

This is a review of context-free grammars from the lecture on Tuesday; in each example, the grammar itself is on the left; the explanation for each non-terminal is on the right.

- Properly nested strings of parentheses.

$$S \rightarrow \varepsilon \mid S(S) \quad \text{properly nested parentheses}$$

Here is a different grammar for the same language:

$$S \rightarrow \varepsilon \mid (S) \mid SS \quad \text{properly nested parentheses}$$

- $\{0^m 1^n \mid m \neq n\}$ . This is the set of all binary strings composed of some number of 0s followed by a different number of 1s.

$S \rightarrow A \mid B$	$\{0^m 1^n \mid m \neq n\}$
$A \rightarrow 0A \mid 0C$	$\{0^m 1^n \mid m > n\}$
$B \rightarrow B1 \mid C1$	$\{0^m 1^n \mid m < n\}$
$C \rightarrow \varepsilon \mid 0C1$	$\{0^m 1^n \mid m = n\}$

Give context-free grammars for each of the following languages. For each grammar, describe *in English* the language for each non-terminal, and in the examples above. As usual, we won't get to all of these in section.

1.  $\{0^{2n} 1^n \mid n \geq 0\}$

2.  $\{0^m 1^n \mid m \neq 2n\}$

[Hint: If  $m \neq 2n$ , then either  $m < 2n$  or  $m > 2n$ . Extend the previous grammar, but pay attention to parity. This language contains the string **01**.]

3.  $\{0, 1\}^* \setminus \{0^{2n} 1^n \mid n \geq 0\}$

[Hint: Extend the previous grammar. What's missing?]

**Work on these later:**

4.  $\{w \in \{0, 1\}^* \mid \#(0, w) = 2 \cdot \#(1, w)\}$  — Binary strings where the number of 0s is exactly twice the number of 1s.

5.  $\{0, 1\}^* \setminus \{ww \mid w \in \{0, 1\}^*\}$ .

[Anti-hint: The language  $\{ww \mid w \in \{0, 1\}^*\}$  is **not** context-free. Thus, the complement of a context-free language is not necessarily context-free!]

6. Prove that every regular language is context free.