

Give regular expressions for each of the following languages over the alphabet $\{0, 1\}$.

1. All strings containing the substring 000 .

Solution: $(0 + 1)^*000(0 + 1)^*$ ■

2. All strings *not* containing the substring 000 .

Solution: $(1 + 01 + 001)^*(\epsilon + 0 + 00)$ ■

Solution: $(\epsilon + 0 + 00)(1(\epsilon + 0 + 00))^*$ ■

3. All strings in which every run of 0 s has length at least 3.

Solution: $(1 + 0000^*)^*$ ■

Solution: $(\epsilon + 1)((\epsilon + 0000^*)1)^*(\epsilon + 0000^*)$ ■

4. All strings in which 1 does not appear after a substring 000 .

Solution: $(1 + 01 + 001)^*0^*$ ■

5. All strings containing at least three 0 s.

Solution: $(0 + 1)^*0(0 + 1)^*0(0 + 1)^*0(0 + 1)^*$ ■

Solution (clever): $1^*01^*01^*0(0 + 1)^*$ or $(0 + 1)^*01^*01^*01^*$ ■

6. Every string except 000 . [*Hint: Don't try to be clever.*]

Solution: Every string $w \neq 000$ satisfies one of three conditions: Either $|w| < 3$, or $|w| = 3$ and $w \neq 000$, or $|w| > 3$. The first two cases include only a finite number of strings, so we just list them explicitly. The last case includes *all* strings of length at least 4.

$$\begin{aligned} & \epsilon + 0 + 1 + 00 + 01 + 10 + 11 \\ & + 001 + 010 + 011 + 100 + 101 + 110 + 111 \\ & + (1 + 0)(1 + 0)(1 + 0)(1 + 0)(1 + 0)^* \end{aligned}$$

Solution (clever): $\epsilon + 0 + 00 + (1 + 01 + 001 + 000(1 + 0))(1 + 0)^*$ ■

7. All strings w such that *in every prefix of w* , the number of **0**s and **1**s differ by at most 1.

Solution: Equivalently, strings that alternate between **0**s and **1**s: $(\mathbf{01} + \mathbf{10})^*(\varepsilon + \mathbf{0} + \mathbf{1})$ ■

*8. All strings containing at least two **0**s and at least one **1**.

Solution: There are three possibilities for how such a string can begin:

- Start with **00**, then any number of **0**s, then **1**, then anything.
- Start with **01**, then any number of **1**s, then **0**, then anything.
- Start with **1**, then a substring with exactly two **0**s, then anything.

All together: $\mathbf{000}^*\mathbf{1}(\mathbf{0} + \mathbf{1})^* + \mathbf{011}^*\mathbf{0}(\mathbf{0} + \mathbf{1})^* + \mathbf{11}^*\mathbf{01}^*\mathbf{0}(\mathbf{0} + \mathbf{1})^*$

Or equivalently: $(\mathbf{000}^*\mathbf{1} + \mathbf{011}^*\mathbf{0} + \mathbf{11}^*\mathbf{01}^*\mathbf{0})(\mathbf{0} + \mathbf{1})^*$ ■

Solution: There are three possibilities for how the three required symbols are ordered:

- Contains a **1** before two **0**s: $(\mathbf{0} + \mathbf{1})^*\mathbf{1}(\mathbf{0} + \mathbf{1})^*\mathbf{0}(\mathbf{0} + \mathbf{1})^*\mathbf{0}(\mathbf{0} + \mathbf{1})^*$
- Contains a **1** between two **0**s: $(\mathbf{0} + \mathbf{1})^*\mathbf{0}(\mathbf{0} + \mathbf{1})^*\mathbf{1}(\mathbf{0} + \mathbf{1})^*\mathbf{0}(\mathbf{0} + \mathbf{1})^*$
- Contains a **1** after two **0**s: $(\mathbf{0} + \mathbf{1})^*\mathbf{0}(\mathbf{0} + \mathbf{1})^*\mathbf{0}(\mathbf{0} + \mathbf{1})^*\mathbf{1}(\mathbf{0} + \mathbf{1})^*$

So putting these cases together, we get the following:

$$\begin{aligned} & (\mathbf{0} + \mathbf{1})^*\mathbf{1}(\mathbf{0} + \mathbf{1})^*\mathbf{0}(\mathbf{0} + \mathbf{1})^*\mathbf{0}(\mathbf{0} + \mathbf{1})^* \\ & + (\mathbf{0} + \mathbf{1})^*\mathbf{0}(\mathbf{0} + \mathbf{1})^*\mathbf{1}(\mathbf{0} + \mathbf{1})^*\mathbf{0}(\mathbf{0} + \mathbf{1})^* \\ & + (\mathbf{0} + \mathbf{1})^*\mathbf{0}(\mathbf{0} + \mathbf{1})^*\mathbf{0}(\mathbf{0} + \mathbf{1})^*\mathbf{1}(\mathbf{0} + \mathbf{1})^* \end{aligned}$$
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Solution (clever): $(\mathbf{0} + \mathbf{1})^*(\mathbf{101}^*\mathbf{0} + \mathbf{010} + \mathbf{01}^*\mathbf{01})(\mathbf{0} + \mathbf{1})^*$ ■

*9. All strings w such that *in every prefix of w* , the number of **0**s and **1**s differ by at most 2.

Solution: $(\mathbf{0}(\mathbf{01})^*\mathbf{1} + \mathbf{1}(\mathbf{10})^*\mathbf{0})^* \cdot (\varepsilon + \mathbf{0}(\mathbf{01})^*(\mathbf{0} + \varepsilon) + \mathbf{1}(\mathbf{10})^*(\mathbf{1} + \varepsilon))$ ■

- ★10. All strings in which the substring **000** appears an even number of times.
(For example, **0001000** and **0000** are in this language, but **00000** is not.)

Solution: Every string in $\{0, 1\}^*$ alternates between (possibly empty) blocks of **0**s and individual **1**s; that is, $\{0, 1\}^* = (0^*1)^*0^*$. Trivially, every **000** substring is contained in some block of **0**s. Our strategy is to consider which blocks of **0**s contain an even or odd number of **000** substrings.

Let X denote the set of all strings in 0^* with an even number of **000** substrings. We easily observe that $X = \{0^n \mid n = 1 \text{ or } n \text{ is even}\} = 0 + (00)^*$.

Let Y denote the set of all strings in 0^* with an *odd* number of **000** substrings. We easily observe that $Y = \{0^n \mid n > 1 \text{ and } n \text{ is odd}\} = 000(00)^*$.

We immediately have $0^* = X + Y$ and therefore $\{0, 1\}^* = ((X + Y)1)^*(X + Y)$.

Finally, let L denote the set of all strings in $\{0, 1\}^*$ with an even number of **000** substrings. A string $w \in \{0, 1\}^*$ is in L if and only if an odd number of blocks of **0**s in w are in Y ; the remaining blocks of **0**s are all in X .

$$L = ((X1)^*Y1 \cdot (X1)^*Y1)^*(X1)^*X$$

Plugging in the expressions for X and Y gives us the following regular expression for L :

$$\left(((0 + (00)^*)1)^* \cdot 000(00)^*1 \cdot ((0 + (00)^*)1)^* \cdot 000(00)^*1 \right)^* \cdot ((0 + (00)^*)1)^* \cdot (0 + (00)^*)$$

Whew! ■