



# CS 225

**Welcome to Lab BTrees!**

# lab\_btree **Belligerent BTrees**

1. B-Tree Definition
2. Insertion
3. Implementation

# B-Tree Definition

A B-Tree of order  $m$

1. Maintains ordering within nodes
2. Each node has one more child than keys
3. All leaves are on the same level

Order = $m$	Possible number of keys	Possible number of children
Root node	$[1, m-1]$	$[2, m]$
Non-root nodes	$[\text{ceil}(m/2)-1, m-1]$	$[\text{ceil}(m/2), m]$

Order = 5	Possible number of keys	Possible number of children
Root node		
Non-root nodes		

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A B-Tree of order  $m$

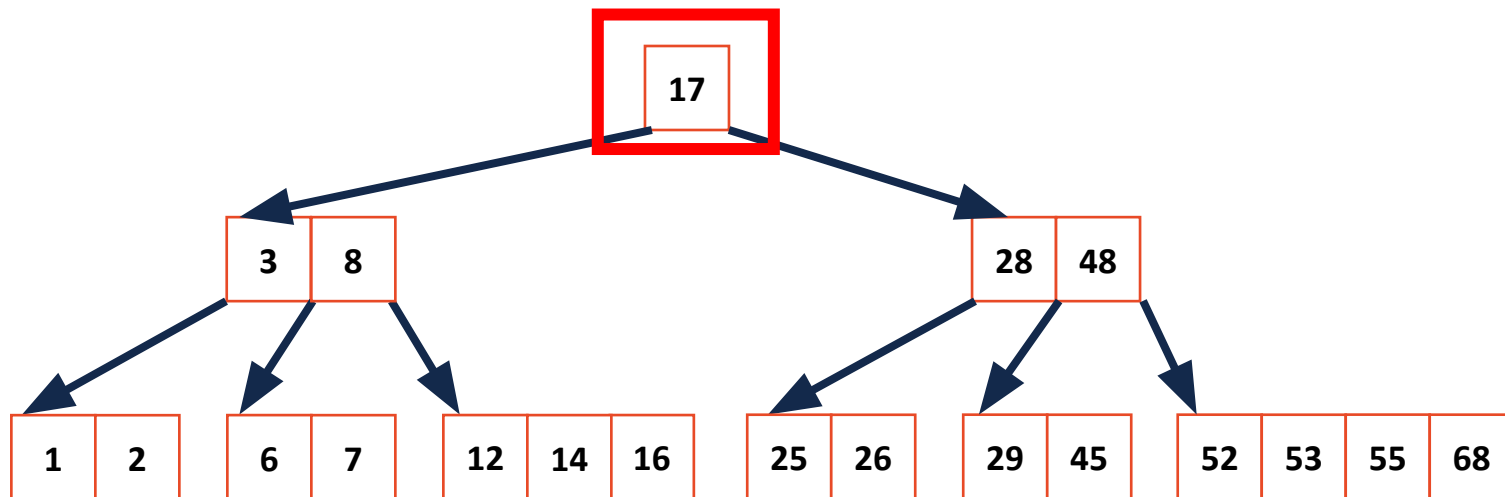
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Non-root nodes	$[2, 4]$	$[3, 5]$

# Btree Properties

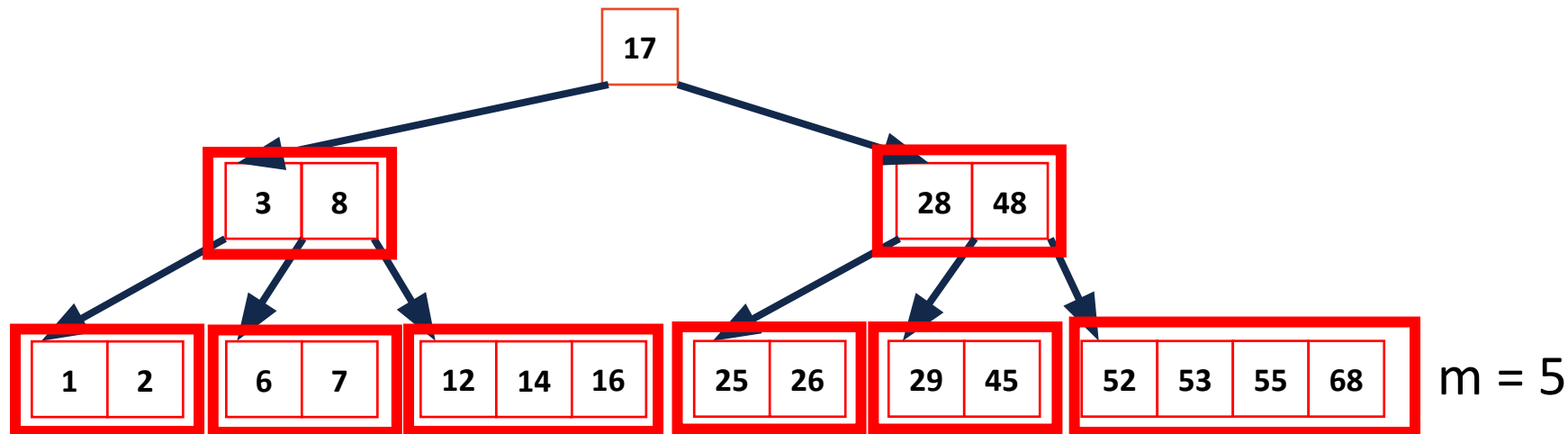
- Root nodes can have:  $[1, m-1]$  keys and  $[2, m]$  children
- All non-root, nodes have  $[\lceil m/2 \rceil - 1, m-1]$  keys.
- All non-root internal nodes have  $[\lceil m/2 \rceil, m]$  children



$m = 5$

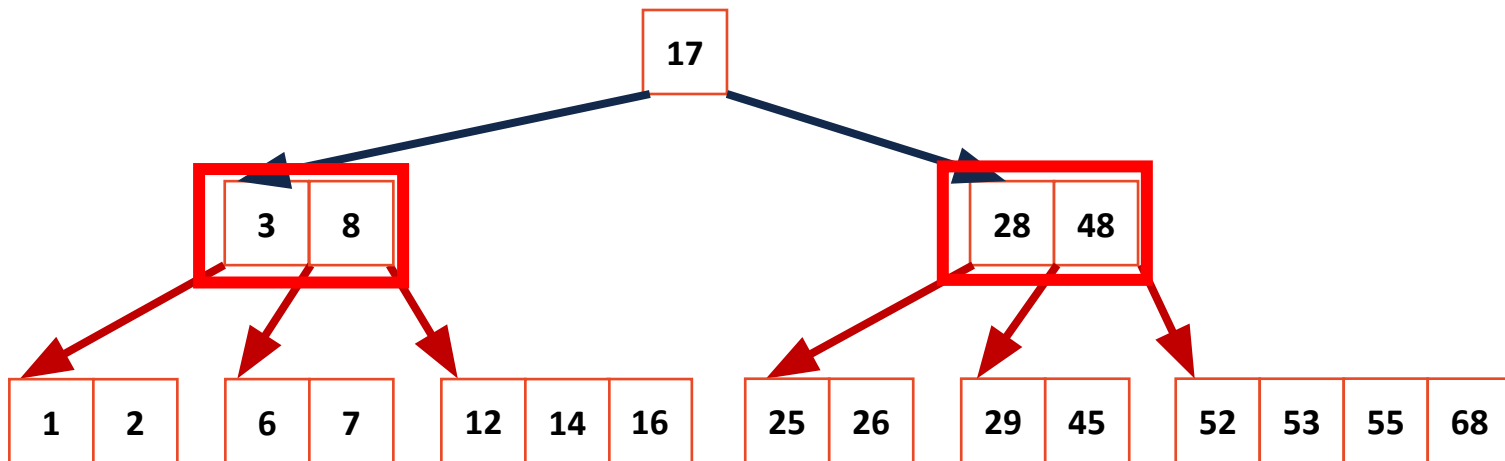
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# Btree Properties

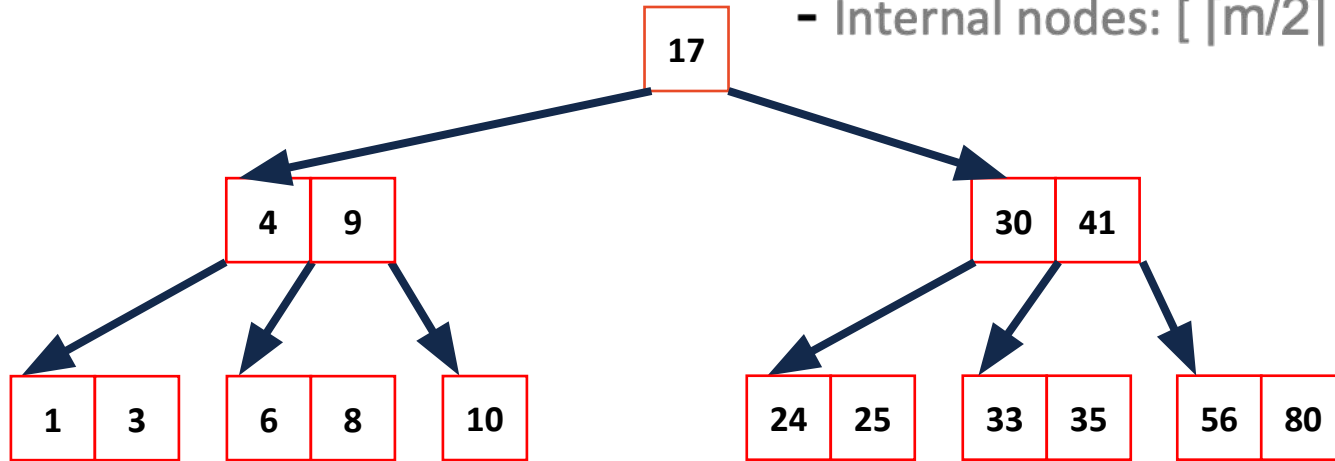
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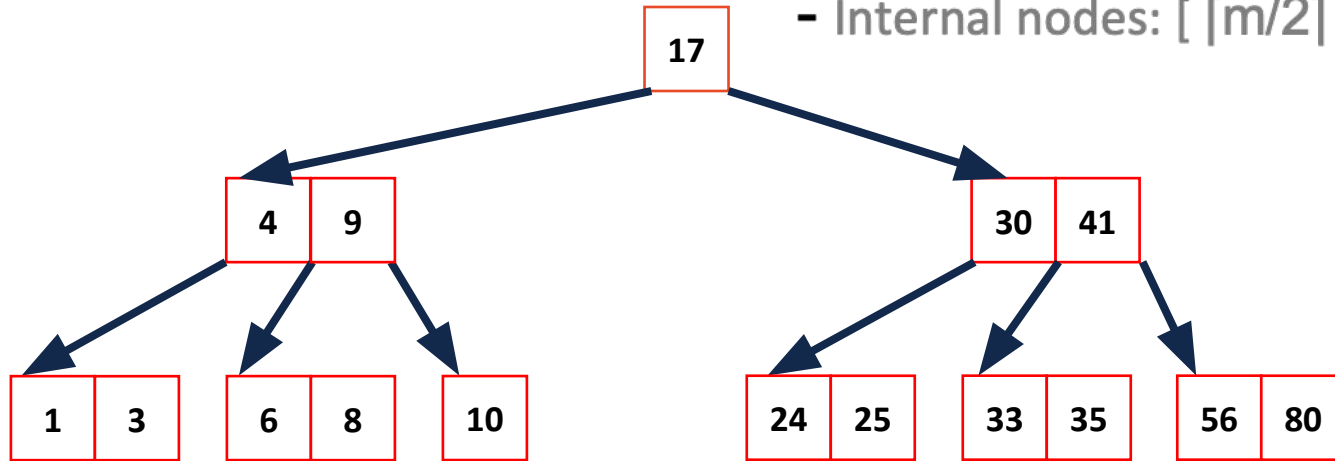
# What are the possible values of $m$ if the following B-Tree is of order $m$ ?

- $[ \lceil m/2 \rceil - 1, m-1 ]$  keys
- Internal nodes:  $[ \lceil m/2 \rceil, m ]$  children



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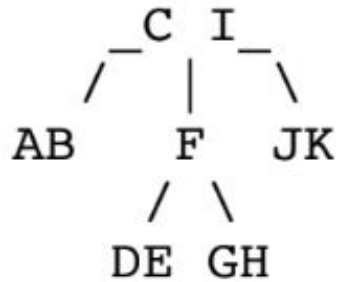


1.  $m \geq 3$  // we have nodes with three children
2. We have node with just one element, that means

$$\text{ceil} \left( \frac{m}{2} \right) - 1 = 1 \Rightarrow \text{ceil} \left( \frac{m}{2} \right) = 2 \Rightarrow m = 3, 4$$

*1 is the minimum number of elements node can have, hence (=)*

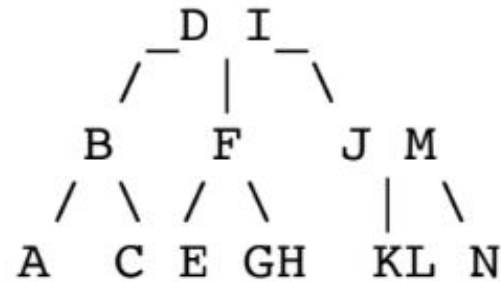
# Which B-tree is a valid B-tree?



1



2

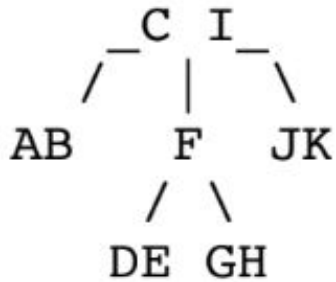


3



4

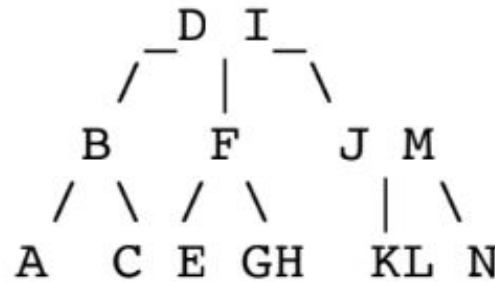
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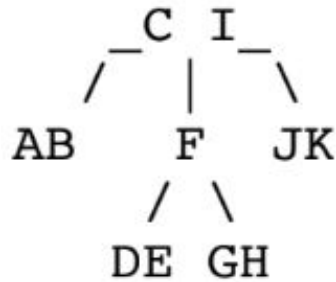
3



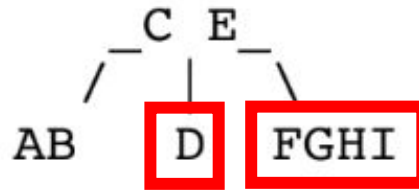
4

1. Leaves are on different level

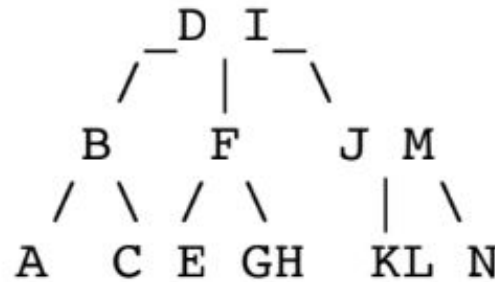
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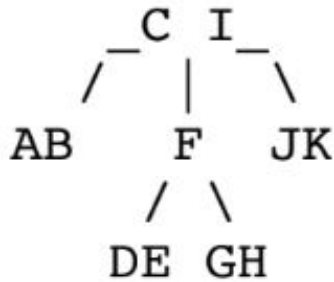
3



4

2. m should be at least 5, but we have node with only one element in it.

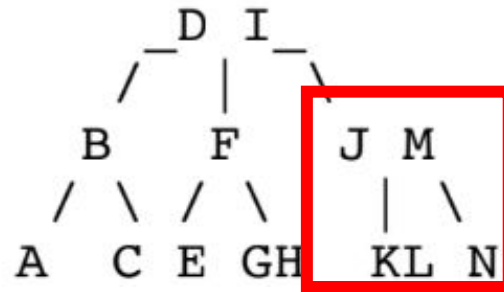
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1



2



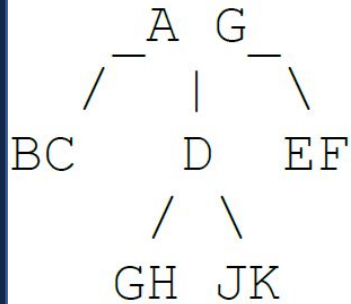
3



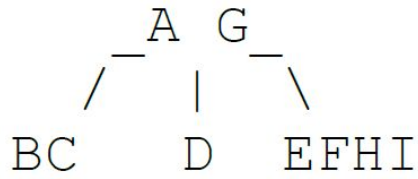
4

3. Node 'JM' does not have enough children

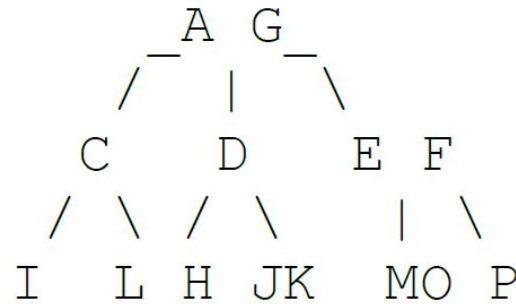
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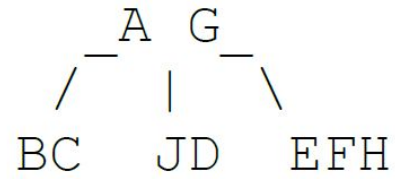
1



2



3



4

4. B-tree:  $m = 4, 5, 6$  ✓

- $[\text{ceil}(m/2)-1, m-1]$  keys
- Internal nodes:  $[\text{ceil}(m/2), m]$  children

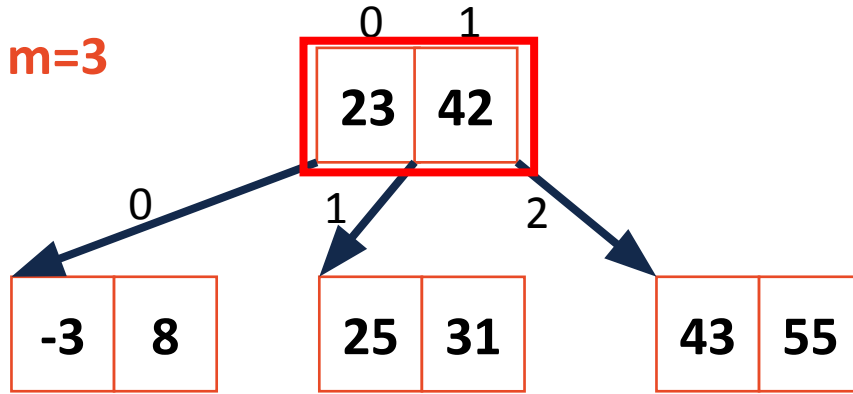
# Insertion

1. Insert the new element into a leaf node (keeping sorted property).
2. If the leaf node has more than  $m-1$  elements: we must ***split*** the node by ***throwing up*** the middle element to the parent node.  
Check parent node for overflow;



# Insert(28)

m=3

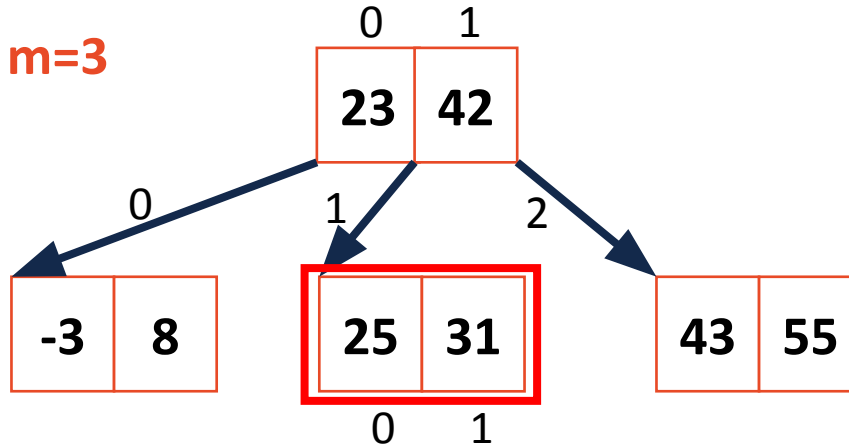


1. Find the corresponding leaf node, start from the root:

- $23 < \mathbf{28} < 42 \Rightarrow$  follow pointer 1;

# Insert(28)

**m=3**

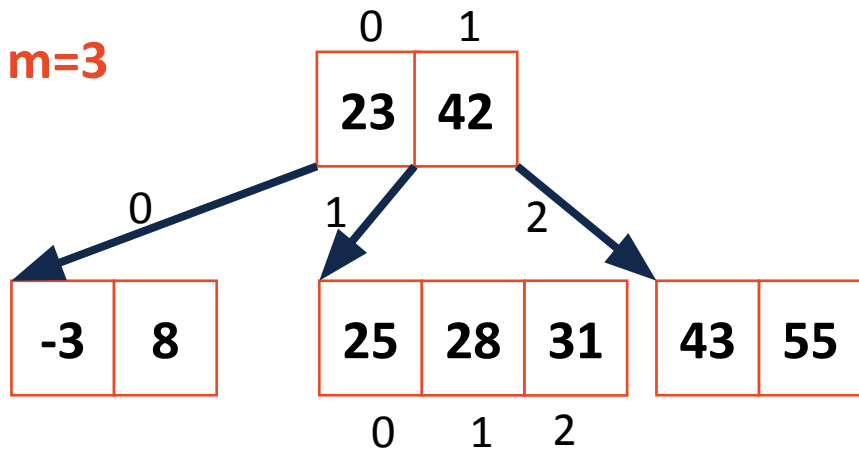


3. Current node is the leaf - How do you know?

4. Insert 28 on the right index:  $25 < \mathbf{28} < 31$

# Insert(28)

**m=3**



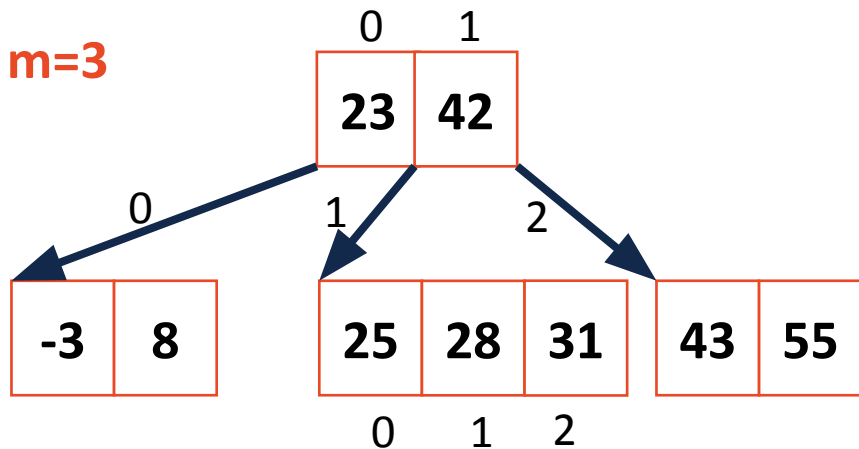
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# Insert(28)

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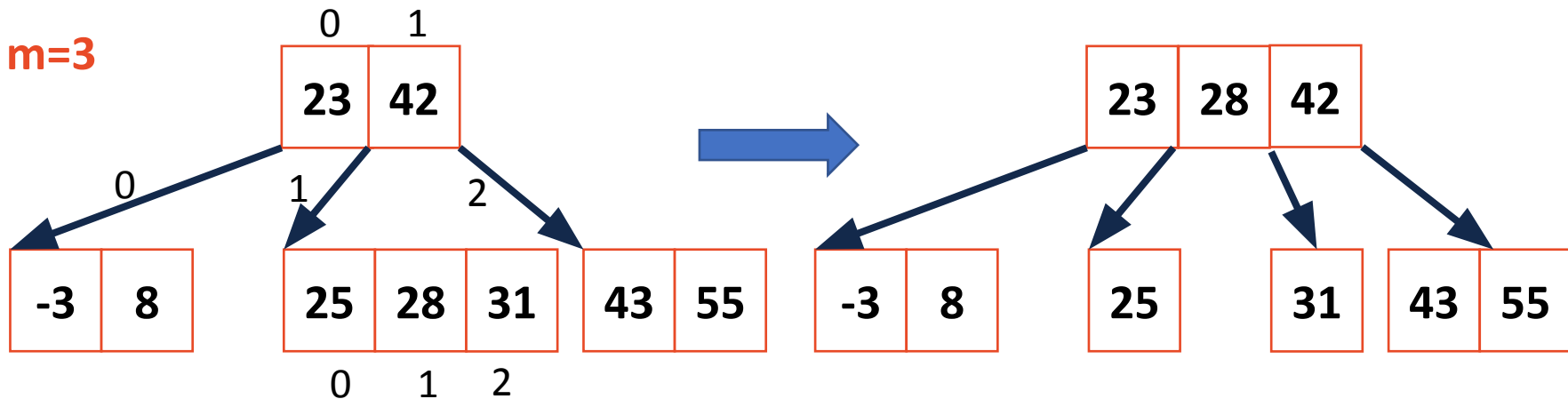
3. Current node is the leaf

4. Insert 28 on the right index:  $25 < \mathbf{28} < 31$

5. Check node for overflow: since  $m=3$ , we have to 'throw up' middle element and split the node.

# Insert(28)

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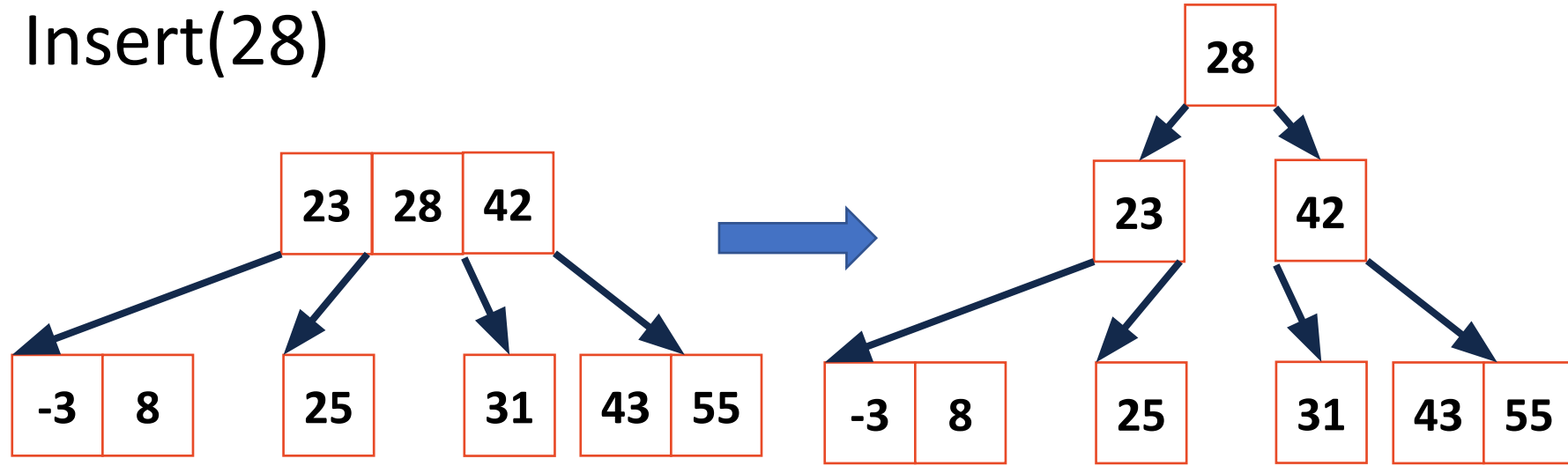


Split node:

Split the node and insert middle element into the parent node;

Check the parent node for overflow => split parent;

# Insert(28)



Split node (parent):

Split the node and 'throw up' middle element (create new node with 28);

Check new node for overflow => we are done!

```
class BTree
```

```
{
```

```
private:
```

```
    struct DataPair {
```

```
        K key;
```

```
        V value;
```

```
    ...
```

```
};
```

```
    struct BTreeNode {
```

```
        bool is_leaf;
```

```
        vector<DataPair> elements;
```

```
        vector<BTreeNode*> children;
```

```
        BTreeNode(bool is_leaf, unsigned int order);
```

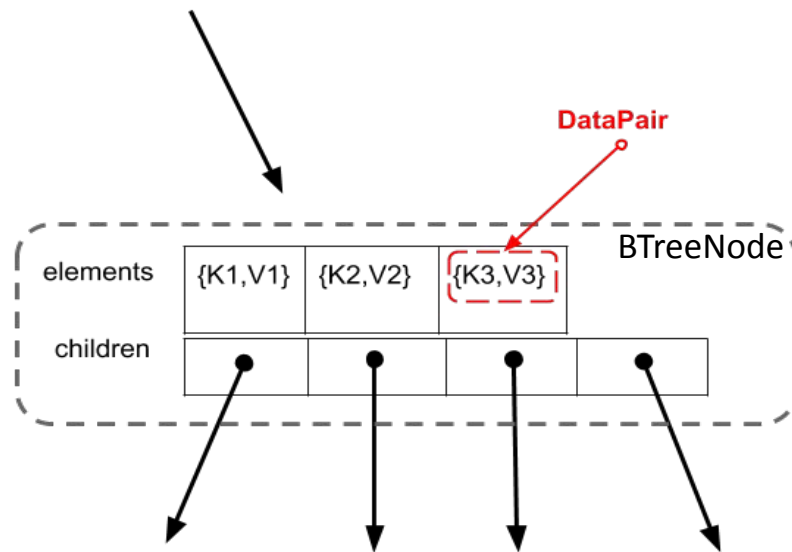
```
        ...
```

```
};
```

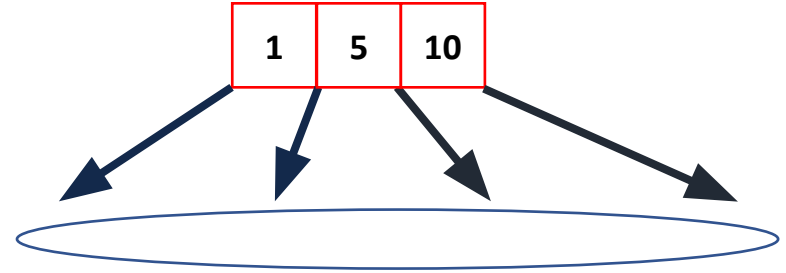
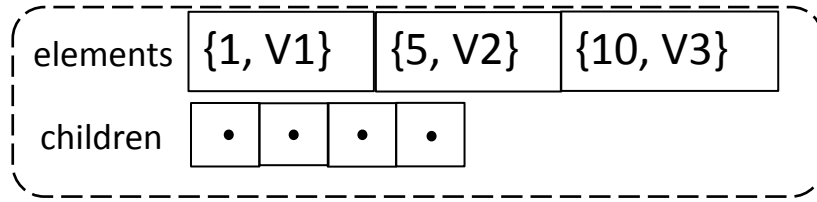
```
    unsigned int order;
```

```
    BTreeNode* root; ... }
```

# Implementation

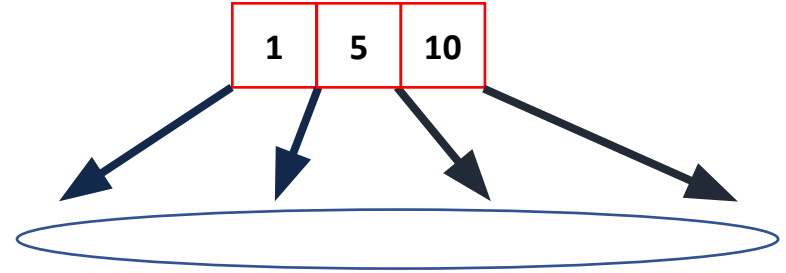
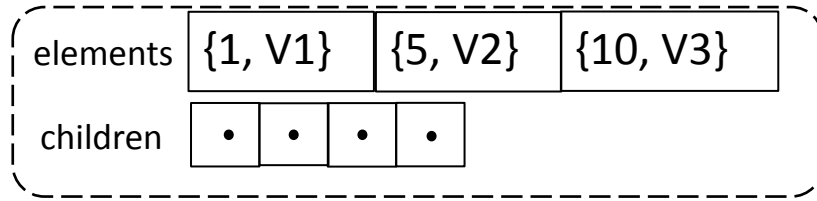


Suppose we are searching for the element with key = 7, which **child pointer** in the BTreeNode should we follow? Give its index in the **children** vector.





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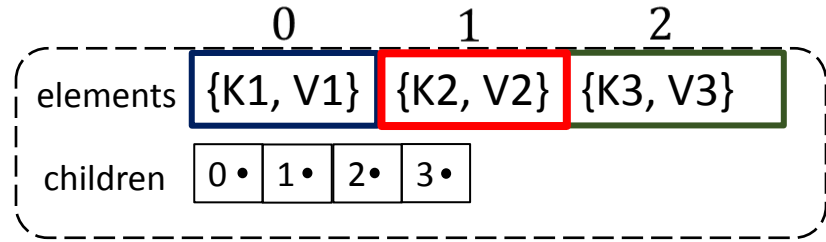


Since  $5 < 7 < 10$

We need a child node between 5 and 10:

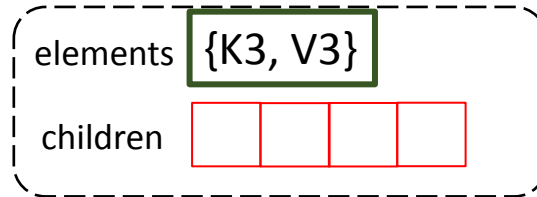
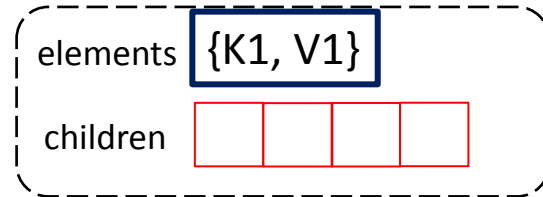
(1 is index of 5 and since  $5 < 7$  we need child on the right side of the 5, thus index 2)

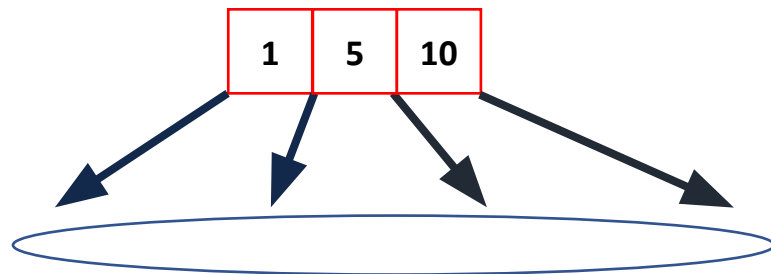
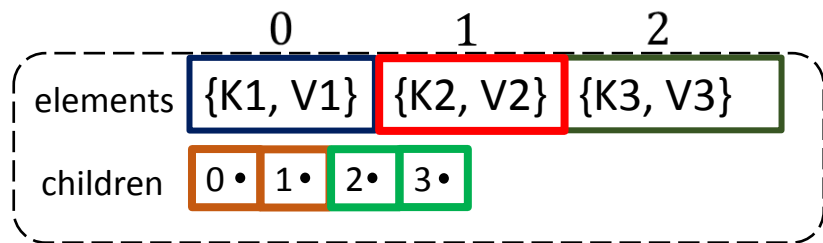
Suppose we need to split the given BTreeNode (so we will throw up the element  $\{K2, V2\}$ ) how would we split the children vector between the newly created BTreeNodes after the split?



Index of  $\{k2, V2\}$  element is  $i = 1$

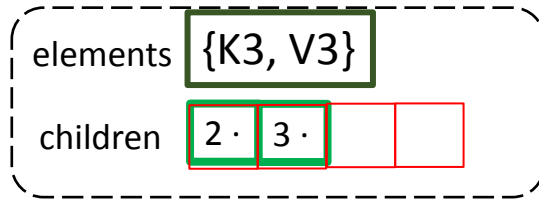
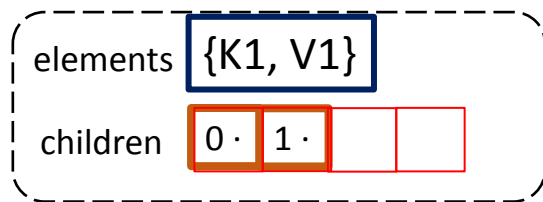
**Split elements:**  $[0, i)$ ,  $[i + 1, \text{end}) \rightarrow [0, 1)$ ,  $[2, 3)$



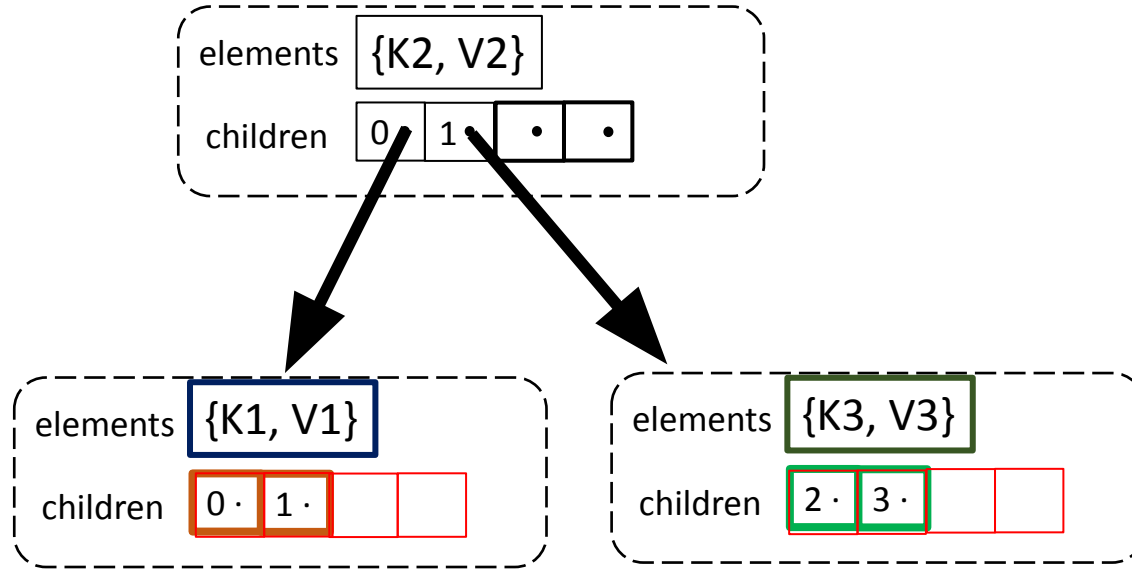


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**Split children:**  $[0, i], [i + 1, \text{end}] \rightarrow [0, 1], [2, 3]$



Create parent node with {K2, V2}, insert pointers to children nodes



# Summary

Insertion:

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    ...
    struct BTreeNode {
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        BTreeNode(bool is_leaf, unsigned int order);
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    };
    unsigned int order;
    BTreeNode* root; ... }
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Implementation

