
Submission instructions as in previous [homeworks](#).

4 (100 PTS.) Regular expressions.

For each of the following languages over the alphabet $\{0, 1\}$, give a regular expression that describes that language, and briefly argue why your expression is correct.

- 4.A. (10 PTS.) All strings that end in **1011**.
- 4.B. (10 PTS.) All strings except **11**.
- 4.C. (10 PTS.) All strings that contain **101** or **010** as a substring.
- 4.D. (10 PTS.) All strings that contain **111** and **000** as a subsequence (the resulting expression is long – describe how you got your expression, instead of writing it out explicitly).
- 4.E. (10 PTS.) The language containing all strings that do not contain **111** as a substring.
- 4.F. (10 PTS.) All strings that do *not* contain **000** as a subsequence.
- 4.G. (10 PTS.) Strings in which every occurrence of the substring **00** appears before every occurrence of the substring **11**.
- 4.H. (10 PTS.) Strings that do not contain the subsequence **010**.
- 4.I. (10 PTS.) Strings that do not contain the subsequence **0101010**.
- 4.J. (10 PTS.) Strings that do not contain the subsequence **10**.
- 4.K. (Not for credit, do not submit a solution.) Strings that do not contain the subsequence **111000**.

5 (100 PTS.) DFA I

Let $\Sigma = \{0, 1\}$. Let L be the set of all strings in Σ^* that contain an even number of **0**s and an even number of **1**s.

- 5.A. (50 PTS.) Describe a DFA over Σ that accepts the language L . Argue that your machine accepts every string in L and nothing else, by explaining what each state in your DFA *means*. (Hint: Zero is even)
You may either draw the DFA or describe it formally, but the states Q , the start state s , the accepting states A , and the transition function δ must be clearly specified, in either case.
- 5.B. (50 PTS.) (Harder.) Give a regular expression for L , and briefly argue why the expression is correct. (Hint: First solve the much easier case where the strings do not contain any consecutive **0**s or **1**s.)

6 (100 PTS.) DFA II

Let L_1, L_2 , and L_3 be regular languages over Σ accepted by DFAs $M_1 = (Q_1, \Sigma, \delta_1, s_1, A_1)$, $M_2 = (Q_2, \Sigma, \delta_2, s_2, A_2)$, and $M_3 = (Q_3, \Sigma, \delta_3, s_3, A_3)$, respectively.

- 6.A.** (20 PTS.) Describe formally the product construction of the DFA M that accepts the language $L_1 \cap L_2 \cap L_3$.
- 6.B.** (30 PTS.) In the DFA M constructed in (??), a state is a triple (q_1, q_2, q_3) . Let δ the transition function of M , and let δ^* be the standard extension of δ to strings. Prove by induction that for any string $w \in \Sigma^*$, we have that

$$\delta^*((q_1, q_2, q_3), w) = (\delta_1^*(q_1, w), \delta_2^*(q_2, w), \delta_3^*(q_3, w)).$$

- 6.C.** (20 PTS.) Describe a DFA $M = (Q, \Sigma, \delta, s, A)$ in terms of M_1, M_2 , and M_3 that accepts $L = \{w \mid w \text{ is in exactly two of } \{L_1, L_2, L_3\}\}$. Formally specify the components Q, δ, s , and A for M in terms of the components of M_1, M_2 , and M_3 . Argue that your construction is correct.
- 6.D.** (30 PTS.) You are given a DFA $M = (Q, \Sigma, \delta, s, A)$, for $\Sigma = \{0, 1\}$. Describe in detail how to build a DFA that accepts the language

$$L = \{w \in \Sigma^* \mid w \notin L(M), \bar{w} \in L(M) \text{ and } 1^{|w|} \in L(M)\}.$$

How many states does your DFA has as a function of $n = |Q|$? Argue that the DFA you constructed indeed accepts the specified language.

Here, for $w = w_1w_2 \dots w_m \in \Sigma^*$, the *complement string* \bar{w} is $\bar{w}_1\bar{w}_2\bar{w}_3 \dots \bar{w}_m$, where $\bar{0} = 1$, and $\bar{1} = 0$.