



CS 225

Data Structures



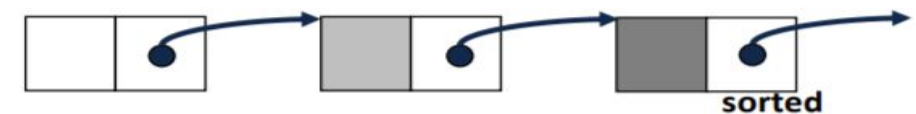
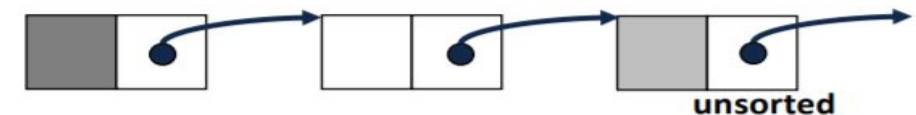
Previously in lectures:

- Priority Queue ADT
- (min)Heap
- insert
- heapifyUp
- removeMin
- heapifyDown
- BuildHeap

Priority Queue Implementation

- **Scenario:** n elements, each with an associated priority p_i , is given to you one element at a time.
- **Task:** How would you remove the highest priority element?

Insert	removeMin
$O(1)^*$	$O(n)$
$O(1)$	$O(n)$
$O(n)$	$O(1)$
$O(n)$	$O(1)$



We Can Do Better! Enter Heaps!

Priority Queue Implementation

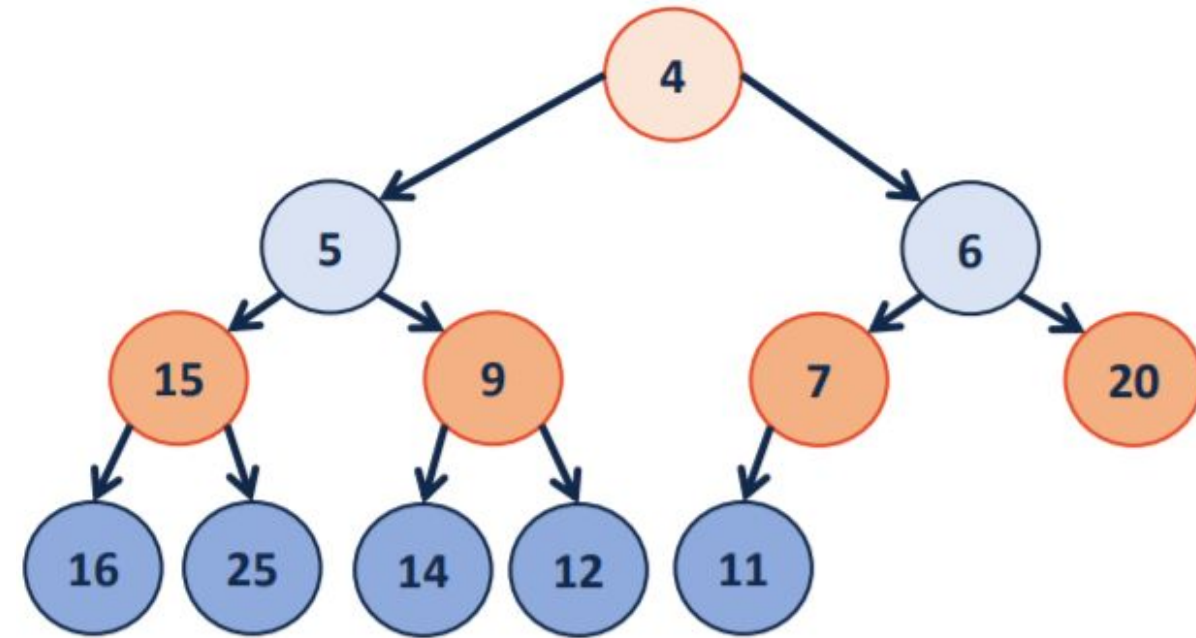
- **Scenario:** n elements, each with an associated priority p_i , is given to you one element at a time.
- **Task:** How would you remove the highest priority element?

Insert	removeMin
$O(1)^*$	$O(n)$
$O(1)$	$O(n)$
$O(n)$	$O(1)$
$O(n)$	$O(1)$



Heaps

- Complexity
 - *Insert*: $O(\log n)$
 - *removeMin*: $O(\log n)$
 - *building heap*: $O(n)$ [One-time]
- each node is an element with an associated priority
- root *always* has the highest priority (expressed as number)
- Binary Heaps: Complete binary tree
 - *min-heap*: minimum between two numbers; root is smallest element.
 - *max-heap*: maximum between two numbers; root is largest element.
- Not BST: in-order traversal may not yield an ordered list.

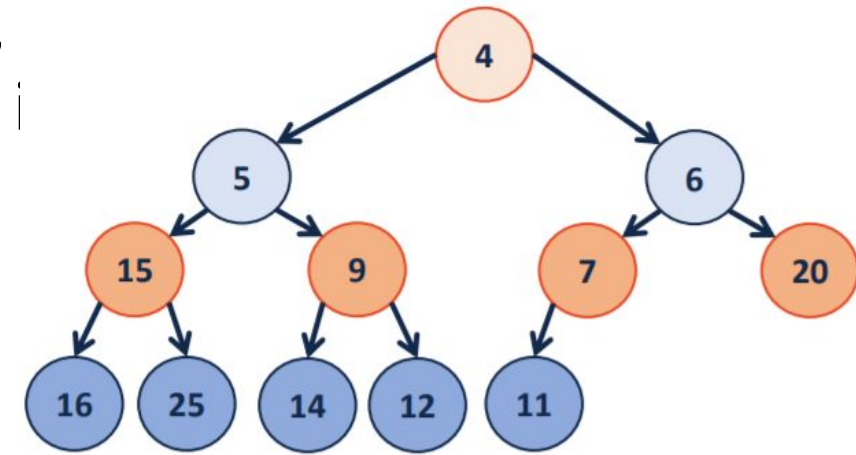


Remember: The parent has higher priority than any of its children



Worksheet Exercise #1

Exercise 1: Suppose the current node is at index i , fill in the blank for the indices of certain locations in the tree. You will be implementing these as functions in your lab. Remember, these formulas depend on whether you are populating the 0th index of the array or not.



Current node = i

Root index = 0

Right child =

Left child =

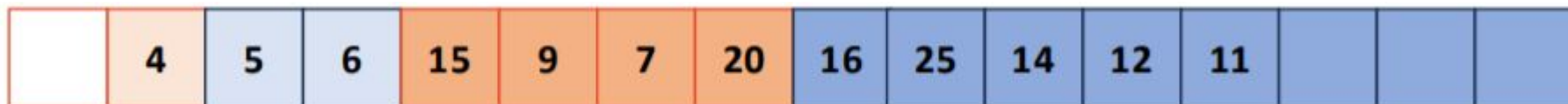
Parent =

Root index = 1

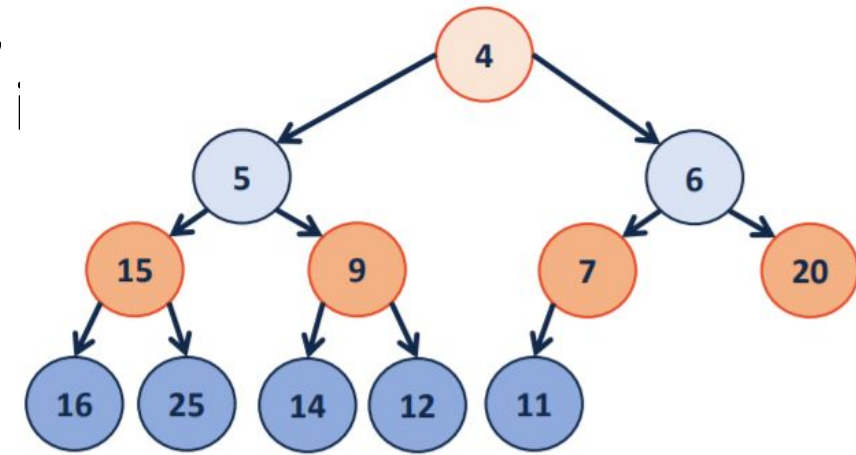
Right child =

Left child =

Parent =



Exercise 1: Suppose the current node is at index i , fill in the blank for the indices of certain locations in the tree. You will be implementing these as functions in your lab. Remember, these formulas depend on whether you are populating the 0th index of the array or not.



Current node = i

Root index = 0

Right child = $2*i + 2$

Root index = 1

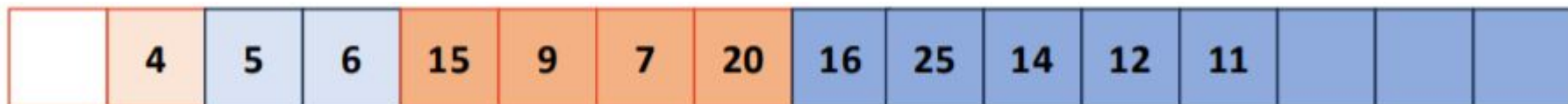
Right child =

Left child = $2*i + 1$

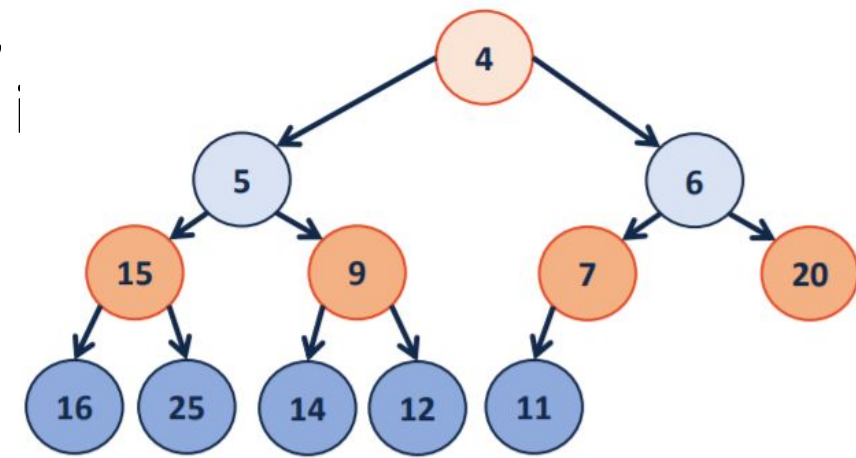
Parent = $\text{floor}[(i-1)/2]$

Left child =

Parent =



Exercise 1: Suppose the current node is at index i , fill in the blank for the indices of certain locations in the tree. You will be implementing these as functions in your lab. Remember, these formulas depend on whether you are populating the 0th index of the array or not.



Current node = i

Root index = 0

Right child = $2*i + 2$

Root index = 1

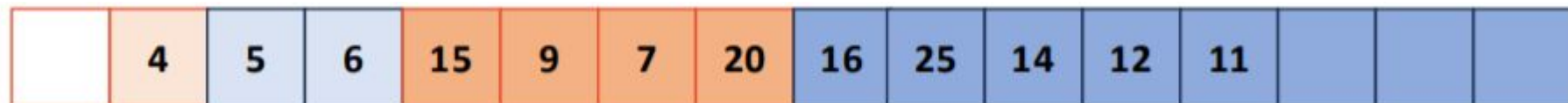
Right child = $2*i + 1$

Left child = $2*i + 1$

Parent = $\text{floor}[(i-1)/2]$

Left child = $2*i$

Parent = $\text{floor}[i/2]$



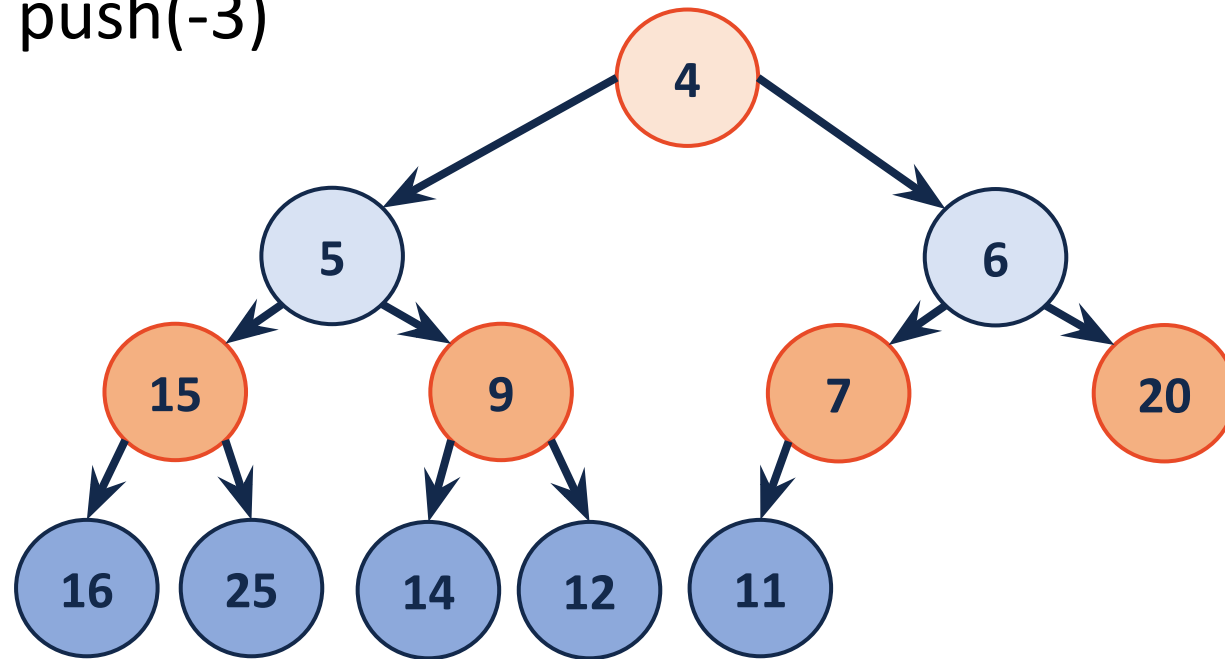


Insert/Push

push(key)

1. Add new element in the end of the vector (index i)
2. Reset heap property from index i – `heapifyUp(i)`

push(-3)



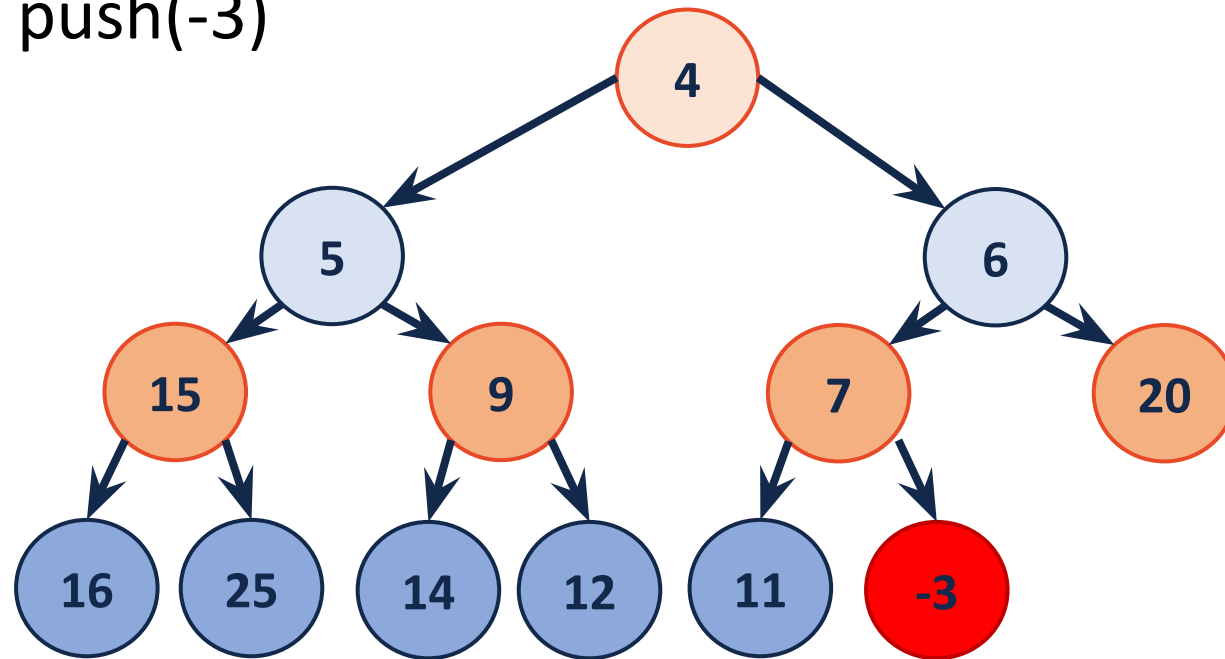
	4	5	6	15	9	7	20	16	25	14	12	11
--	---	---	---	----	---	---	----	----	----	----	----	----

push(key)

1. Add new element in the end of the vector (index i)
2. Reset heap property from index i using

`heapifyUp(i)`

push(-3)



	4	5	6	15	9	7	20	16	25	14	12	11	-3
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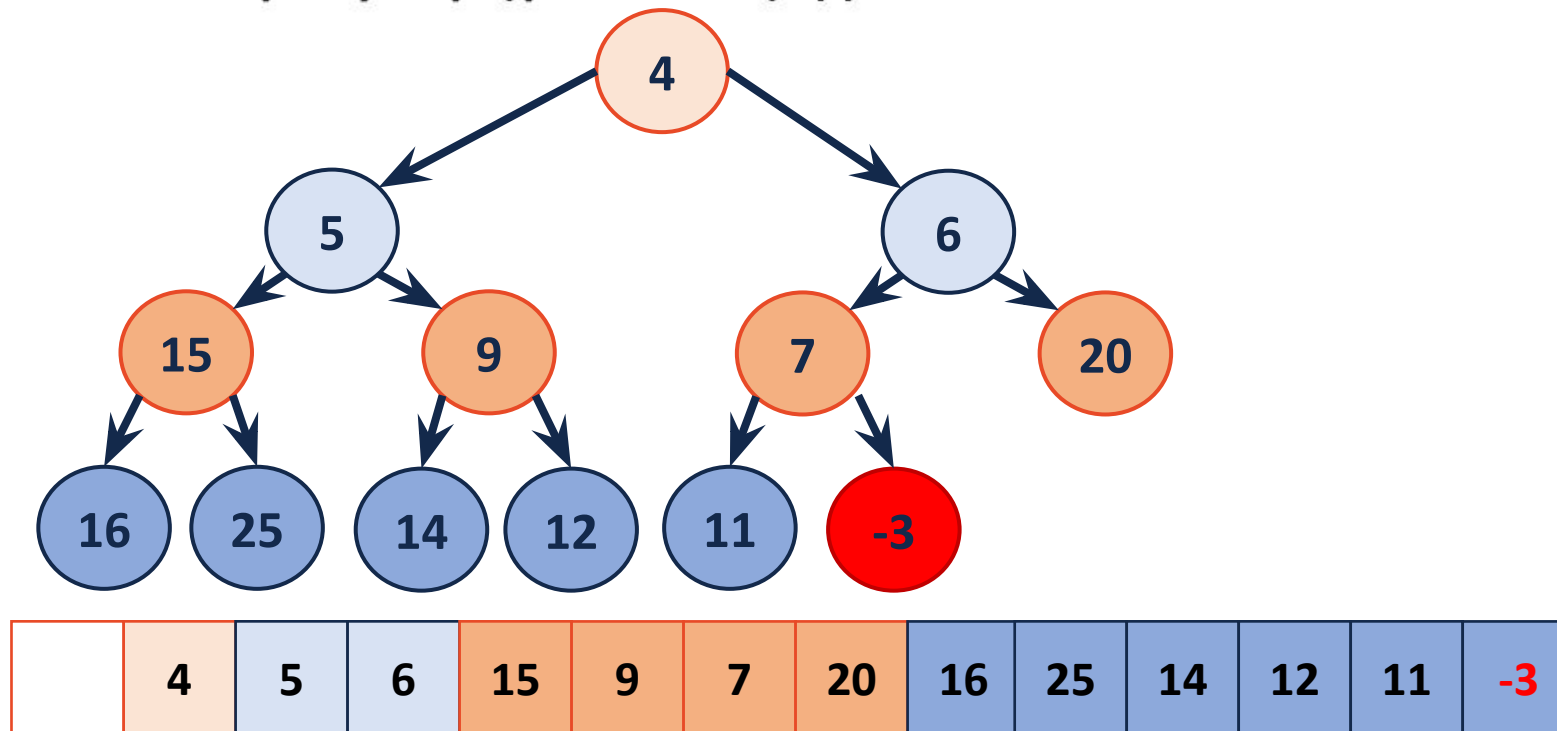
While the new element on given index has a smaller value than its parent, swap the element and its parent.

heapifyUp(*i*)

if $i \neq \text{rootIndex} \ \&\& \ A[i] < A[\text{parent}(i)]$

 swap(*i*, parent(*i*))

 heapifyUp(parent(*i*))



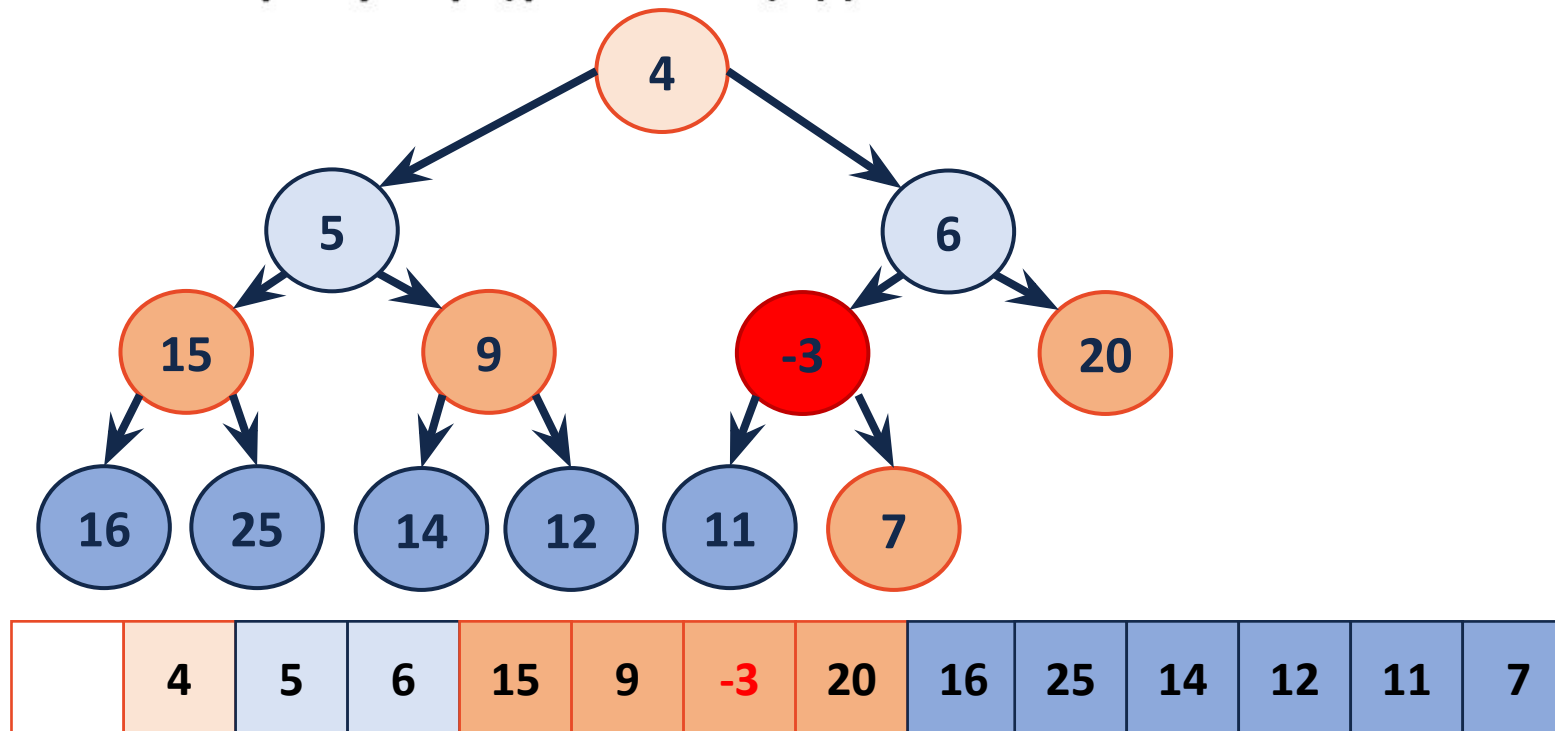
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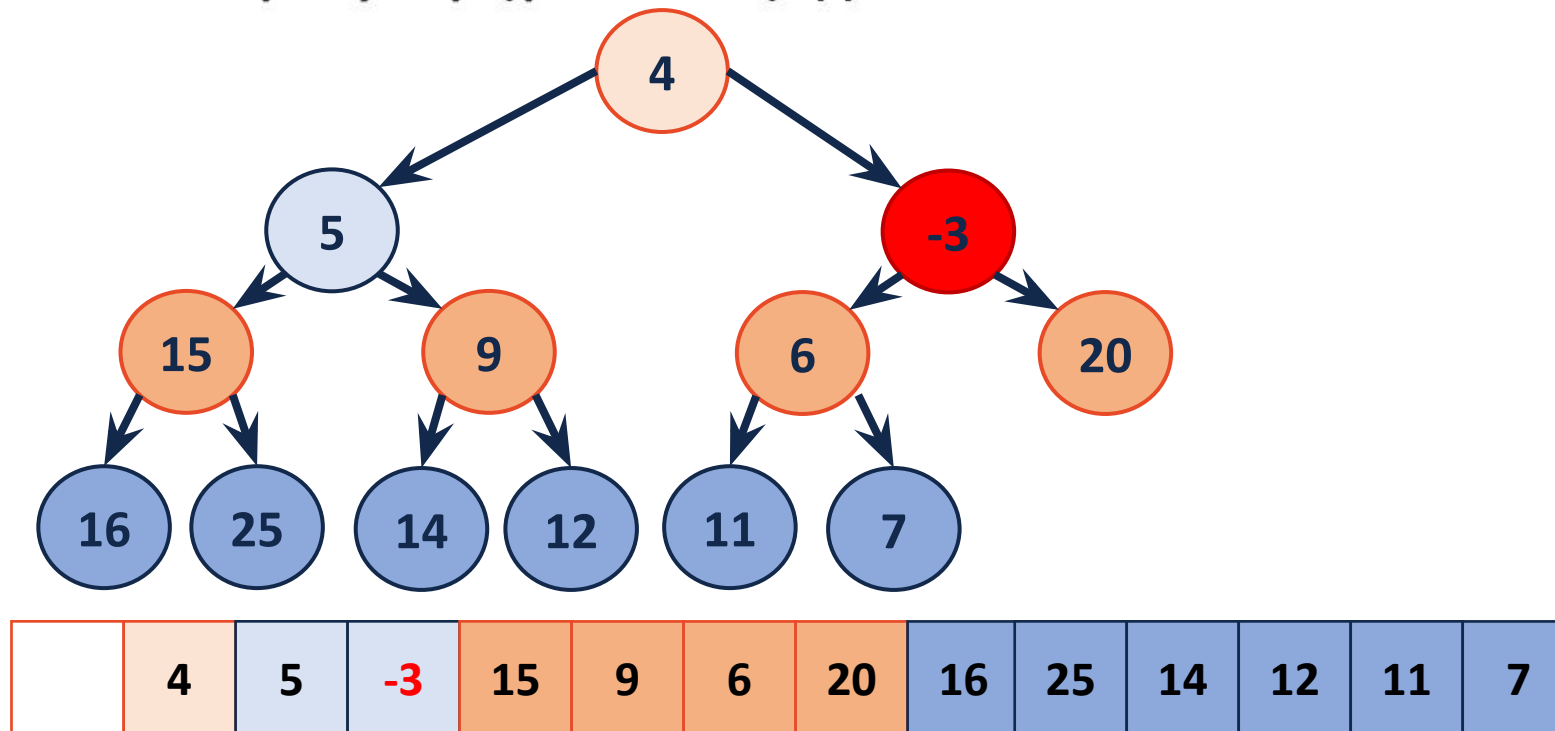
While the new element on given index has a smaller value than its parent, swap the element and its parent.

heapifyUp(i)

if $i \neq \text{rootIndex} \ \&\& \ A[i] < A[\text{parent}(i)]$

 swap(i , parent(i))

 heapifyUp(parent(i))



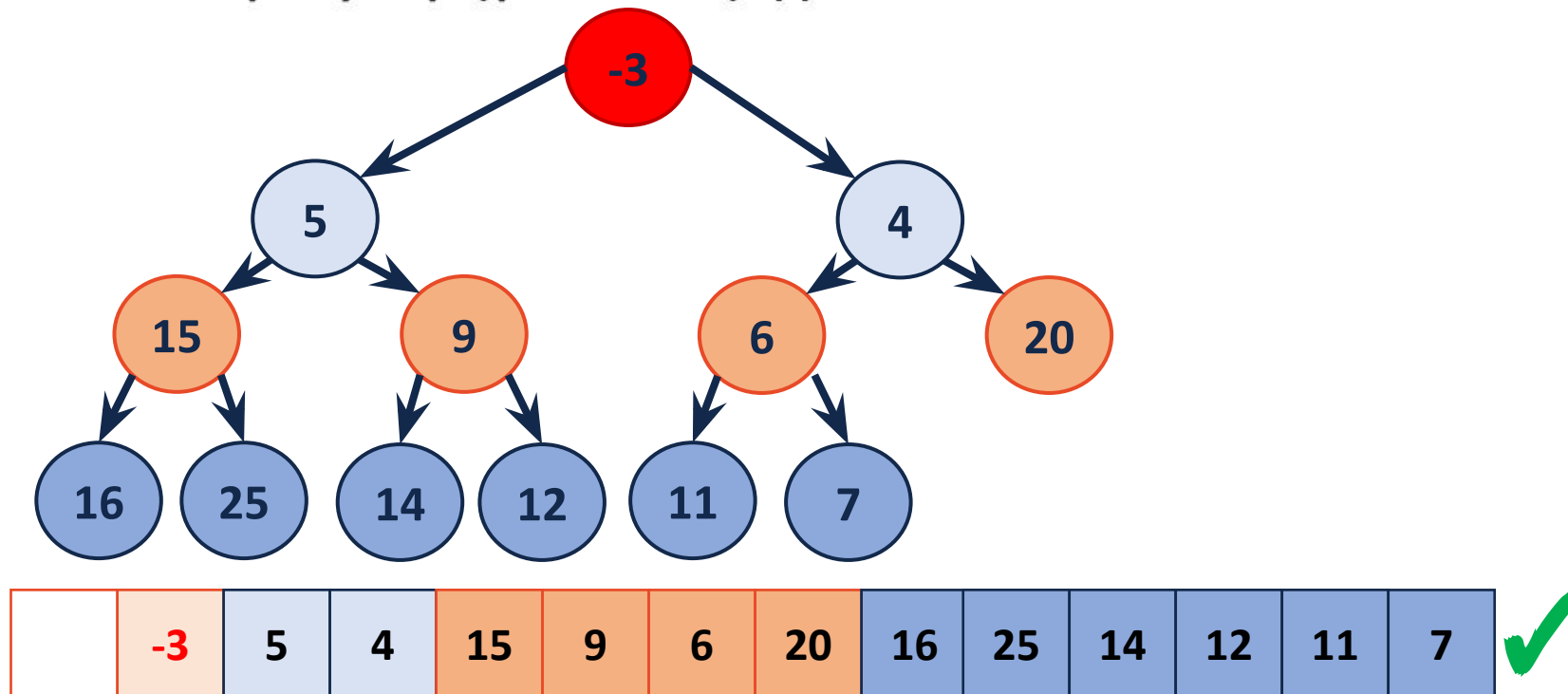
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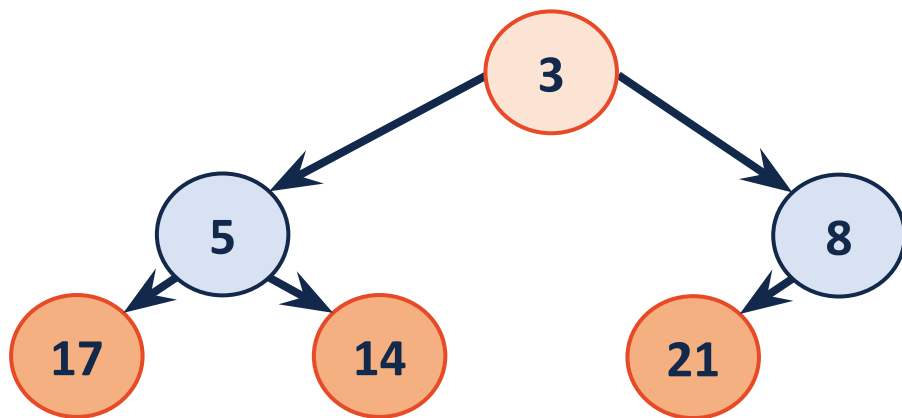


Worksheet

Exercise #2.1 and #2.2

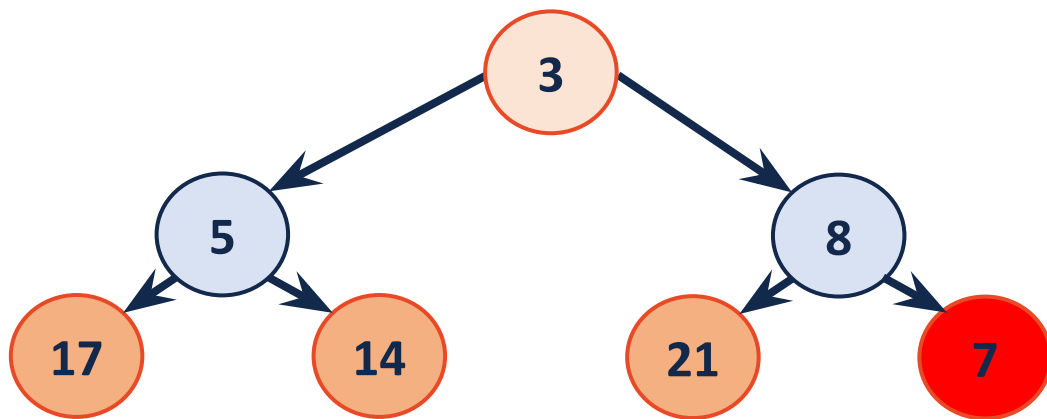
Suppose 7 is inserted into the heap below.

- **Exercise 2.1:** Suppose 7 is inserted into the heap below. What will 7's left and right children be?
- **Exercise 2.2:** What is the array representation of the tree after 7 is inserted?



Suppose 7 is inserted into the heap below.

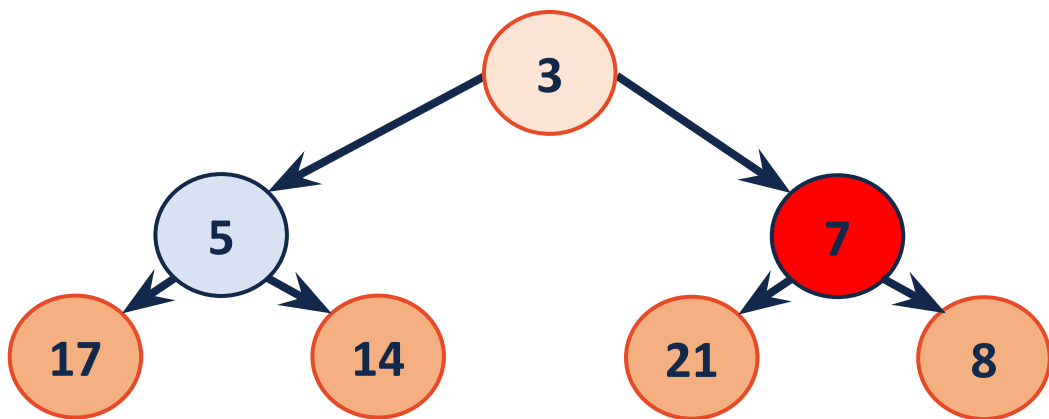
- **Exercise 2.1:** Suppose 7 is inserted into the heap below. What will 7's left and right children be?
- **Exercise 2.2:** What is the array representation of the tree after 7 is inserted?



	3	5	8	17	14	21	7
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Suppose 7 is inserted into the heap below.

- **Exercise 2.1:** Suppose 7 is inserted into the heap below. What will 7's left and right children be? Left child = 21; Right Child = 8;
- **Exercise 2.2:** What is the array representation of the tree after 7 is inserted?

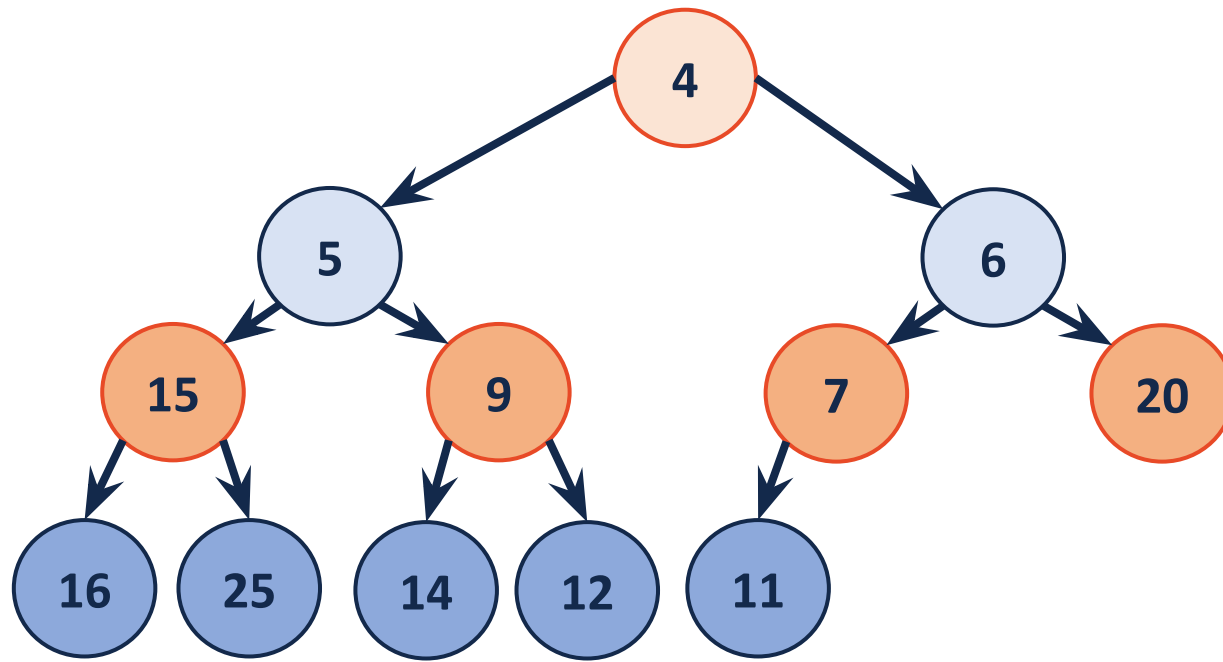




Delete/Pop

pop()

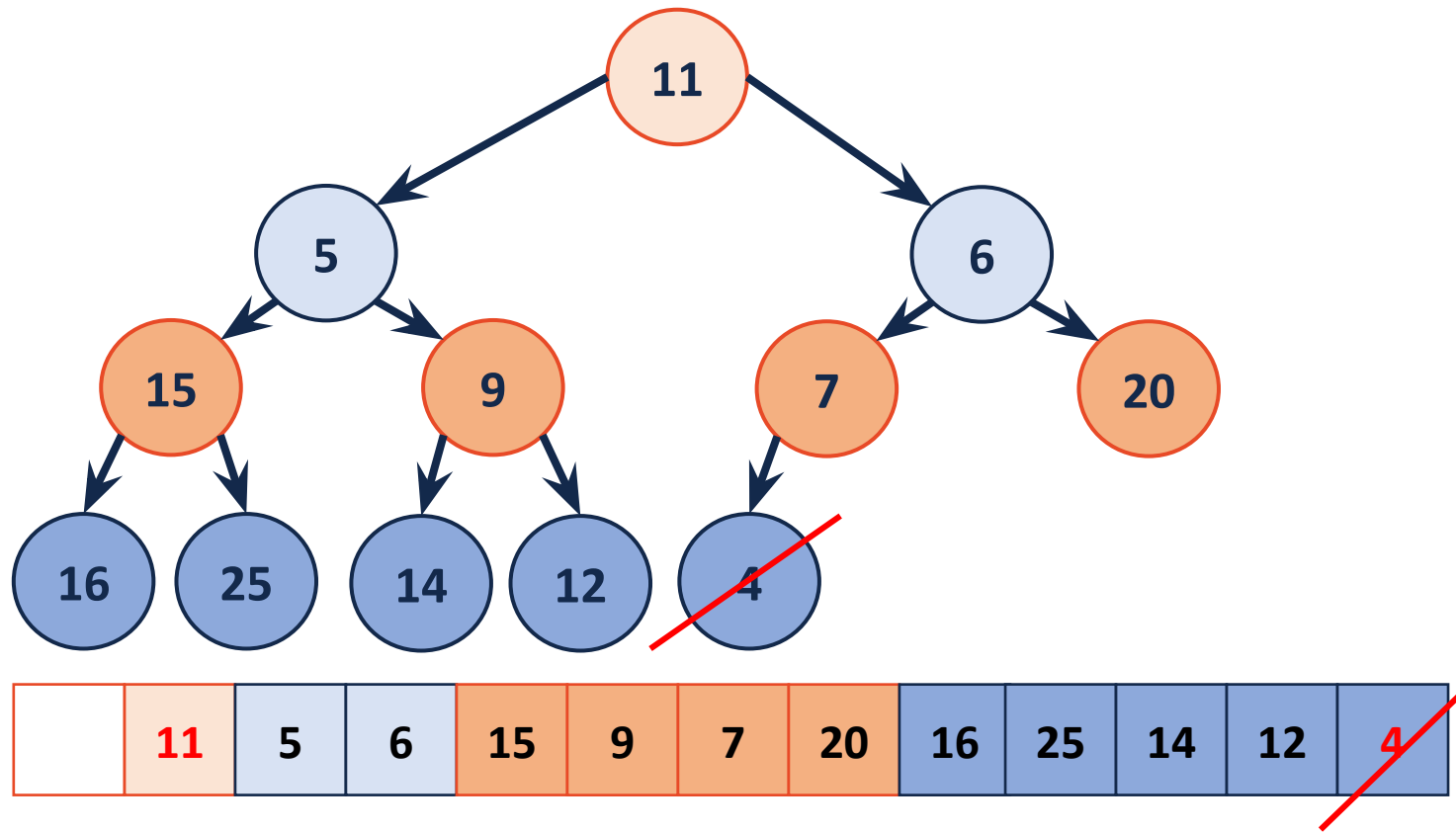
1. Swap the last element(index i) for the root
2. `result=_elems[i]`
3. Remove the last element: `(_elems.pop_back())`
4. `HeapifyDown(root);`
5. Return `result`;



	4	5	6	15	9	7	20	16	25	14	12	11
--	---	---	---	----	---	---	----	----	----	----	----	----

pop()

1. Swap the last element(index i) for the root
2. `result=_elems[i]`
3. Remove the last element: `(_elems.pop_back())`
- ➔ 4. `HeapifyDown(root);`
5. Return `result`;



HeapifyDown(current)

If ! isLeaf(current)

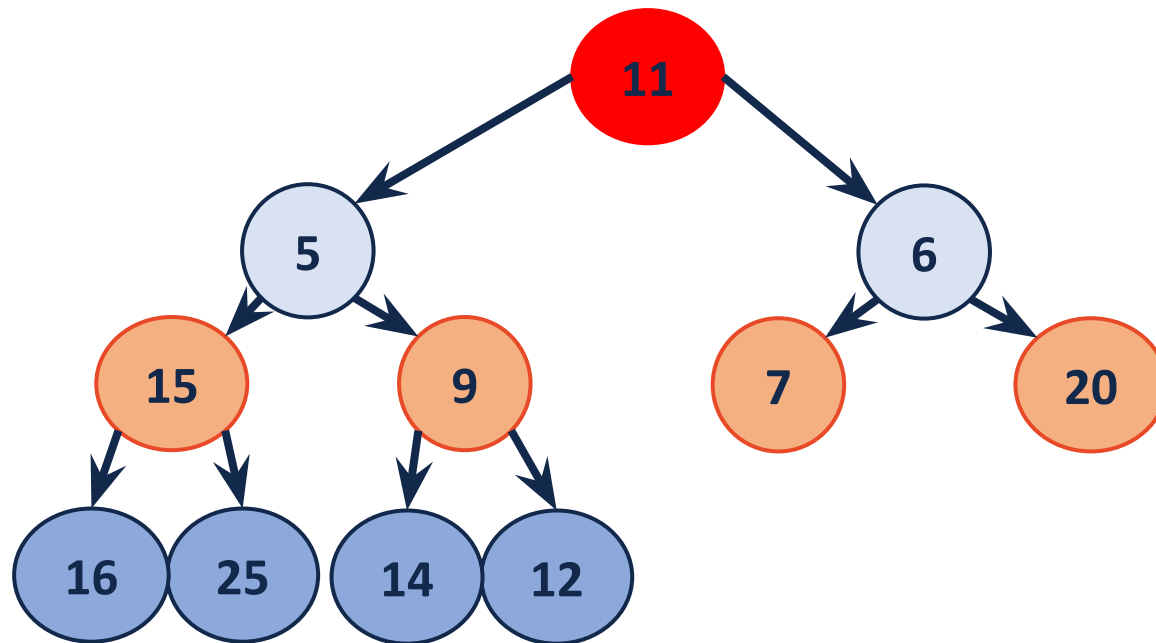
Find the min = index of min child of current

If $A[\text{current}] > A[\text{min}]$ child

swap $A[\text{current}]$ and $A[\text{min}]$

HeapifyDown(min)

//heapifyDown restores heap
property only when left subtree and
right subtree are already heaps!



	11	5	6	15	9	7	20	16	25	14	12
--	----	---	---	----	---	---	----	----	----	----	----



HeapifyDown(current)

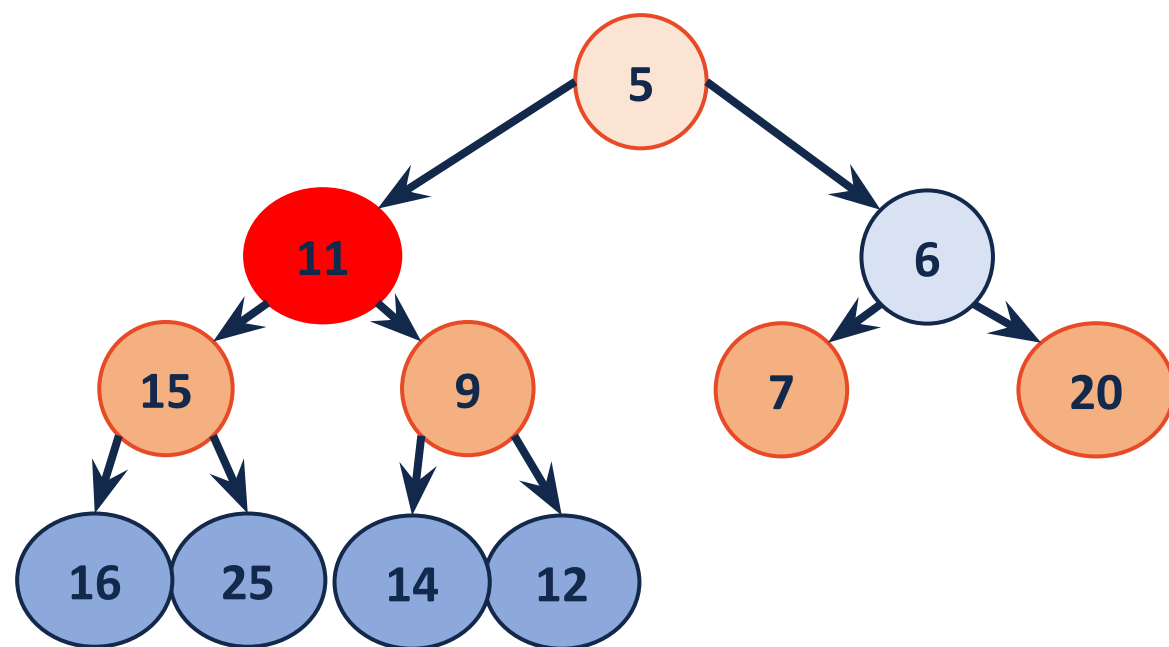
If ! isLeaf(current)

Find the min = index of min child of current

If $A[\text{current}] > A[\text{min}]$ child

swap $A[\text{current}]$ and $A[\text{min}]$

HeapifyDown(min)



	5	11	6	15	9	7	20	16	25	14	12
--	---	----	---	----	---	---	----	----	----	----	----

HeapifyDown(current)

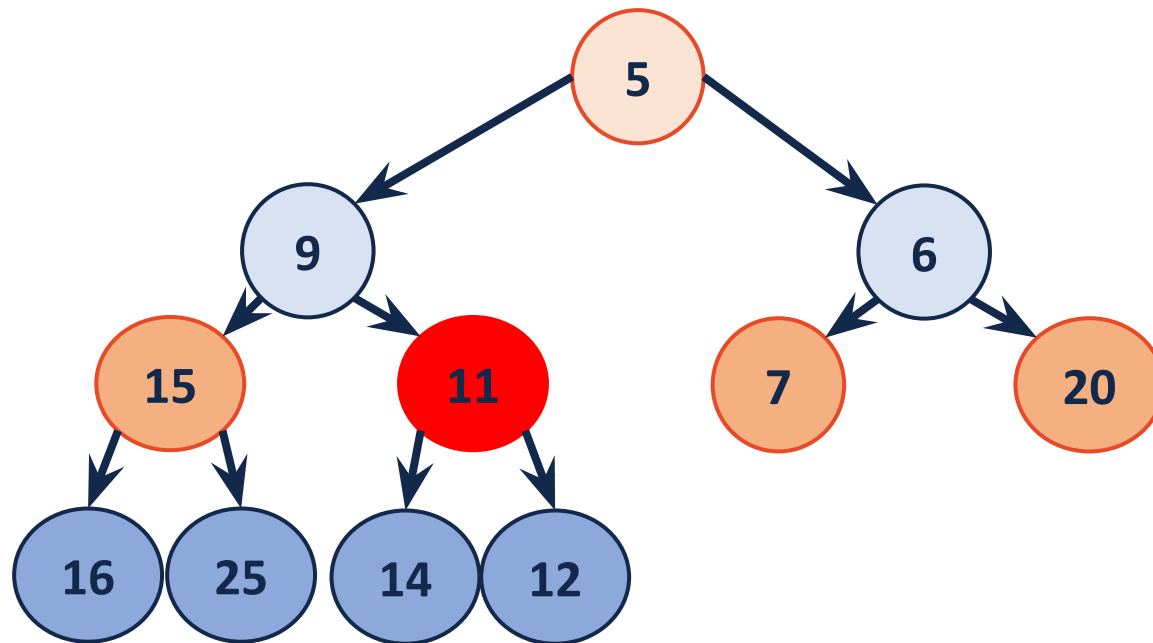
If ! isLeaf(current)

Find the min = index of min child of current

If $A[\text{current}] > A[\text{min}]$ child

swap $A[\text{current}]$ and $A[\text{min}]$

HeapifyDown(min)



	5	9	6	15	11	7	20	16	25	14	12
--	---	---	---	----	----	---	----	----	----	----	----

HeapifyDown(current)

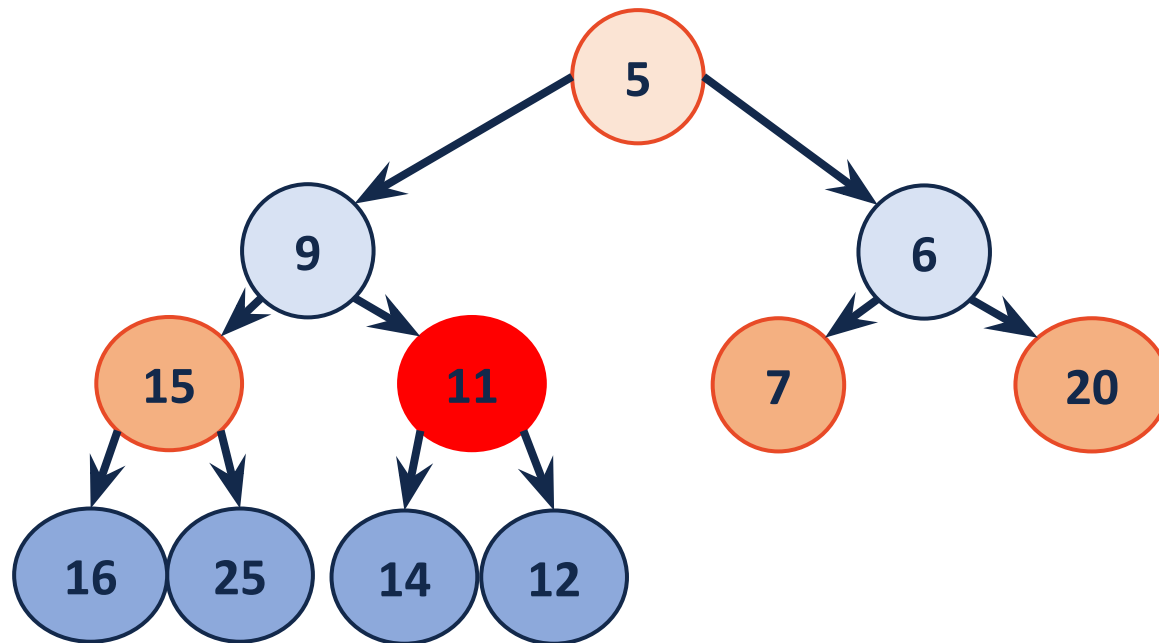
If ! isLeaf(current)

Find the min = index of min child of current

If $A[\text{current}] > A[\text{min}]$ child

swap $A[\text{current}]$ and $A[\text{min}]$

HeapifyDown(min)



	5	9	6	15	11	7	20	16	25	14	12
--	---	---	---	----	----	---	----	----	----	----	----

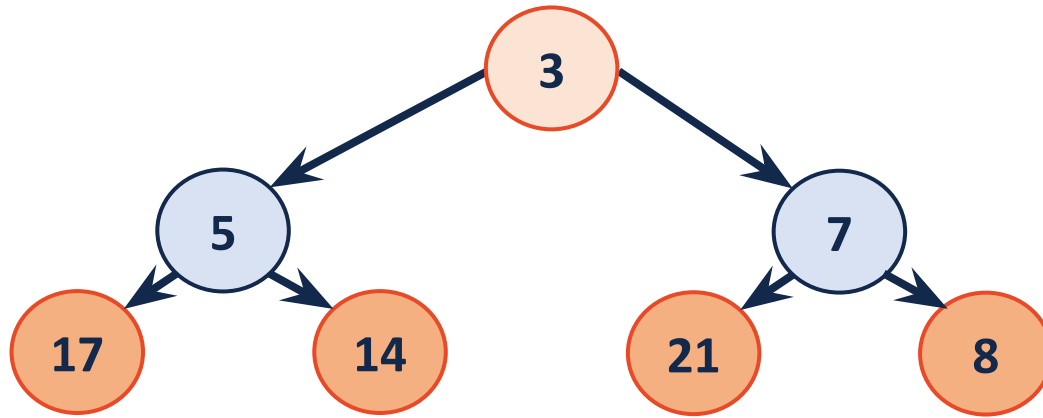




Worksheet

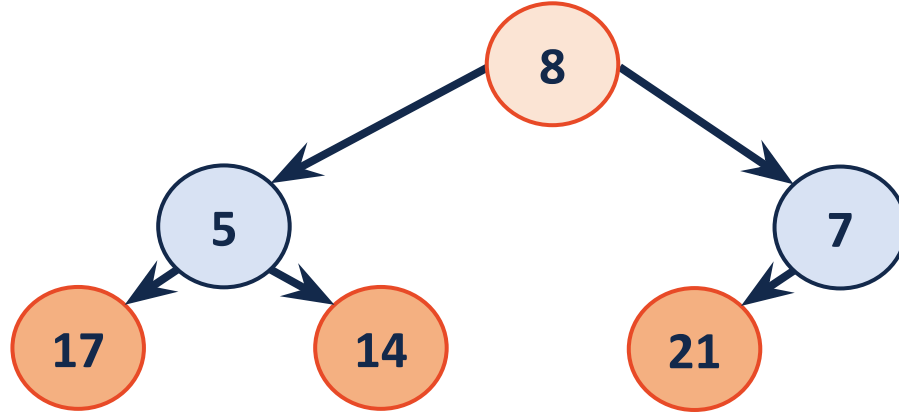
Exercise #2.3

Exercise 2.3: What is the array representation of the tree if **3** is removed?



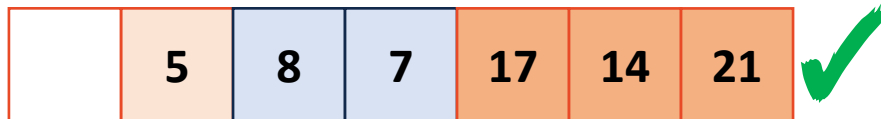
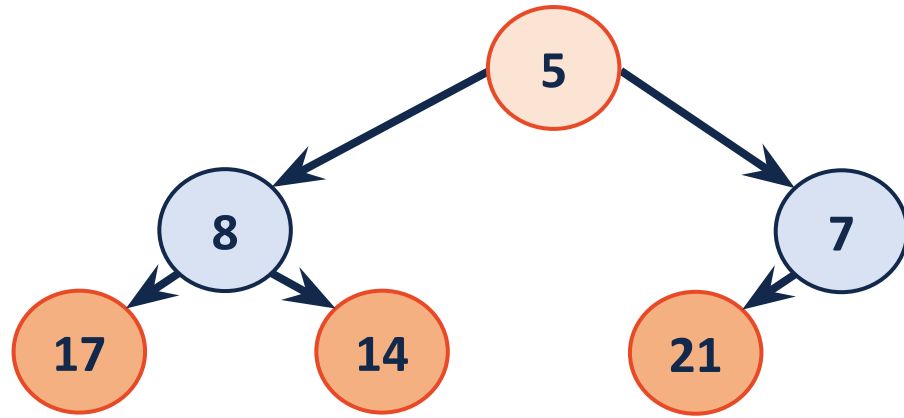
	3	5	7	17	14	21	8
--	---	---	---	----	----	----	---

Exercise 2.3: What is the array representation of the tree if **3** is removed?



	8	5	7	17	14	21
--	---	---	---	----	----	----

Exercise 2.3: What is the array representation of the tree if **3** is removed?

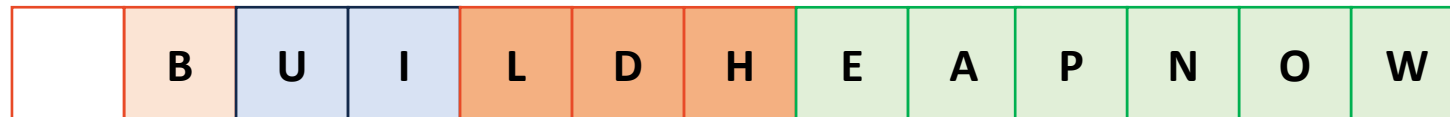
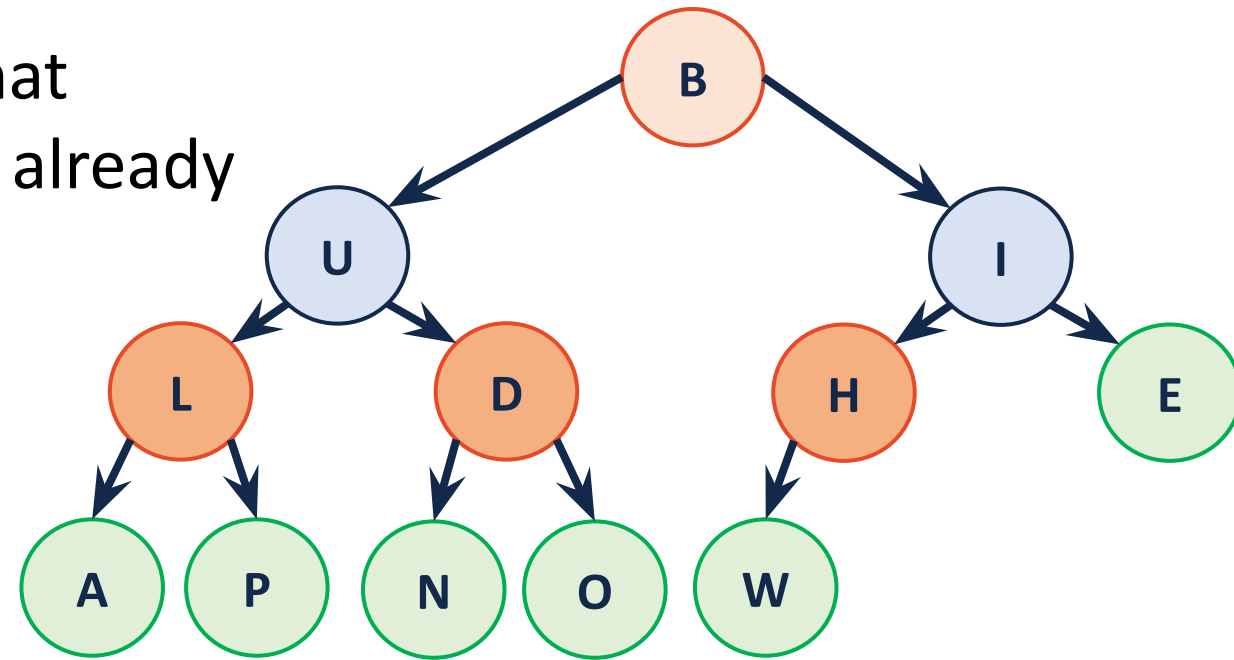




Build a Heap

buildHeap

We can take advantage of the fact that subtrees containing only a leaf node already satisfy heap property! ---- (1)



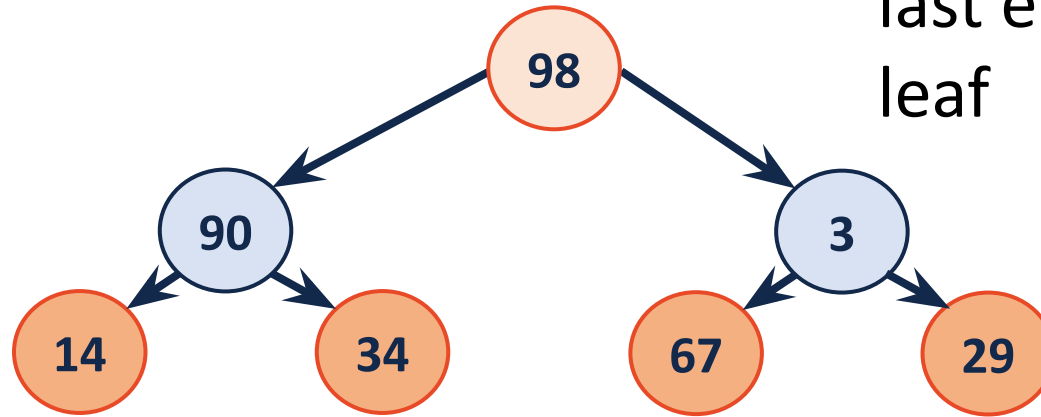
We have to restore heap properties in every subtree started from last element's parent (i), because of (1) we can call HeapfyDown on every element started from index i to root!



Worksheet Exercise #3

Exercise 3: After using the buildHeap algorithm on the array below, draw the corresponding tree.

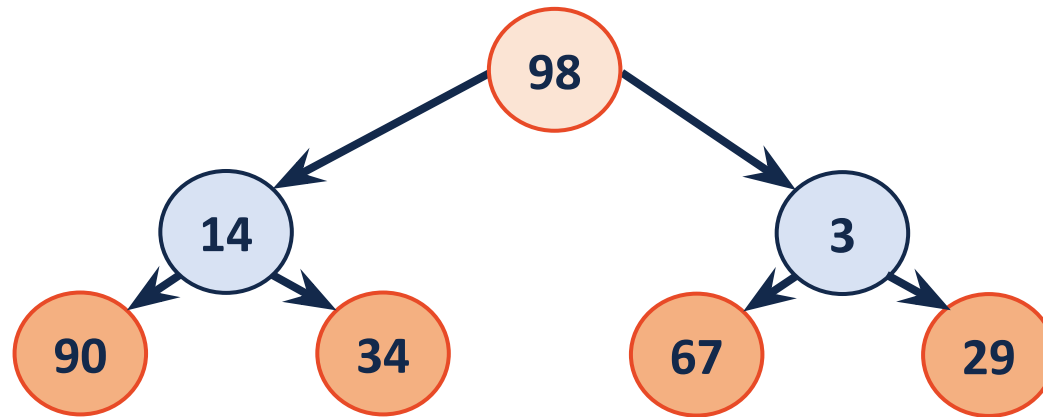
heapifyDown() from the last element that is not a leaf



The right half of the array (All of the nodes that are leaves) is already in a heap structure

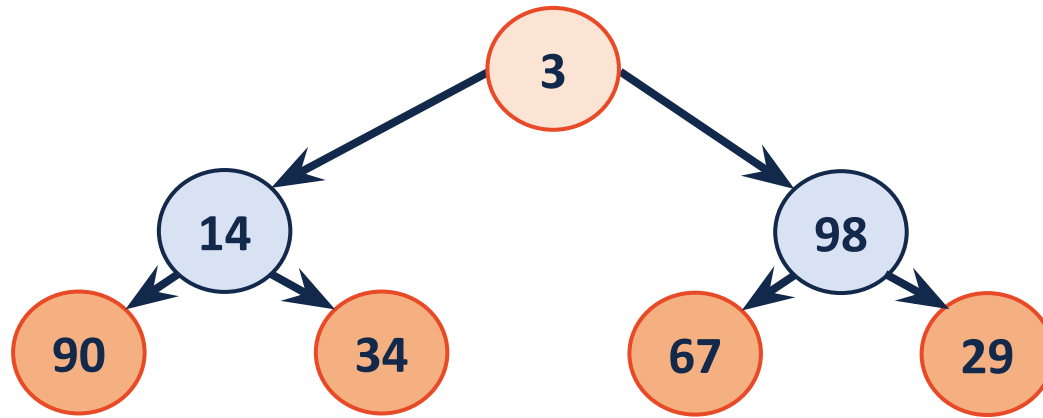
	98	90	3	14	34	67	29
--	----	----	---	----	----	----	----

Exercise 3: After using the buildHeap algorithm on the array above, draw the corresponding tree.



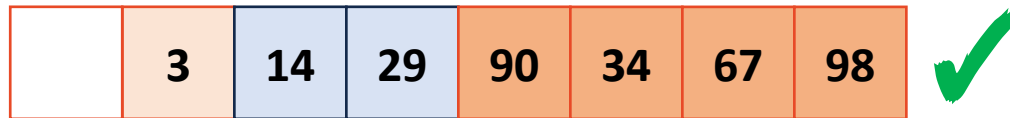
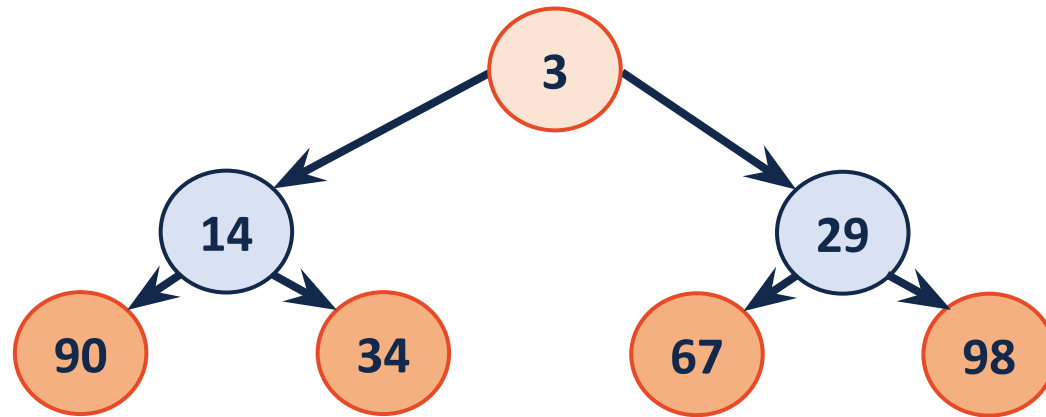
	98	14	3	90	34	67	29
--	----	----	---	----	----	----	----

Exercise 3: After using the buildHeap algorithm on the array above, draw the corresponding tree.



	3	14	98	90	34	67	29
--	---	----	----	----	----	----	----

Exercise 3: After using the buildHeap algorithm on the array above, draw the corresponding tree.



Helper function that restores the heap property by bubbling a node up the tree as necessary.
- already defined for you

heapifyUp(*i*)

```
if i != rootIndex && A[i] < A[parent(i)]  
    swap(i, parent(i))  
    heapifyUp(parent(i))
```

Helper function that restores the heap property by sinking a node down the tree as necessary.

heapifyDown(current)

```
If ! isLeaf(current)  
Find the min = index of min child of current  
If A[current] > A[min] child  
    swap A[current] and A[min]  
    HeapifyDown(min)
```