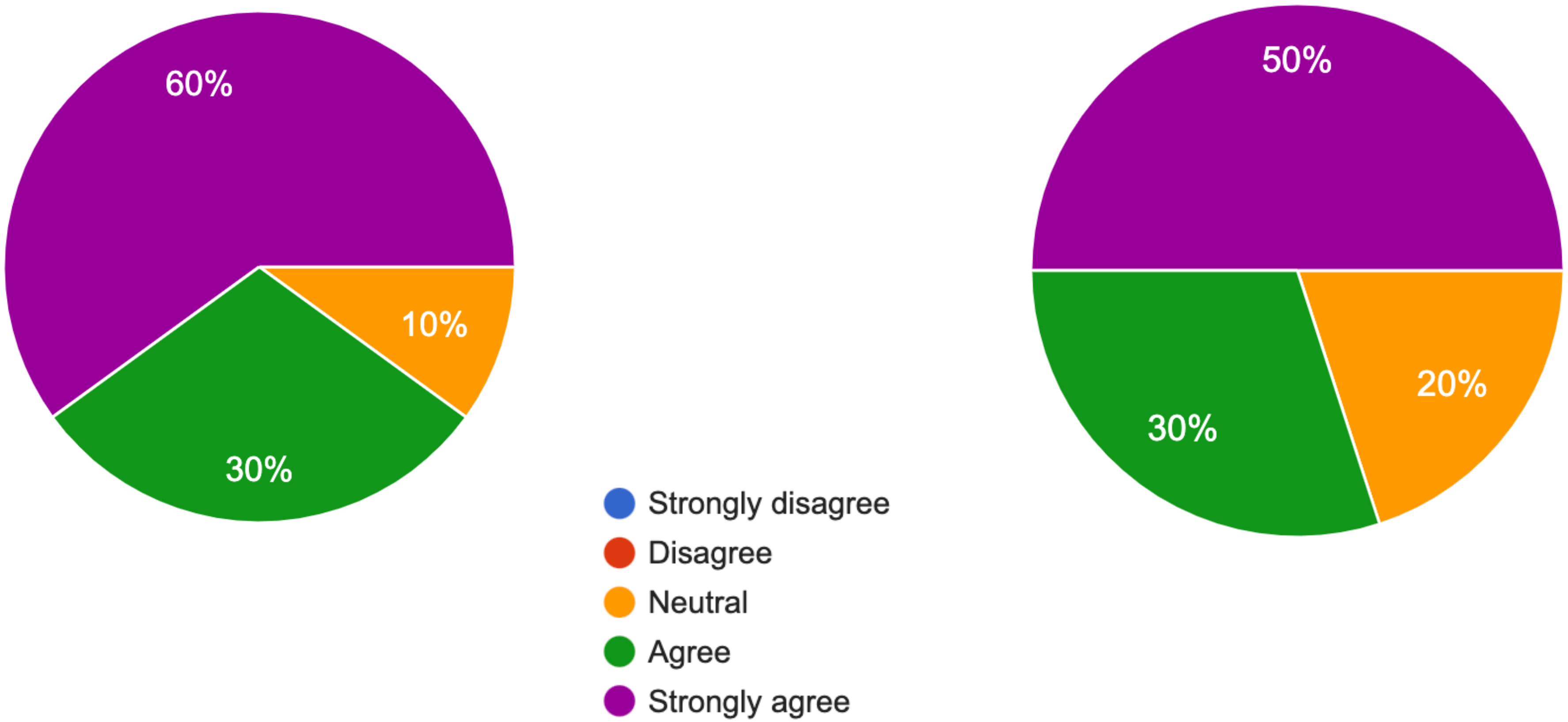
Informal Early Feedback



Informal Early Feedback

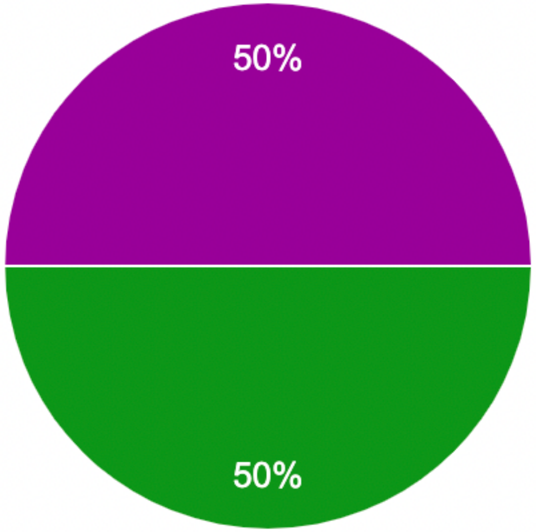
I feel that I can actively participate in lecture
10 responses

I feel that I can actively participate in class in general
10 responses

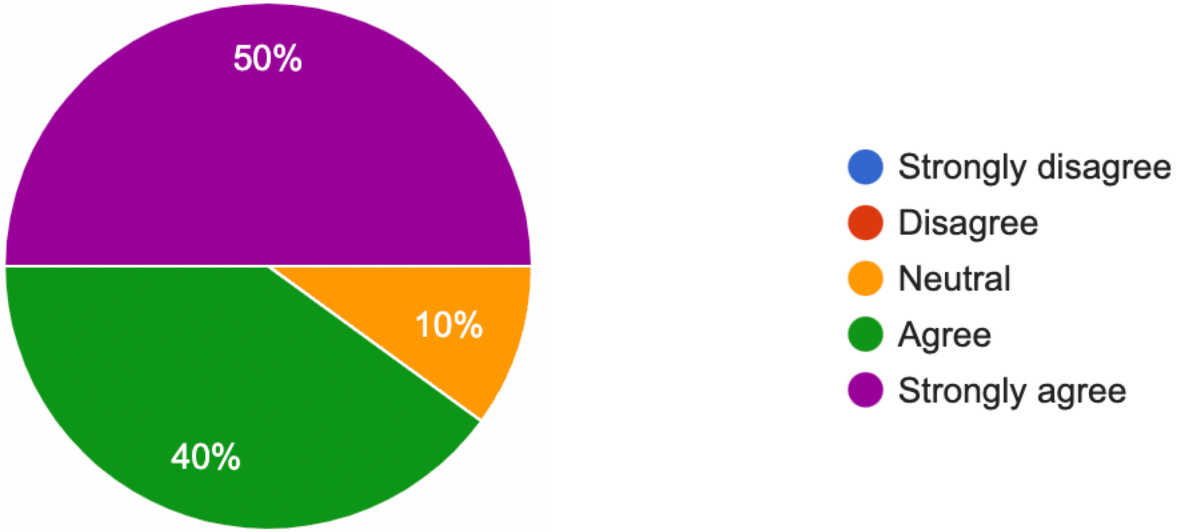


Informal Early Feedback

I receive helpful and complete answers to my questions

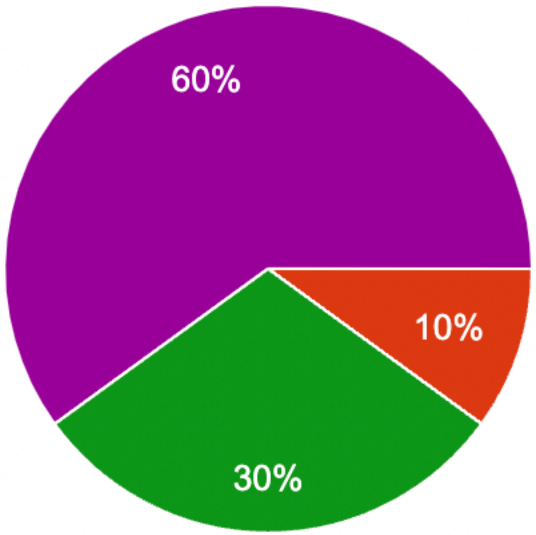


During lecture

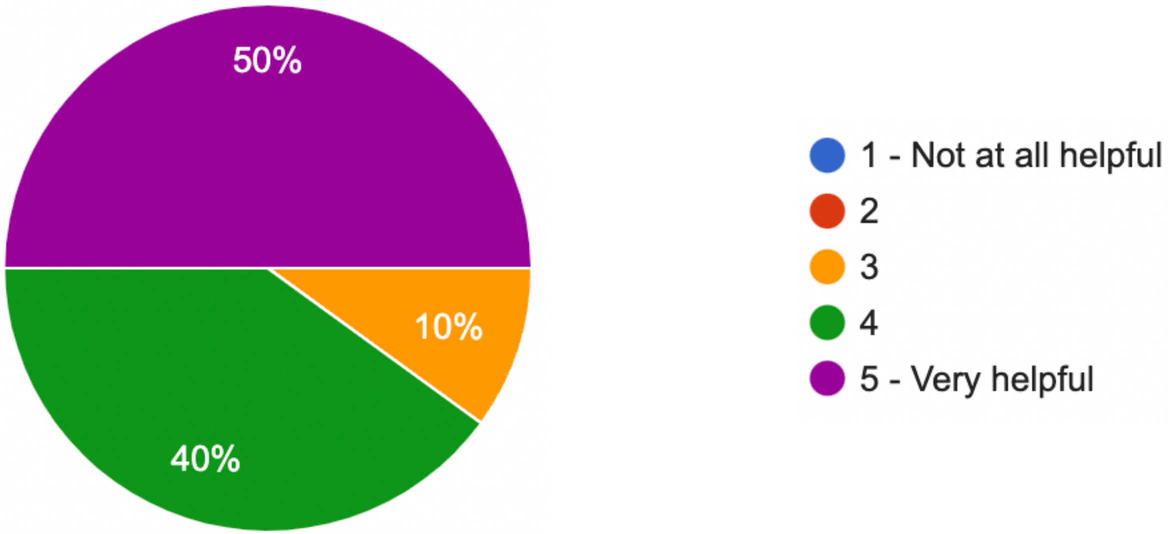


Outside lecture

Informal Early Feedback



Lecture helpfulness



Assignment helpfulness

Informal Early Feedback

I think it was very useful in learning no necessarily string algorithms themselves, but algorithms in general. It feels like a **good introduction to how to optimize a seemingly simple naive problem.**

Learning different interesting algorithms for operation on string help me to have better ability to handle problems.

I do enjoy having a highlight reel of string matching algorithms tailored to different use cases with **an opportunity to try my hand at implementing them.**

I am not sure what g**roup activity could be added** but something that allows us to interact with other students more.

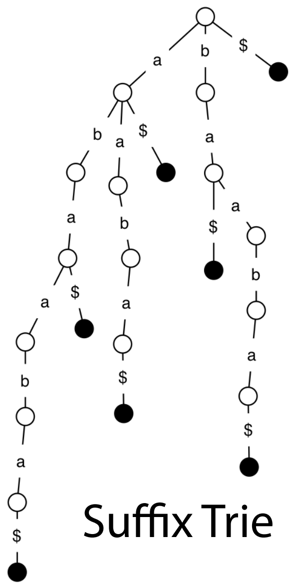
not just string algorithm? challenging problems for bonus points?

I wouldn't mind **having some segments that goes deeper into the material and theory covered in the parent CS 225 course** e.g. fuller implementations of DSAs and their associated algorithms. But string matching is pretty cool as well!

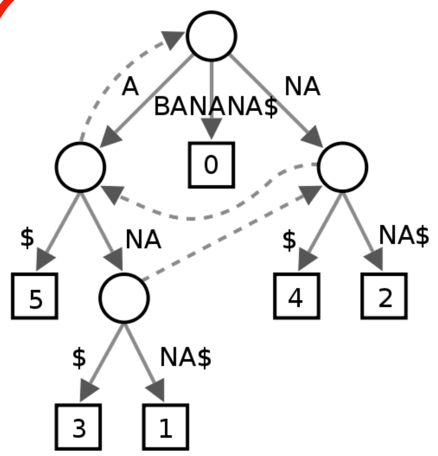
Exact pattern matching *w/ indexing*

There are many data structures built on ***suffixes***

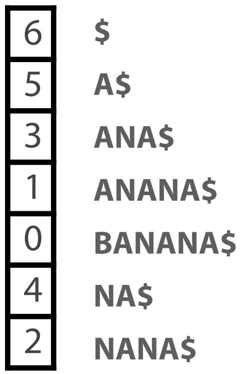
We have now seen both of these data structures



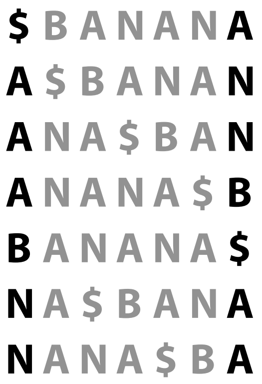
Suffix Trie



Suffix Tree



Suffix Array



FM Index

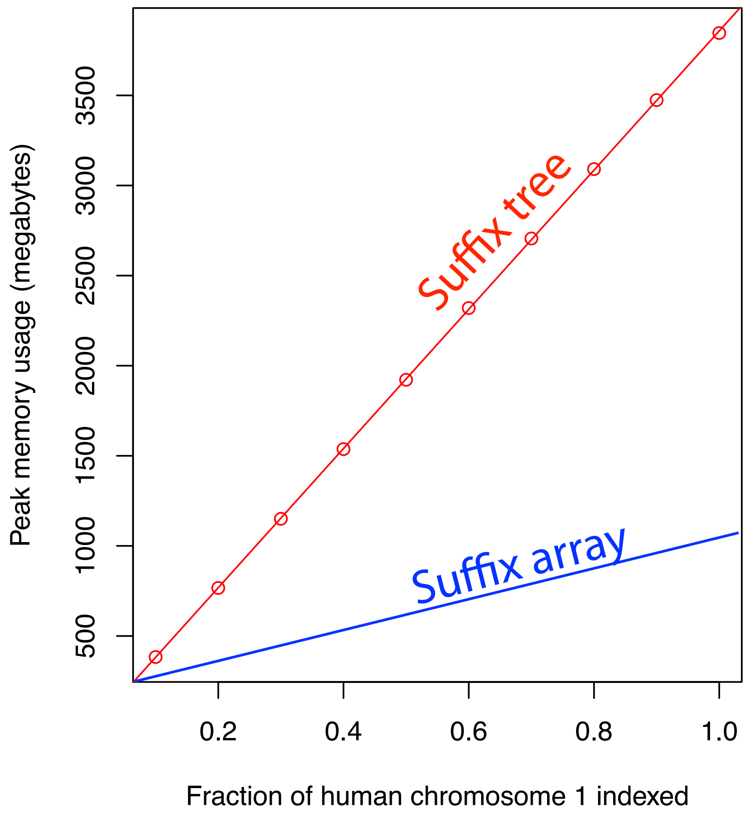
Exact pattern matching *w/ indexing*

	Suffix tree	**Suffix array**
Time: Does P occur?		
Time: Report k locations of P		
Space		

$m = |T|,\ n = |P|,\ k = \#$ occurrences of P in T

Suffix tree vs suffix array: size

The suffix array has a smaller constant factor than the tree



Suffix tree: ~16 bytes per character

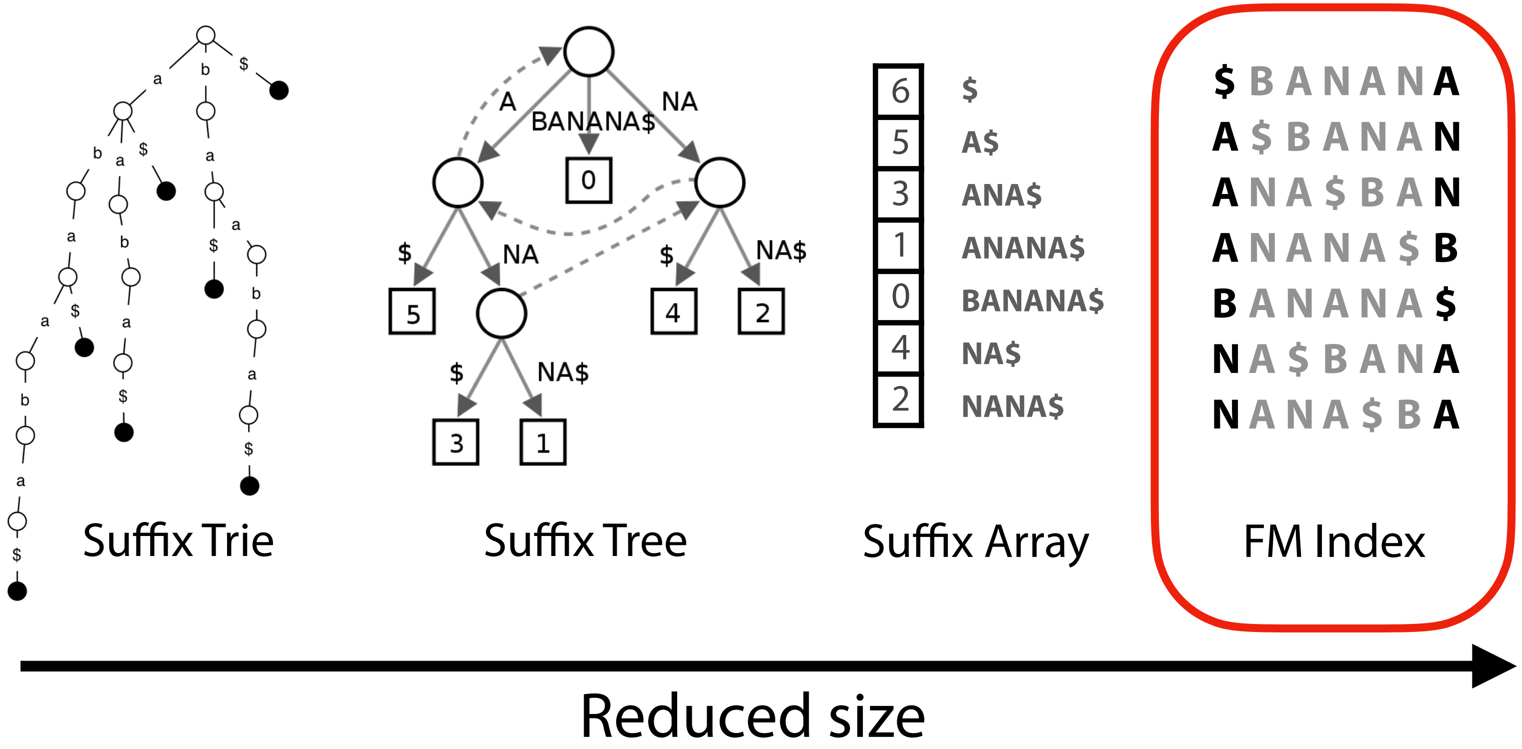
Suffix array: ~4 bytes per character

Raw text: 2 bits per character

Exact pattern matching *w/ indexing*

There are many data structures built on ***suffixes***

The FM index is a compressed self-index (smaller* than original text)!



Exact pattern matching *w/ indexing*

The basis of the FM index is a *transformation*

B A N A N A $

A N N B $ A A

Burrows-Wheeler Transform

Reversible permutation of the characters of a string

T		BWT(T)
B A N A N A $	⟷	A N N B $ A A

1) How to encode?

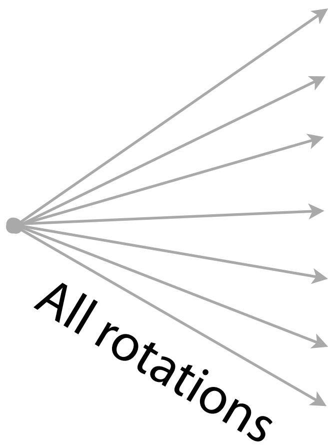
2) How to decode?

3) How is it useful for search?

Burrows-Wheeler Transform

Reversible permutation of the characters of a string

a b a a b a $
T



???

Burrows M, Wheeler DJ: A block sorting lossless data compression algorithm. *Digital Equipment Corporation, Palo Alto, CA* 1994, Technical Report 124; 1994

Text rotations

A string is a ‘rotation’ of another string if it can be reached by wrap-around shifting the characters

a b c d e f $

b c d e f $ a

c d e f $ a b

d e f $ a b c

e f $ a b c d

f $ a b c d e

$ a b c d e f

(after this they repeat)

Text Rotations

A string is a ‘rotation’ of another string if it can be reached by wrap-around shifting the characters

Which of these are rotations of ‘ABCD’?

A) BCDA

B) BACD

C) DCAB

D) CDAB

Burrows-Wheeler Transform

Reversible permutation of the characters of a string

a b a a b a $
T

All rotations

a b a a b a $
b a a b a $ a
a a b a $ a b
a b a $ a b a
b a $ a b a a
a $ a b a a b
$ a b a a b a

(after this they repeat)

Burrows M, Wheeler DJ: A block sorting lossless data compression algorithm. *Digital Equipment Corporation, Palo Alto, CA* 1994, Technical Report 124; 1994

Burrows-Wheeler Transform

Reversible permutation of the characters of a string

a b a a b a $

T

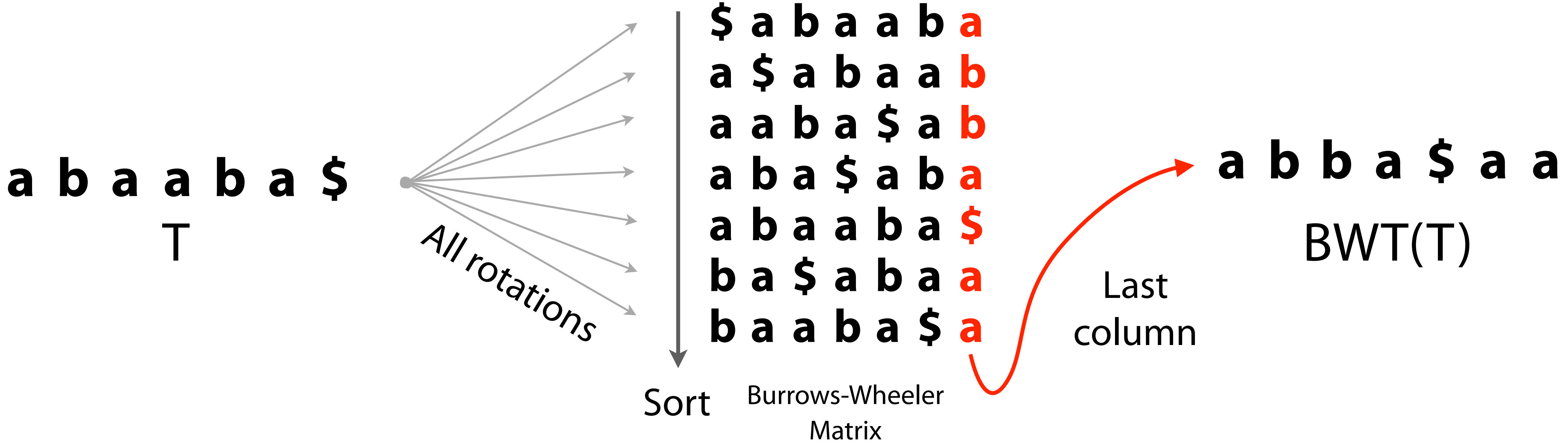
All rotations

a b a a b a $
$ a b a a b a
a $ a b a a b
b a $ a b a a
a b a $ a b a
a a b a $ a b
b a a b a $ a

Burrows M, Wheeler DJ: A block sorting lossless data compression algorithm. *Digital Equipment Corporation, Palo Alto, CA* 1994, Technical Report 124; 1994

Burrows-Wheeler Transform

Reversible permutation of the characters of a string



Burrows M, Wheeler DJ: A block sorting lossless data compression algorithm. *Digital Equipment Corporation, Palo Alto, CA* 1994, Technical Report 124; 1994

Burrows-Wheeler Transform

(1) Build all rotations

(2) Sort all rotations

(3) Take last column

T = **c a r $**

Burrows-Wheeler Transform

(1) Build all rotations

(2) Sort all rotations

(3) Take last column

T = **c a r $** $\xrightarrow{\text{All rotations}}$

$	**c**	**a**	**r**
a	**r**	**$**	**c**
c	**a**	**r**	**$**
r	**$**	**c**	**a**

Sort

$\xrightarrow{\text{Last column}}$ **r c $ a**

Assignment 8: a_bwt

Learning Objective:

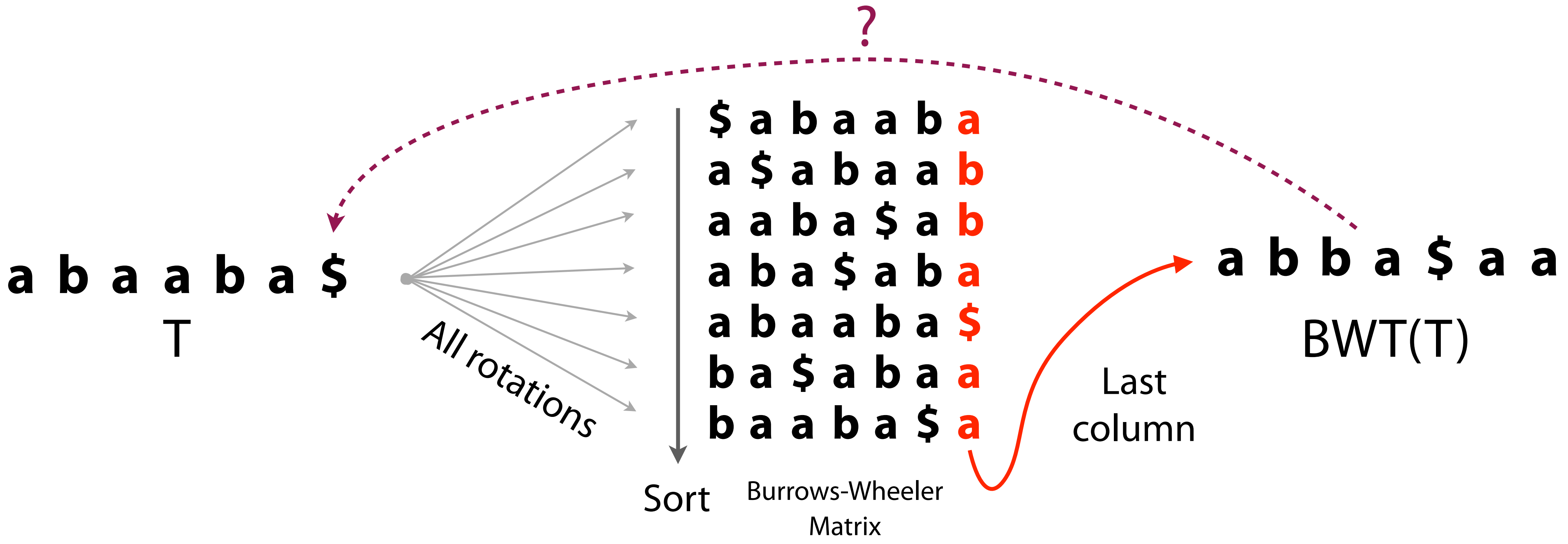
Implement the Burrows-Wheeler Transform on text

Reverse the Burrows-Wheeler Transform to reproduce text

Consider: How can the BWT be stored *smaller* than the original text?

Burrows-Wheeler Transform

How to reverse the BWT?



Burrows-Wheeler Transform

BWT(T) = **r c $ a** T = **c a r $**

Burrows-Wheeler Transform

BWT(T) = **r c $ a**   T = **c a r $**

1) Prepend the BWT as a column

2) Sort the full matrix as rows

3) Repeat 1 and 2 until full matrix

4) Pick the row ending in '$'

$$\begin{matrix} \$ \\ a \\ c \\ r \end{matrix} \Rightarrow \begin{matrix} \$ & c \\ a & r \\ c & a \\ r & \$ \end{matrix} \Rightarrow \begin{matrix} \$ & c & a \\ a & r & \$ \\ c & a & r \\ r & \$ & c \end{matrix} \Rightarrow \begin{matrix} \$ & c & a & r \\ a & r & \$ & c \\ c & a & r & \$ \\ r & \$ & c & a \end{matrix}$$

Burrows-Wheeler Transform

BWT(T) = **r c $ a** T = **c a r $**

$	c	a	**r**	**$**
a	r	$	**c**	**a**
c	a	r	**$**	**c**
r	$	c	**a**	**r**

Burrows-Wheeler Transform

BWT(T) = **r c $ a** T = **c a r $**

$	c	a	**r**
a	r	$	**c**
c	a	r	**$**
r	$	c	**a**

$	**c**
a	**r**
c	**a**
r	**$**

Burrows-Wheeler Transform

BWT(T) = **r c $ a**   T = **c a r $**

$	c	a	**r**
a	r	$	**c**
c	a	r	**$**
r	$	c	**a**

$	**c**	**a**
a	**r**	**$**
c	**a**	**r**
r	**$**	**c**

Burrows-Wheeler Transform

What is the right context of **a p $\boxed{p}$ l e $** ? **l e $ a p**

A letter always has the same right context.

$	**a**	**p**	**p**	**l**	**e**
a	**p**	**p**	**l**	**e**	**$**
e	**$**	**a**	**p**	**p**	**l**
l	**e**	**$**	**a**	**p**	**p**
p	**l**	**e**	**$**	**a**	**p**
p	**p**	**l**	**e**	**$**	**a**

Burrows-Wheeler Transform: T-ranking

To continue, we have to be able to uniquely identify each character in our text.

Give each character in *T* a rank, equal to # times the character occurred previously in *T*. Call this the *T-ranking*.

a b a a b a $

Ranks aren't explicitly stored; they are just for illustration

Burrows-Wheeler Transform

BWM with T-ranking:

F						L
$\$$	a_0	b_0	a_1	a_2	b_1	a_3
a_3	$\$$	a_0	b_0	a_1	a_2	b_1
a_1	a_2	b_1	a_3	$\$$	a_0	b_0
a_2	b_1	a_3	$\$$	a_0	b_0	a_1
a_0	b_0	a_1	a_2	b_1	a_3	$\$$
b_1	a_3	$\$$	a_0	b_0	a_1	a_2
b_0	a_1	a_2	b_1	a_3	$\$$	a_0

Look at first and last columns, called *F* and *L* (and look at just the **a**s)

as occur in the same order in *F* and *L*. As we look down columns, in both cases we see: $\mathbf{a}_3$, $\mathbf{a}_1$, $\mathbf{a}_2$, $\mathbf{a}_0$

Burrows-Wheeler Transform

BWM with T-ranking:

F						*L*
\$	a_0	b_0	a_1	a_2	b_1	a_3
a_3	\$	a_0	b_0	a_1	a_2	$\mathbf{b}_1$
a_1	a_2	b_1	a_3	\$	a_0	$\mathbf{b}_0$
a_2	b_1	a_3	\$	a_0	b_0	a_1
a_0	b_0	a_1	a_2	b_1	a_3	\$
$\mathbf{b}_1$	a_3	\$	a_0	b_0	a_1	a_2
$\mathbf{b}_0$	a_1	a_2	b_1	a_3	\$	a_0

Same with **b**s: $\mathbf{b}_1$, $\mathbf{b}_0$

Burrows-Wheeler Transform: LF Mapping

BWM with T-ranking:

F						*L*
\$	a_0	b_0	a_1	a_2	b_1	$\mathbf{a_3}$
$\mathbf{a_3}$	\$	a_0	b_0	a_1	a_2	$\mathbf{b_1}$
$\mathbf{a_1}$	a_2	b_1	a_3	\$	a_0	$\mathbf{b_0}$
$\mathbf{a_2}$	b_1	a_3	\$	a_0	b_0	$\mathbf{a_1}$
$\mathbf{a_0}$	b_0	a_1	a_2	b_1	a_3	**\$**
$\mathbf{b_1}$	a_3	\$	a_0	b_0	a_1	$\mathbf{a_2}$
$\mathbf{b_0}$	a_1	a_2	b_1	a_3	\$	$\mathbf{a_0}$

LF Mapping: The i^{th} occurrence of a character c in L and the i^{th} occurrence of c in F correspond to the *same* occurrence in T (i.e. have same rank)

Burrows-Wheeler Transform: LF Mapping

Why does this work?

Right context:
a b a $ a b

$ a b a a b a
a $ a b a a b
a a b a $ a b
a b a $ a b a
a b a a b a $
b a $ a b a a
b a a b a $ a

Right context:
a b a $ a b

These characters have the same right contexts!

These characters *are the same character!* $\mathbf{a}_0\ \mathbf{b}_0\ \mathbf{a}_1\ \mathbf{a}_2\ \mathbf{b}_1\ \mathbf{a}_3\ \$$

Burrows-Wheeler Transform: LF Mapping

Why does this work?

Why are these **a**s in this order relative to each other?

$\$\ a\ b\ a\ a\ b\ a_3$
$a_3\ \$\ a\ b\ a\ a\ b_1$
$a_1\ a\ b\ a\ \$\ a\ b_0$
$a_2\ b\ a\ \$\ a\ b\ a_1$
$a_0\ b\ a\ a\ b\ a\ \$$
$b_1\ a\ \$\ a\ b\ a\ a_2$
$b_0\ a\ a\ b\ a\ \$\ a_0$

They're sorted by right-context

$\$\ a\ b\ a\ a\ b\ a_3$
$a_3\ \$\ a\ b\ a\ a\ b_1$
$a_1\ a\ b\ a\ \$\ a\ b_0$
$a_2\ b\ a\ \$\ a\ b\ a_1$
$a_0\ b\ a\ a\ b\ a\ \$$
$b_1\ a\ \$\ a\ b\ a\ a_2$
$b_0\ a\ a\ b\ a\ \$\ a_0$

They're sorted by right-context

Why are these **a**s in this order relative to each other?

Occurrences of *c* in *F* are sorted by right-context. Same for *L*!

Any ranking we give to characters in *T* will match in *F* and *L*

Burrows-Wheeler Transform: LF Mapping

LF Mapping can be used to recover our original text too!

Given BWT = $\mathbf{a}_3\ \mathbf{b}_1\ \mathbf{b}_0\ \mathbf{a}_1\ \$\ \mathbf{a}_2\ \mathbf{a}_0$

What is L?

What is F?

Burrows-Wheeler Transform: LF Mapping

LF Mapping can be used to recover our original text too!

Start in first row. *F* must have **$**.
L contains character just prior to **$**: $\mathbf{a}_3$

Jump to row *beginning* with $\mathbf{a}_3$.
L contains character just prior to $\mathbf{a}_3$: $\mathbf{b}_1$.

Repeat for $\mathbf{b}_1$, get $\mathbf{a}_2$

Repeat for $\mathbf{a}_2$, get $\mathbf{a}_1$

Repeat for $\mathbf{a}_1$, get $\mathbf{b}_0$

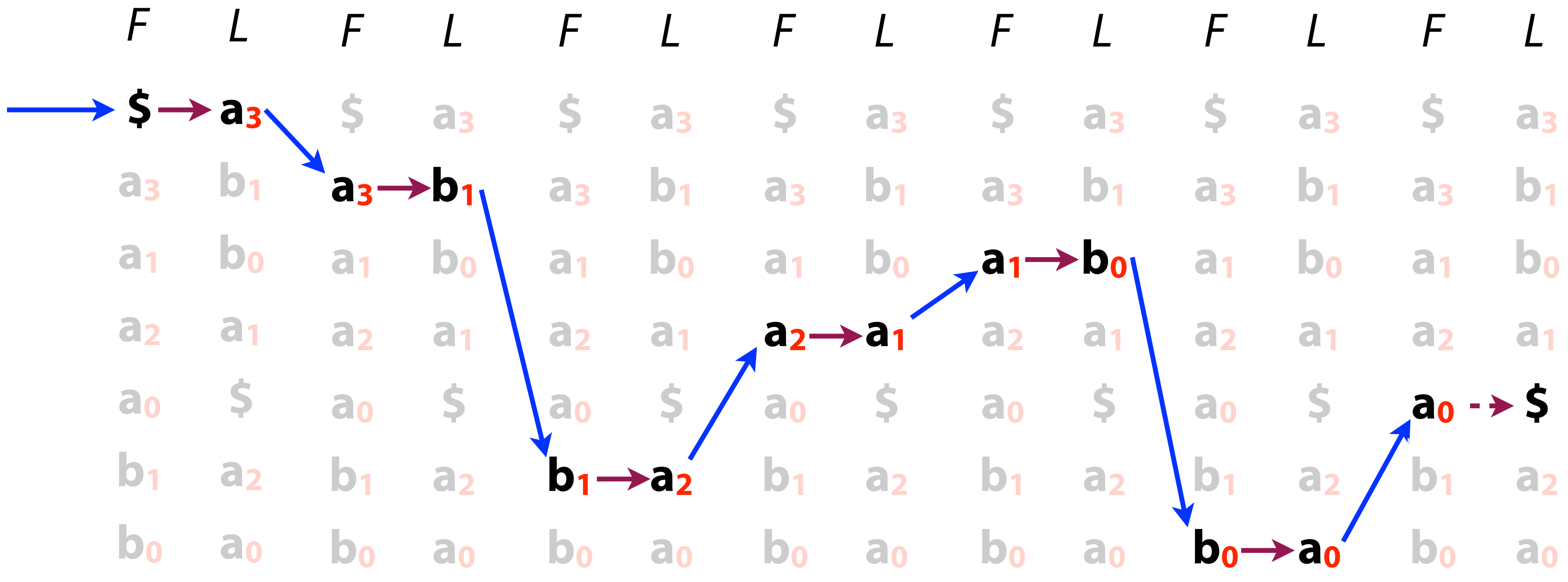
Repeat for $\mathbf{b}_0$, get $\mathbf{a}_0$

Repeat for $\mathbf{a}_0$, get **$** (done)

F	*L*
$	$\mathbf{a}_3$
$\mathbf{a}_3$	$\mathbf{b}_1$
$\mathbf{a}_1$	$\mathbf{b}_0$
$\mathbf{a}_2$	$\mathbf{a}_1$
$\mathbf{a}_0$	**$**
$\mathbf{b}_1$	$\mathbf{a}_2$
$\mathbf{b}_0$	$\mathbf{a}_0$

Burrows-Wheeler Transform: LF Mapping

Another way to visualize:



T: $a_0\ b_0\ a_1\ a_2\ b_1\ a_3\ \$$

Assignment 8: a_bwt

Learning Objective:

Implement the Burrows-Wheeler Transform on text

Reverse the Burrows-Wheeler Transform to reproduce text

Consider: You can use either LF mapping or prepend-sort to reverse. Which do you think would be easier to implement (or more efficient)?

Burrows-Wheeler Transform: A better ranking

Any ranking we give to characters in *T* will match in *F* and *L*

T-Rank: Order by T

F	*L*
$	a_3
a_3	b_1
a_1	b_0
a_2	a_1
a_0	$
b_1	a_2
b_0	a_0

F-Rank: Order by F

F	*L*
$	a_0
a_0	b_0
a_1	b_1
a_2	a_1
a_3	$
b_1	a_2
b_0	a_3

What is good about f-rank?

Burrows-Wheeler Transform: A better ranking

T = **a b b c c d $**

What is the BWM index for my first instance of C? **(C_0)** [0-base for answer]

F						*L*
$	a	b	b	c	c	**d**
a	b	b	c	c	d	**$**
b	b	c	c	d	$	**a**
b	c	c	d	$	a	**b**
c	c	d	$	a	b	**b**
c	d	$	a	b	b	**c**
d	$	a	b	b	c	**c**

Burrows-Wheeler Transform: A better ranking

Say *T* has 300 **A**s, 400 **C**s, 250 **G**s and 700 **T**s and **$** < **A** < **C** < **G** < **T**

What is the BWM index for my 100th instance of G? **(G_{99})** [0-base for answer]

Skip row starting with **$** (1 row)
Skip rows starting with **A** (300 rows)
Skip rows starting with **C** (400 rows)
Skip first 99 rows starting with **G** (99 rows)

Answer: skip 800 rows -> **index 800 contains my 100th G**

With a little preprocessing we can find any character in O(1) time!

FM Index

(Next week's material)

An index combining the BWT *with a few small auxiliary data structures*

Core of index is ***first (F)*** and ***last (L) rows*** from BWM:

F						*L*
$	a	b	a	a	b	**a**
a	$	a	b	a	a	**b**
a	a	b	a	$	a	**b**
a	b	a	$	a	b	**a**
a	b	a	a	b	a	**$**
b	a	$	a	b	a	**a**
b	a	a	b	a	$	**a**

L is the same size as *T*

F can be represented as array of $|\Sigma|$ integers (or not stored at all!)

We're discarding *T* — *we can recover it from L!*

FM Index: Querying

Can we query like the suffix array?

$	a	b	a	a	b	a
a	$	a	b	a	a	b
a	a	b	a	$	a	b
a	b	a	$	a	b	a
a	b	a	a	b	a	$
b	a	$	a	b	a	a
b	a	a	b	a	$	a

6	$
5	a $
2	a a b a $
3	a b a $
0	a b a a b a $
4	b a $
1	b a a b a $

We don't have these columns, and we don't have T.
Binary search not possible.

FM Index: Querying

The BWM is a lot like the suffix array — maybe we can query the same way?

$	a	b	a	a	b	a
a	$	a	b	a	a	b
a	a	b	a	$	a	b
a	b	a	$	a	b	a
a	b	a	a	b	a	$
b	a	$	a	b	a	a
b	a	a	b	a	$	a

BWM(T)

6	$
5	a $
2	a a b a $
3	a b a $
0	a b a a b a $
4	b a $
1	b a a b a $

SA(T)

FM Index: Querying

The BWM is a lot like the suffix array — maybe we can query the same way?

$	a	b	a	a	b	a
a	$	a	b	a	a	b
a	a	b	a	$	a	b
a	b	a	$	a	b	a
a	b	a	a	b	a	$
b	a	$	a	b	a	a
b	a	a	b	a	$	a

6	$
5	a $
2	a a b a $
3	a b a $
0	a b a a b a $
4	b a $
1	b a a b a $

We don't have these columns, and we don't have T.

FM Index: Querying

Look for range of rows of BWM(T) with *P* as prefix

Start with shortest suffix, then match successively longer suffixes

P = **aba**

Easy to find all the rows beginning with **a**

F						*L*
$	a	b	a	a	b	**a_0**
a_0	$	a	b	a	a	**b**
a_1	a	b	a	$	a	**b**
a_2	b	a	$	a	b	**a_1**
a_3	b	a	a	b	a	**$**
b	a	$	a	b	a	**a_2**
b	a	a	b	a	$	**a_3**

FM Index: Querying

We have rows beginning with **a**, now we want rows beginning with **ba**

P = **aba**

F						*L*
\$	a	b	a	a	b	$\mathbf{a}_0$
$\mathbf{a}_0$	\$	a	b	a	a	$\mathbf{b}_0$
$\mathbf{a}_1$	a	b	a	\$	a	$\mathbf{b}_1$
$\mathbf{a}_2$	b	a	\$	a	b	$\mathbf{a}_1$
$\mathbf{a}_3$	b	a	a	b	a	**\$**
$\mathbf{b}_0$	a	\$	a	b	a	$\mathbf{a}_2$
$\mathbf{b}_1$	a	a	b	a	\$	$\mathbf{a}_3$

← Look at those rows in *L*.

$\mathbf{b}_0$, $\mathbf{b}_1$ are **b**s occuring just to left.

P = **aba**

F						*L*
\$	a	b	a	a	b	$\mathbf{a}_0$
$\mathbf{a}_0$	\$	a	b	a	a	$\mathbf{b}_0$
$\mathbf{a}_1$	a	b	a	\$	a	$\mathbf{b}_1$
$\mathbf{a}_2$	b	a	\$	a	b	$\mathbf{a}_1$
$\mathbf{a}_3$	b	a	a	b	a	**\$**
$\mathbf{b}_0$	a	\$	a	b	a	$\mathbf{a}_2$
$\mathbf{b}_1$	a	a	b	a	\$	$\mathbf{a}_3$

Use LF Mapping. Let new range delimit those **b**s →

Note: We still aren't storing the characters in grey, we just know they exist.

FM Index: Querying

We have rows beginning with **ba**, now we seek rows beginning with **aba**

P = **aba**

F						L
\$	a	b	a	a	b	a_0
a_0	\$	a	b	a	a	b_0
a_1	a	b	a	\$	a	b_1
a_2	b	a	\$	a	b	a_1
a_3	b	a	a	b	a	\$
b_0	a	\$	a	b	a	a_2
b_1	a	a	b	a	\$	a_3

← **a_2**, **a_3** occur just to left.

P = **aba**

F						L
\$	a	b	a	a	b	a_0
a_0	\$	a	b	a	a	b_0
a_1	a	b	a	\$	a	b_1
a_2	b	a	\$	a	b	a_1
a_3	b	a	a	b	a	\$
b_0	a	\$	a	b	a	a_2
b_1	a	a	b	a	\$	a_3

Use LF Mapping →

Now we have the rows with prefix **aba**

FM Index: Querying

When *P* does not occur in *T*, we eventually fail to find next character in *L*:

P = **bba**

F						*L*
$	a	b	a	a	b	a_0
a_0	$	a	b	a	a	b_0
a_1	a	b	a	$	a	b_1
a_2	b	a	$	a	b	a_1
a_3	b	a	a	b	a	$
b_0	a	$	a	b	a	a_2
b_1	a	a	b	a	$	a_3

Rows with **ba** prefix

No **b**s!

FM Index: Querying

Problem 1: If we *scan* characters in the last column, that can be slow, $O(m)$

P = **aba**

F						L
$	a	b	a	a	b	a_0
a_0	$	a	b	a	a	b_0
a_1	a	b	a	$	a	b_1
a_2	b	a	$	a	b	a_1
a_3	b	a	a	b	a	$
b_0	a	$	a	b	a	a_2
b_1	a	a	b	a	$	a_3

Scan, looking for **b**s

FM Index: Querying

Problem 2: We don't immediately know *where* the matches are in T...

$P =$ **aba**

Got the same range, [3, 5), we would have got from suffix array

	F						L
	$	a	b	a	a	b	a_0
	a_0	$	a	b	a	a	b_0
	a_1	a	b	a	$	a	b_1
[3, 5)	a_2	b	a	$	a	b	a_1
[3, 5)	a_3	b	a	a	b	a	$
	b_0	a	$	a	b	a	a_2
	b_1	a	a	b	a	$	a_3

Where are the values?

	6	$
	5	a $
	2	a a b a $
[3, 5)	3	a b a $
[3, 5)	0	a b a a b a $
	4	b a $
	1	b a a b a $

Bonus Slides

Burrows-Wheeler Transform

Reversible permutation of the characters of a string

T		BWT(T)
B A N A N A $	⟷	A N N B $ A A

1) How to encode?

2) How to decode?

3) How is it useful for compression?

4) How is it useful for search?

Burrows-Wheeler Transform

```
Tomorrow_and_tomorrow_and_tomorrow
```

```
w$wwdd__nnoooaattTmmmrrrrrrooo__ooo
```

```
It_was_the_best_of_times_it_was_the_worst_of_times$
```

```
s$esttssfftteww_hhmmbootttt_ii__woeeaaressIi_______
```

“bzip”: compression w/ a BWT to better organize text

Burrows-Wheeler Transform

```
orrow_and_tomorrow_and_tomorrow$tom
ow$tomorrow_and_tomorrow_and_tomorr
ow_and_tomorrow$tomorrow_and_tomorr
ow_and_tomorrow_and_tomorrow$tomorr
row$tomorrow_and_tomorrow_and_tomor
row_and_tomorrow$tomorrow_and_tomor
row_and_tomorrow_and_tomorrow$tomor
rrow$tomorrow_and_tomorrow_and_tomo
```

Ordered by the ***context*** to the ***right*** of each character

Burrows-Wheeler Transform

In English (and most languages), the next character in a word is not independent of the previous.

In general, if text structured BWT(T) more compressible

final char (*L*)	sorted rotations
a	n to decompress. It achieves compression
o	n to perform only comparisons to a depth
o	n transformation} This section describes
o	n transformation} We use the example and
o	n treats the right-hand side as the most
a	n tree for each 16 kbyte input block, enc
a	n tree in the output stream, then encodes
i	n turn, set \$L[i]\$ to be the
i	n turn, set \$R[i]\$ to the
o	n unusual data. Like the algorithm of Man
a	n use a single set of probabilities table
e	n using the positions of the suffixes in
i	n value at a given point in the vector \$R
e	n we present modifications that improve t
e	n when the block size is quite large. Ho
i	n which codes that have not been seen in
i	n with \$ch\$ appear in the {\em same order
i	n with \$ch\$. In our exam
o	n with Huffman or arithmetic coding. Bri
o	n with figures given by Bell~\cite{bell}.

Figure 1: Example of sorted rotations. Twenty consecutive rotations from the sorted list of rotations of a version of this paper are shown, together with the final character of each rotation.

Burrows M, Wheeler DJ: A block sorting lossless data compression algorithm. *Digital Equipment Corporation, Palo Alto, CA* 1994, Technical Report 124; 1994

Burrows-Wheeler Transform

Lets compare the SA with the BWT…

T = **a b a a b a $**

6
5
2
3
0
4
1

SA(T)

$ a b a a b a
a $ a b a a b
a a b a $ a b
a b a $ a b a
a b a a b a $
b a $ a b a a
b a a b a $ a

BWM(T)

Suffix Array is O(m)

Burrows-Wheeler Transform

Lets compare the SA with the BWT…

T = **a b a a b a $**

6
5
2
3
0
4
1

SA(T)

a
b
b
a
$
a
a

BWT(T)

Suffix Array is O(m)

BWT is O(m)

The BWT has a better constant factor!

Burrows-Wheeler Transform

BWM is related to the suffix array

$ a b a a b a
a $ a b a a b
a a b a $ a b
a b a $ a b a
a b a a b a $
b a $ a b a a
b a a b a $ a

BWM(T)

6	$
5	a $
2	a a b a $
3	a b a $
0	a b a a b a $
4	b a $
1	b a a b a $

SA(T)

Same order whether rows are rotations or suffixes

Burrows-Wheeler Transform

In fact, this gives us a new definition / way to construct BWT(T):

$$BWT[i] = \begin{cases} T[SA[i]-1] & \text{if } SA[i] > 0 \\ \$ & \text{if } SA[i] = 0 \end{cases}$$

"BWT = characters just to the left of the suffixes in the suffix array"

a b a a b a $

T

6	**$**
5	**a $**
2	**a a b a $**
3	**a b a $**
0	**a b a a b a $**
4	**b a $**
1	**b a a b a $**

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“BWT = characters just to the left of the suffixes in the suffix array”

a b a a b a $

T

6	$	a	T[5]
5	a $	b	T[4]
2	a a b a $	b	T[1]
3	a b a $	a	T[2]
0	a b a a b a $	$	**SA[0]**
4	b a $	a	T[3]
1	b a a b a $	a	T[0]

SA(T)

BWM(T)

$$T \xrightarrow{O(|T|)} SA(T) \xrightarrow{O(|T|)} BWM(T)$$