

## - Don't panic!

- Please print your name and NetID in each page in the appropriate fields, and circle your discussion section in the boxes above. We will return your exam at the indicated section.
- If you brought anything except your writing implements, your double-sided handwritten (in the original) $81 / 2^{\prime \prime} \times 11^{\prime \prime}$ cheat sheet, and your university ID, please put it away for the duration of the exam. In particular, you please turn off and put away all medically unnecessary electronic devices.
- Submit your cheat sheet together with your exam. An exam without your cheat sheet attached to it will not be graded.
- If you are NOT using a cheat sheet, please indicate so in large friendly letters on this page.
- Please ask for clarification if any question is unclear.
- This exam lasts 120 minutes. The clock started when you got the questions.
- If you run out of space for an answer, feel free to use the blank pages at the back of this booklet, but please tell us where to look.
- As usual, answering any (sub)problem with I don't know (and nothing else) is worth $25 \%$ partial credit. Correct, complete, but sub-optimal solutions are always worth more than $25 \%$. A blank answer is not the same as I don't know.
- Total IDK points for the whole exam would not exceed 10.
- Beware the Three Deadly Sins. Give complete solutions, not examples. Declare all your variables. If you don't know the answer admit it and use IDK.
- Style counts. Please use the backs of the pages or the blank pages at the end for scratch work, so that your actual answers are clear.
- Please return all paper with your answer booklet: your cheat sheet, and all scratch paper. We will return everything with your graded exam.


## - Good luck!

1 For each statement below, check "True" if the statement is always true and "False" otherwise. Each correct answer is worth 2 points; each incorrect answer is worth 0 points; and flipping a coin is (on average) worth 1 point. There is no IDK for this question.
1.A. For all languages $L$, if $L$ is not regular, then $L$ does not have False: $\square$ True: $\square$
a finite fooling set.
1.B. For all context-free languages $L$ and $L^{\prime}$, the language $L \cup L^{\prime}$ is $\quad$ False: $\square$ True: $\square$
also context-free.
1.C. $\left\{0^{i} 1^{j} 0^{k} 1^{\ell} \mid i \geq j \geq k \geq \ell \geq 0\right\}$ is not regular.

False:


True:


Let $M=(\Sigma, Q, s, A, \delta)$ and $M^{\prime}=(\Sigma, Q, s, Q \backslash A, \delta)$ be ar-
1.D. bitrary NFAs with identical alphabets, states, starting states, and transition functions, but with complementary accepting

False: $\square$ True: $\square$ states. Then $L(M) \backslash L\left(M^{\prime}\right)=L(M)$.

For all languages $L, L^{\prime} \subset \Sigma^{*}$, if $L$ and $L^{\prime}$ are recognized by
1.E. DFAs $M$ and $M^{\prime}$, respectively, then $L^{\prime} \backslash L$ can be represented by a regular expression.
1.F. If $L_{1}, L_{2}, \ldots$ are all context free languages, then $\bigcup_{i=0}^{\infty} L_{i}$ is context free.

Let $L$ be a regular language over alphabet $\Sigma$, and consider the language
1.G.

$$
L^{\prime}=\left\{x \alpha y \mid x, y \in \Sigma^{*}, \alpha \in \Sigma, \text { and } x y \in L\right\} .
$$

False: $\square$ True: $\square$

True: $\square$

False:


True: $\square$
language $L^{\prime}$ is not regular.
The strings 010 and 101 are distinguishable by the language
1.H.
$L=\left\{x \in \Sigma^{*} \mid\right.$ the number of 1 s in $x$ is even $\}$.
1.I. If a language $L \subseteq\{0\}^{*}$ then the language $L^{*}$ is regular.

False: $\square$ True: $\square$
1.J. Consider the logical statement "If pigs can fly then $2+2=5$."

This expression is:
False: $\square$ True: $\square$

2 For each of the following languages over the alphabet $\Sigma=\{0,1\}$, either prove that the language is regular or prove that the language is not regular. Exactly one of these two languages is regular. [This is a tricky question.]
2.A. $L=\left\{1^{n} w 0^{m-n} \mid w \in \Sigma^{+}\right.$and $\left.m \geq n>0\right\}$.
2.B. $L=\left\{x 0^{n} x^{R} \mid x \in \Sigma^{+}\right.$and $\left.n>0\right\}$, where $x^{R}$ is the reverse string of $x$.

3 For any language $L$, let $\operatorname{PS}(L)=\left\{x z \mid x y z \in L\right.$ for some $\left.x, y, z \in \Sigma^{*}\right\}$ be the language containing all strings in $L$ after one is allowed to delete some continuous middle portion. For example, if $L=\{000,10101\}$, then we have the string $10 \not 1 \notin 1=101 \in \operatorname{PS}(L)$. In this case, we have

$$
\operatorname{PS}(L)=\{\varepsilon, 0,00,000,10101,0101,1101,1001,1011,1010,101,01,11,10,1\} .
$$

Prove that for any regular language $L$, the language $\operatorname{PS}(L)$ is also regular. Your proof should be self contained, and not use any languages seen in class/lectures (suggestion: first describe the necessary construction, and then prove the correctness of the construction).

4 In the following, you do not need to prove that your answers are correct.
4.A. Provide a DFA and a regular expression for the following language: The set of all strings in $\{a, b\}^{*}$ that do not contain the substring aaaaa.
4.B. Provide a regular expression for the following language: The set of all strings in $\{a, b\}^{*}$ that contain both $a b$ and $a a$ as substrings.

5 A CFG $G$ is idiotic if all production rules in $G$ are of the form $A \rightarrow w B$ or $A \rightarrow \varepsilon$, where $w \in \Sigma^{+}$ and $A$ and $B$ are variables (i.e.., non-terminals) of $G$. Here $\Sigma=\{0\}$. For a given idiotic CFG grammar $G$, provide a construction that shows that the language of $L(G)$ is regular. Provide a convincing argument why this is true.
[No need to provide a formal proof here - it is too tedious for an exam.]
[If you do not know the answer - just use IDK - no need to waste your and our time.]

