

IDK

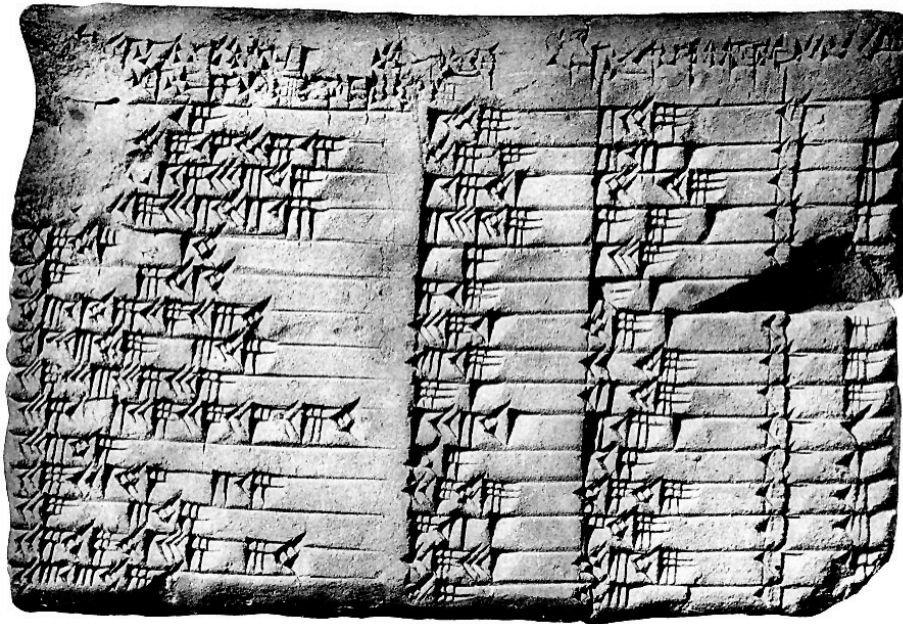
Deadly sins

No proofs by example

Declare your variables

NO WEAK INDUCTION

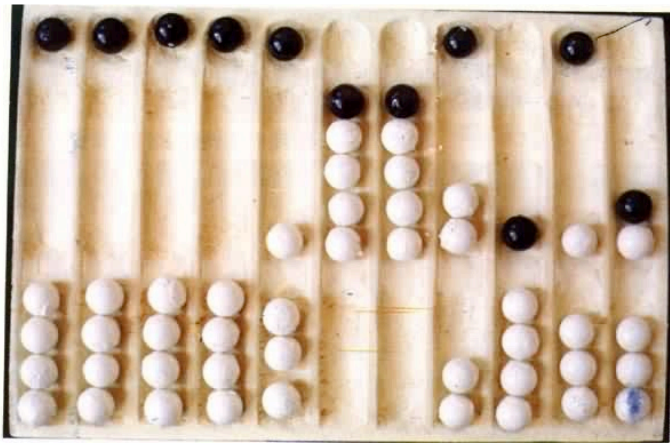
Cheating ~ Don't!



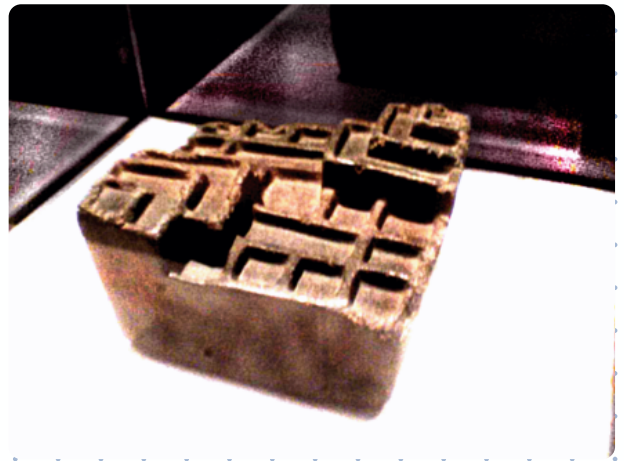
Plimpton 322 (Mesopotamia, 1800 BCE)



Rhind Papyrus (Egypt, 1550 BCE)



Roman abacus



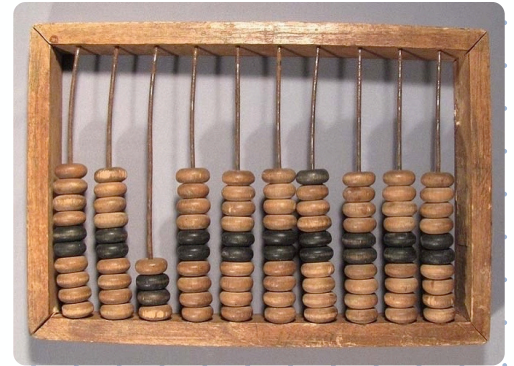
Incan Yupana



Suan Pan

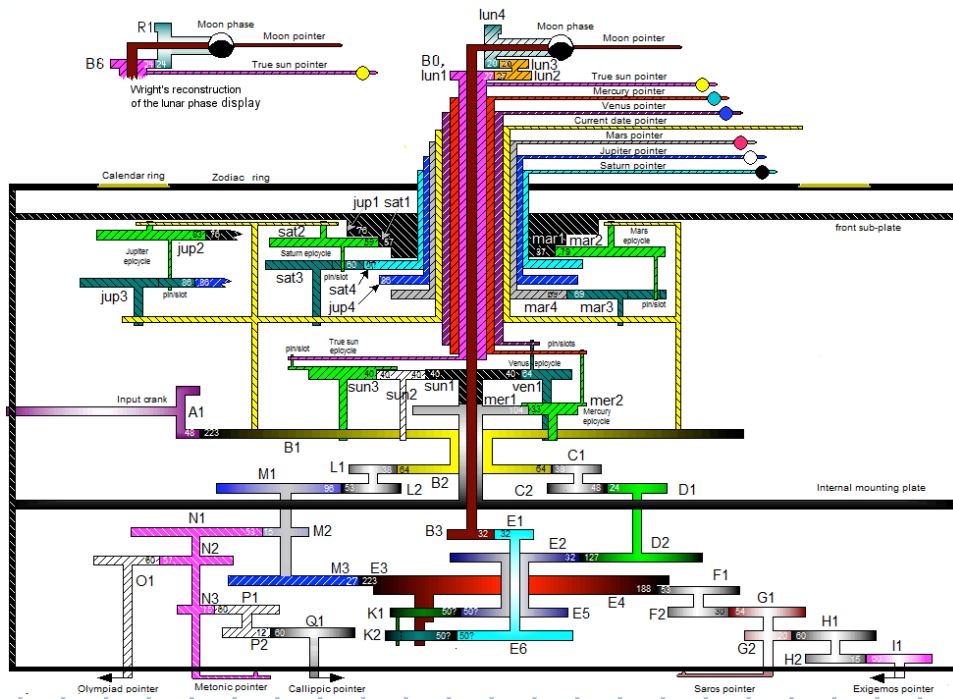


Soroban

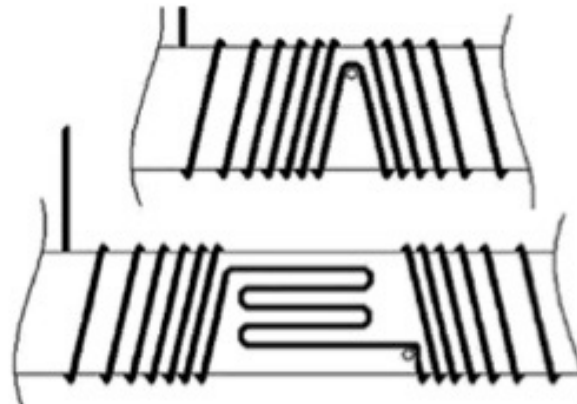
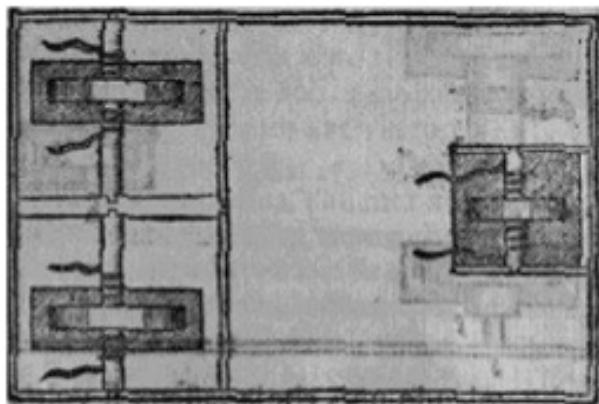


Shoty

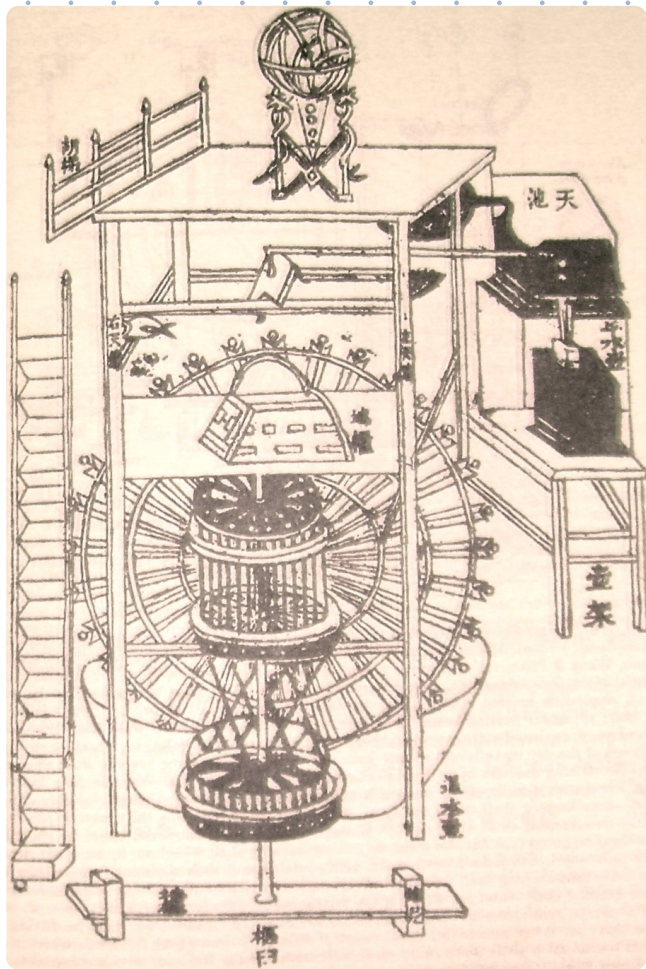




Antikythera Mechanism



Heron of Alexandria's programmable cart



Su Song's "cosmic engine" (1088CE)



↑↑↑↑↑↑

Pask-Tiſtu, Sara-Pick, Cykeln af 532 år

År	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21	22	23	24	25	26	27	28
1	♯	♯	♯	♯	♯	♯	♯	♯	♯	♯	♯	♯	♯	♯	♯	♯	♯	♯	♯	♯	♯	♯	♯	♯	♯	♯	♯	♯
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Påsk-Tiſtu. Sara-Pick. Cykeln af 532 år

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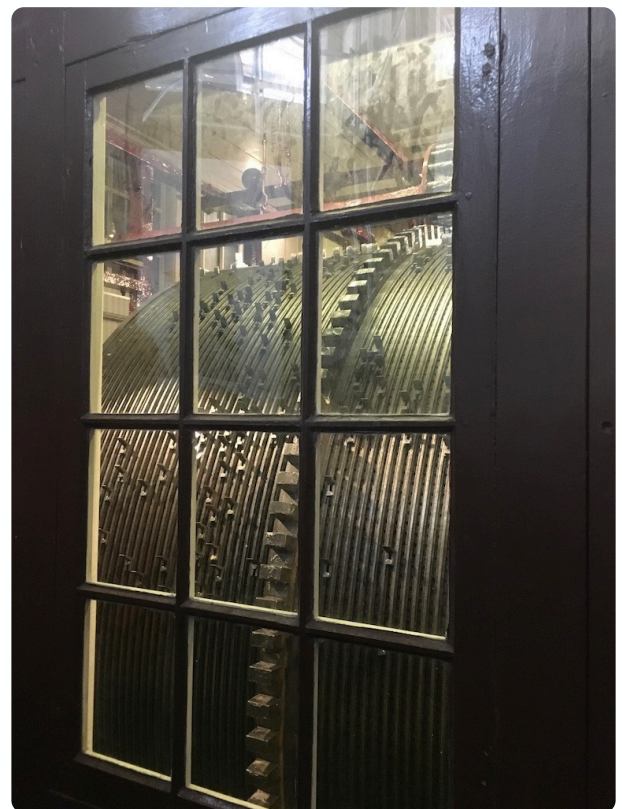
Computus tables (12th C. Sweden)



Strasbourg astronomical clock (16th C.)



Automatic carillon (15th C.)



Utrecht Domtoren
automatic carillon (1975)



Leibniz 1684
(also Newton 1687)

DERIVATIVE DEFINITION

$$\frac{d}{dx}(f(x)) = f'(x) = \lim_{h \rightarrow 0} \frac{f(x+h) - f(x)}{h}$$

BASIC PROPERTIES

$$(cf(x))' = c(f'(x))$$

$$(f(x) \pm g(x))' = f'(x) \pm g'(x)$$

$$\frac{d}{dx}(c) = 0$$

MEAN VALUE THEOREM

If f is differentiable on the interval (a, b) and continuous at the end points there exists a c in (a, b) such that

$$f'(c) = \frac{f(b) - f(a)}{b - a}$$

PRODUCT RULE

$$(f(x)g(x))' = f(x)'g(x) + f(x)g(x)'$$

QUOTIENT RULE

$$\frac{d}{dx} \left(\frac{f(x)}{g(x)} \right) = \frac{f'(x)g(x) - f(x)g'(x)}{[g(x)]^2}$$

POWER RULE

$$\frac{d}{dx}(x^n) = nx^{n-1}$$

CHAIN RULE

$$\frac{d}{dx}(f(g(x))) = f'(g(x))g'(x)$$

COMMON DERIVATIVES

$$\frac{d}{dx}(x) = 1$$

$$\frac{d}{dx}(\sin x) = \cos x$$

$$\frac{d}{dx}(\cos x) = -\sin x$$

$$\frac{d}{dx}(\tan x) = \sec^2 x$$

$$\frac{d}{dx}(\sec x) = \sec x \tan x$$

$$\frac{d}{dx}(\csc x) = -\csc x \cot x$$

$$\frac{d}{dx}(\cot x) = -\csc^2 x$$

$$\frac{d}{dx}(\sin^{-1} x) = \frac{1}{\sqrt{1-x^2}}$$

$$\frac{d}{dx}(\cos^{-1} x) = -\frac{1}{\sqrt{1-x^2}}$$

$$\frac{d}{dx}(\tan^{-1} x) = \frac{1}{1+x^2}$$

$$\frac{d}{dx}(a^x) = a^x \ln(a)$$

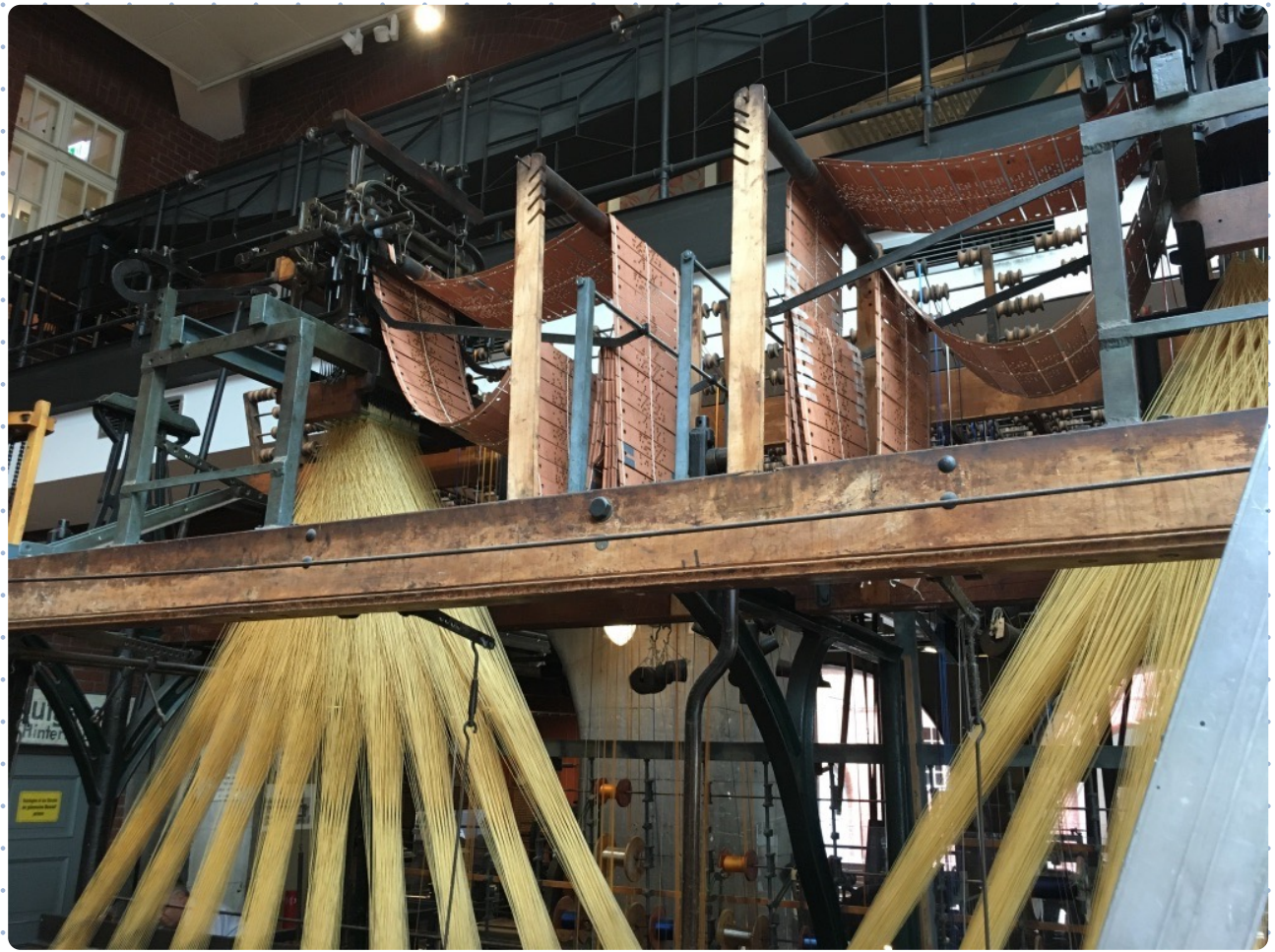
$$\frac{d}{dx}(e^x) = e^x$$

$$\frac{d}{dx}(\ln(x)) = \frac{1}{x}, x > 0$$

$$\frac{d}{dx}(\ln|x|) = \frac{1}{x}$$

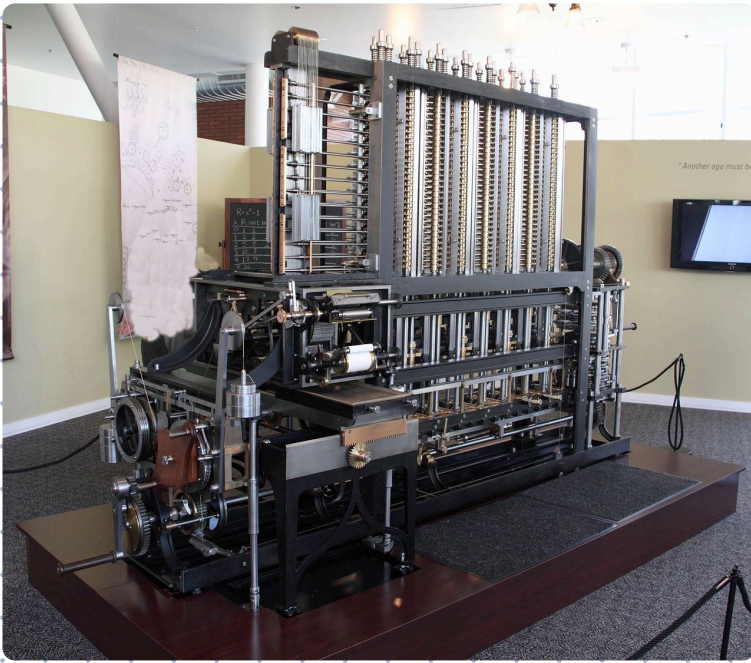
$$\frac{d}{dx}(\log_a(x)) = \frac{1}{x \ln(a)}$$

... equ. (con. ...)
 dinatae respondenti curvae VV) erit dy æqu. dv . Jam *Additio & Sub-*
tractio: si sit $z - y \dagger vv \dagger x$ æqu. v , erit $dz - y \dagger vv \dagger x$ seu dv , æqu.
 $dz - dy \dagger dvv \dagger dx$. *Multiplicatio*, $dx v$ æqu. $x dv \dagger v dx$, seu posito
 y æqu. xv , fiet dy æqu. $x dv \dagger v dx$. In arbitrio enim est vel formulam,
 ... Notandum & ...

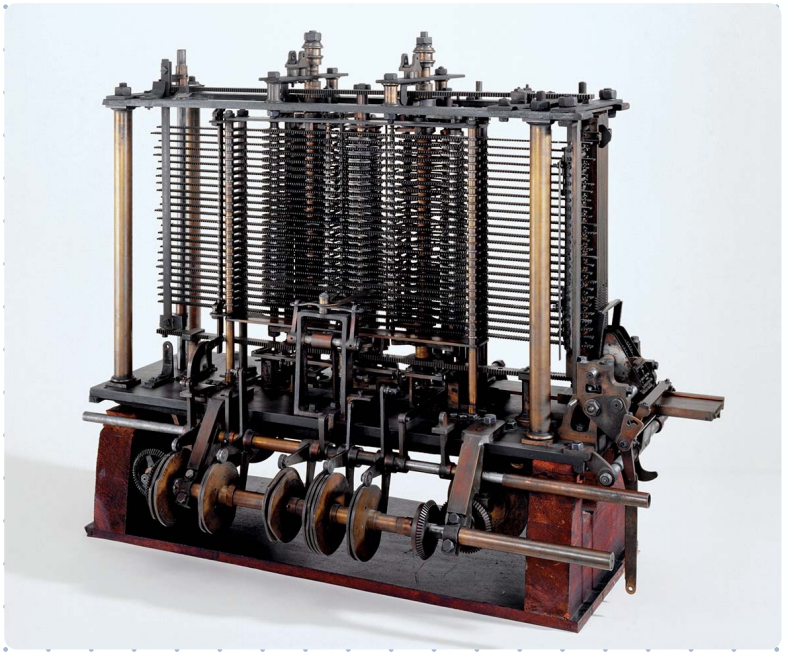


Jacquard loom (1804)

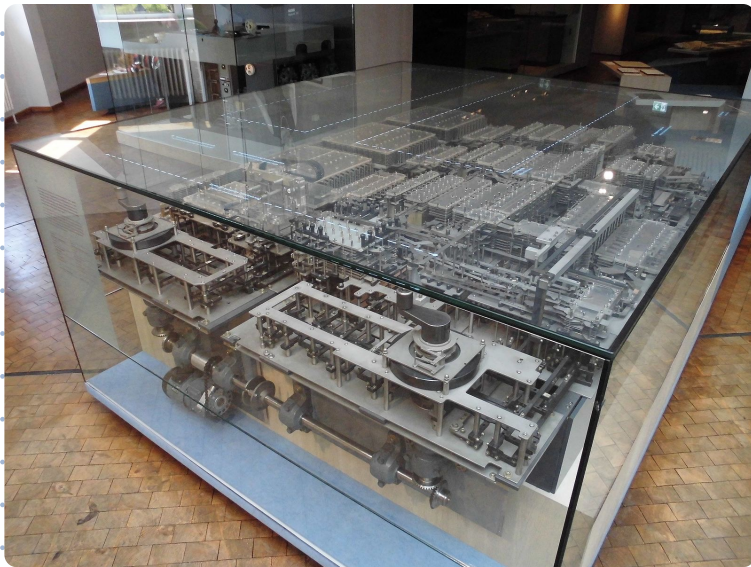




Babbage difference engine
(designed 1822)



Babbage analytical engine
(partially designed by 1877)



Zuse Z1 (1936-38)



Programmsteuerung mit Lochstreifen
program control using punched tape



One of Zuse's
mechanical relays

COLLATE(n):

if $n=1$
return TRUE
else if n is even
 $n \leftarrow n/2$
else
 $n \leftarrow 3n+1$

Sequences

↓
Strings

Let Σ be any finite set $\{0, 1\}$

A string is either

- empty ϵ
- $a \cdot x$ for some symbol $a \in \Sigma$
and some string x

STRING = S.TRING
= S.(T.TRING) = ...

Length $|w|$ of a string w is...

$$|w| = \begin{cases} 0 & \text{if } w = \epsilon \\ 1 + |x| & \text{if } w = a \cdot x \end{cases}$$

Concatenation $w \cdot z$

$$w \cdot z = \begin{cases} z & \text{if } w = \epsilon \\ a \cdot (x \cdot z) & \text{if } w = a \cdot x \end{cases}$$

HEAD · ACHE
 $\underbrace{\quad}_a \quad \underbrace{\quad}_x \quad \underbrace{\quad}_z$

Theorem: $|w \cdot z| = |w| + |z|$ for all strings w and z

Proof: Let w and z be arbitrary strings.

IH: Assume $|x \cdot z| = |x| + |z|$ for all strings x shorter than w .

There are two cases:

$$\begin{aligned} \bullet w = \varepsilon &\Rightarrow |w \cdot z| = |\varepsilon \cdot z| \\ &= |z| \\ &= |\varepsilon| + |z| \\ &= |w| + |z| \end{aligned}$$

$w = \varepsilon$
by def.
by def ||
because $w = \varepsilon$

$$\begin{aligned} \bullet w = ax &\Rightarrow |w \cdot z| = |ax \cdot z| \\ &= |a \cdot (x \cdot z)| \\ &= 1 + |x \cdot z| \\ &\quad \downarrow ? \\ &= 1 + |x| + |z| \\ &= |ax| + |z| \\ &= |w| + |z| \end{aligned}$$

$w = ax$
def.
by def ||

by IH
def ||
 $w = ax$

So $|w \cdot z| = |w| + |z|$