Deterministic Finite Automata (DFAs)

Lecture 3

Tuesday, January 22, 2019

LaTeXed: December 27, 2018 08:25

Part I

DFA Introduction

DFAs also called Finite State Machines (FSMs)

- The "simplest" model for computers?
- State machines that are common in practice.
  - Vending machines
  - Elevators
  - Digital watches
  - Simple network protocols
- Programs with fixed memory

A simple program

Program to check if a given input string w has odd length

```
int n = 0
While input is not finished
    read next character c
    n ← n + 1
endWhile
If (n is odd) output YES
Else output NO
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bit x = 0
While input is not finished
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    x ← flip(x)
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If (x = 1) output YES
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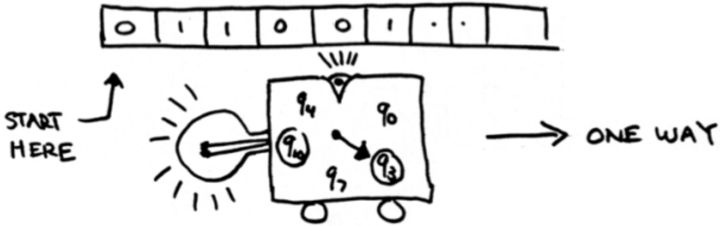
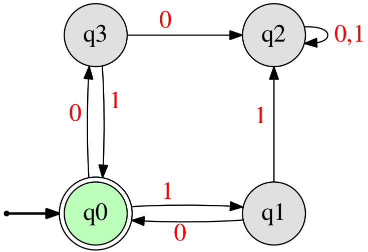
Another view



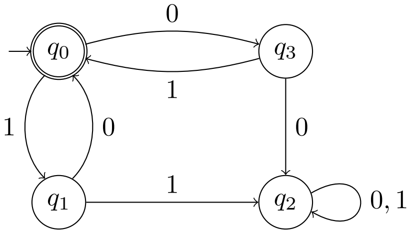
- Machine has input written on a *read-only* tape
- Start in specified start state
- Start at left, scan symbol, change state and move right
- Circled states are *accepting*
- Machine *accepts* input string if it is in an accepting state after scanning the last symbol.

Graphical Representation/State Machine



- Directed graph with nodes representing states and edge/arcs representing transitions labeled by symbols in Σ
- For each state (vertex) q and symbol $a \in \Sigma$ there is *exactly* one outgoing edge labeled by a
- Initial/start state has a pointer (or labeled as s, q_0 or “start”)
- Some states with double circles labeled as accepting/final states

Graphical Representation



- Where does **001** lead? **10010**?
- Which strings end up in accepting state?
- Can you prove it?
- Every string w has a unique walk that it follows from a given state q by reading one letter of w from left to right.

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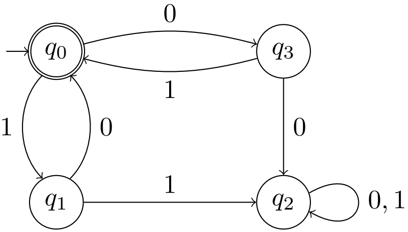

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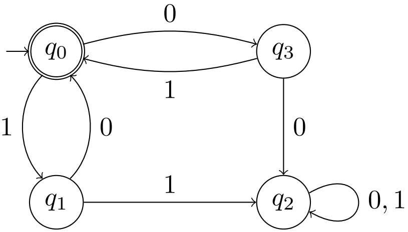

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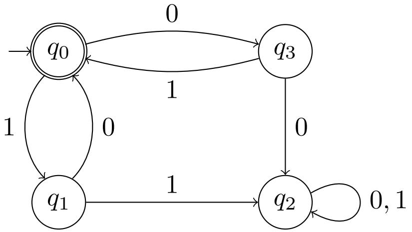

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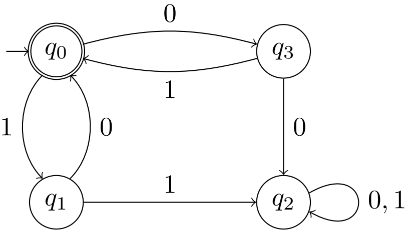

Definition

A DFA M accepts a string w iff the unique walk starting at the start state and spelling out w ends in an accepting state.

Definition

The language accepted (or recognized) by a DFA M is denote by $L(M)$ and defined as: $L(M) = \{w \mid M \text{ accepts } w\}$.

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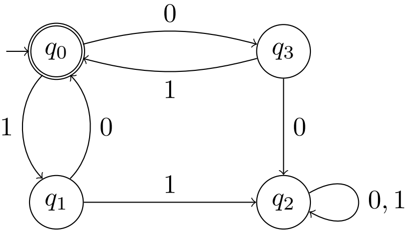

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"M accepts language L" does not mean simply that that M accepts each string in L.

It means that M accepts each string in L and no others. Equivalently M accepts each string in L and does not accept/rejects strings in $\Sigma^* \setminus L$.

M "recognizes" L is a better term but "accepts" is widely accepted (and recognized) (joke attributed to Lenny Pitt)

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Formal Tuple Notation

Definition

A deterministic finite automata (DFA) $M = (Q, \Sigma, \delta, s, A)$ is a five tuple where

- Q is a finite set whose elements are called states,
- Σ is a finite set called the input alphabet,
- $\delta : Q \times \Sigma \to Q$ is the transition function,
- $s \in Q$ is the start state,
- $A \subseteq Q$ is the set of accepting/final states.

Common alternate notation: q_0 for start state, F for final states.

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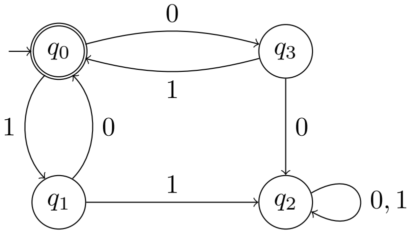
DFA Notation

$$M = \Big(\overbrace{Q}^{\text{set of all states}}, \underbrace{\Sigma}_{\text{alphabet}}, \overbrace{\delta}^{\text{transition func}}, \underbrace{s}_{\text{start state}}, \overbrace{A}^{\text{set of all accept states}} \Big)$$

Example



- $Q = \{q_0, q_1, q_1, q_3\}$
- $\Sigma = \{0, 1\}$
- δ
- $s = q_0$
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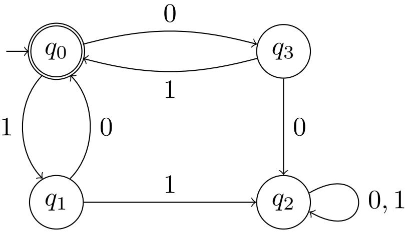

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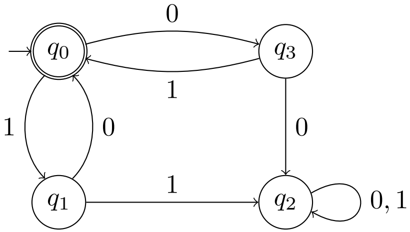

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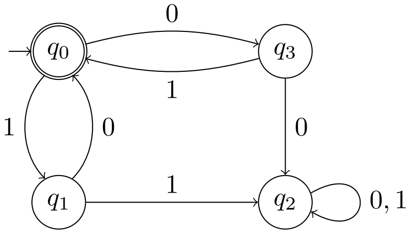

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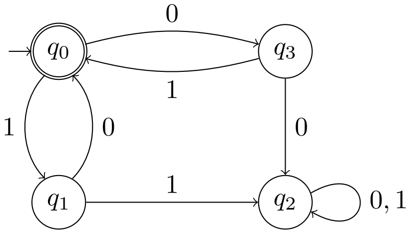

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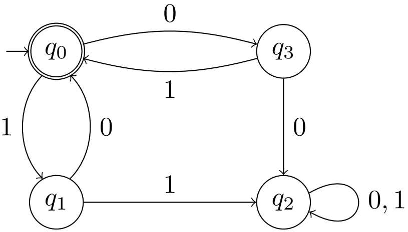

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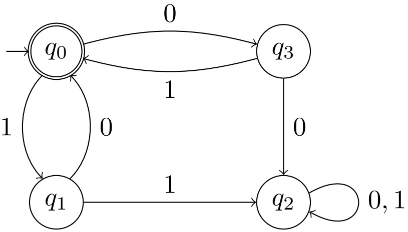

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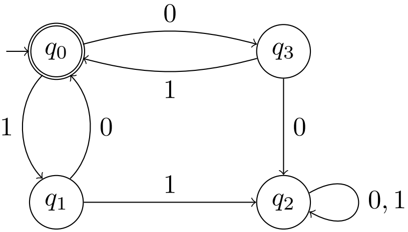

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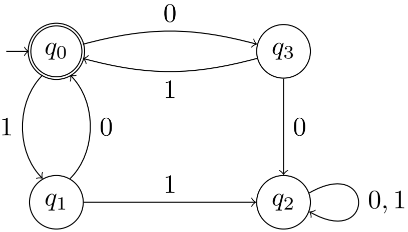

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Extending the transition function to strings

Given DFA $M = (Q, \Sigma, \delta, s, A)$, $\delta(q, a)$ is the state that M goes to from q on reading letter a

Useful to have notation to specify the unique state that M will reach from q on reading *string* w

Transition function $\delta^* : Q \times \Sigma^* \to Q$ defined inductively as follows:

- $\delta^*(q, w) = q$ if $w = \epsilon$
- $\delta^*(q, w) = \delta^*(\delta(q, a), x)$ if $w = ax$.

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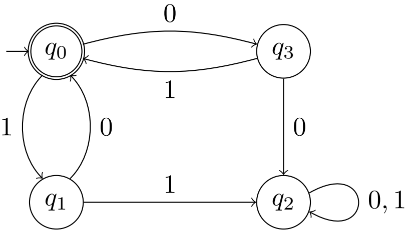
Formal definition of language accepted by M

Definition

The language $L(M)$ accepted by a DFA $M = (Q, \Sigma, \delta, s, A)$ is

$$\{w \in \Sigma^* \mid \delta^*(s, w) \in A\}.$$

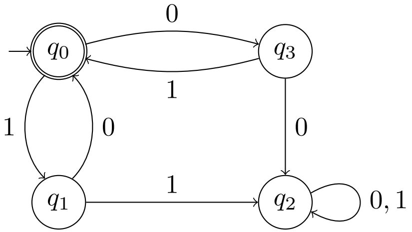
Example



What is:

- $\delta^*(q_1, \epsilon)$
- $\delta^*(q_0, 1011)$
- $\delta^*(q_1, 010)$
- $\delta^*(q_4, 10)$

Example continued



- What is $L(M)$ if start state is changed to q_1?
- What is $L(M)$ if final/accept states are set to $\{q_2, q_3\}$ instead of $\{q_0\}$?

Advantages of formal specification

- Necessary for proofs
- Necessary to specify abstractly for class of languages

Exercise: Prove by induction that for any two strings $\boldsymbol{u}, \boldsymbol{v}$, any state $\boldsymbol{q}$, $\delta^*(q, uv) = \delta^*(\delta^*(q, u), v)$.

Part II

Constructing DFAs

DFAs: State = Memory

How do we design a DFA M for a given language L? That is $L(M) = L$.

- DFA is a like a program that has fixed amount of memory independent of input size.
- The memory of a DFA is encoded in its states
- The state/memory must capture enough information from the input seen so far that it is sufficient for the suffix that is yet to be seen (note that DFA cannot go back)

DFA Construction: Example

Assume $\Sigma = \{0, 1\}$

- $L = \emptyset$, $L = \Sigma^*$, $L = \{\epsilon\}$, $L = \{0\}$.
- $L = \{w \in \{0, 1\}^* \mid |w| \text{ is divisible by } 5\}$
- $L = \{w \in \{0, 1\}^* \mid w \text{ ends with } 01\}$
- $L = \{w \in \{0, 1\}^* \mid w \text{ contains } 001 \text{ as substring}\}$
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DFA Construction: Example

$L = \{$Binary numbers congruent to $\mathbf{0} \mod \mathbf{5}\}$

Example: $\mathbf{1101011} = \mathbf{107} = \mathbf{2} \mod \mathbf{5}$, $\mathbf{1010} = \mathbf{10} = \mathbf{0} \mod \mathbf{5}$

Key observation:

$w0 \mod 5 = a$ implies

$w0 \mod 5 = 2a \mod 5$ and $w1 \mod 5 = (2a + 1) \mod 5$

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Part III

Product Construction and Closure Properties

Part IV

Complement

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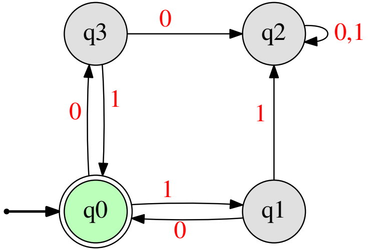
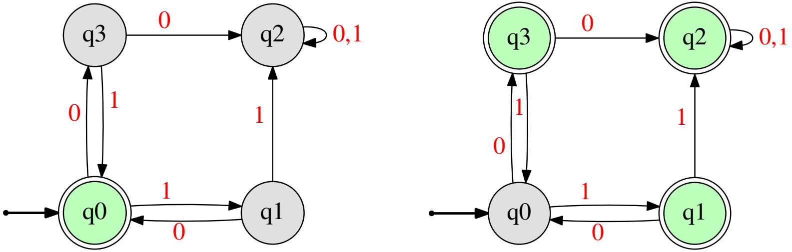
Question: If M is a DFA, is there a DFA M' such that $L(M') = \Sigma^* \setminus L(M)$? That is, are languages recognized by DFAs closed under complement?



Complement

Example...

Just flip the state of the states!



Complement

Theorem

Languages accepted by DFA*s are closed under complement.*

Proof.

Let $M = (Q, \Sigma, \delta, s, A)$ such that $L = L(M)$.
Let $M' = (Q, \Sigma, \delta, s, Q \setminus A)$. Claim: $L(M') = \bar{L}$. Why?
$\delta^*_M = \delta^*_{M'}$. Thus, for every string w, $\delta^*_M(s, w) = \delta^*_{M'}(s, w)$.
$\delta^*_M(s, w) \in A \Rightarrow \delta^*_{M'}(s, w) \notin Q \setminus A$.
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Part V

Product Construction

Union and Intersection

Question: Are languages accepted by DFAs closed under union?
That is, given DFAs M_1 and M_2 is there a DFA that accepts $L(M_1) \cup L(M_2)$?
How about intersection $L(M_1) \cap L(M_2)$?

Idea from programming: on input string w

- Simulate M_1 on w
- Simulate M_2 on w
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Question: Are languages accepted by DFAs closed under union?
That is, given DFAs M_1 and M_2 is there a DFA that accepts $L(M_1) \cup L(M_2)$?
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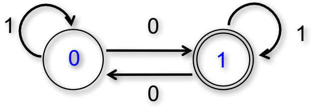
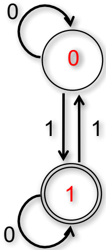
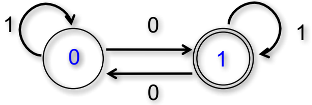
Example



M_1 accepts #0 = odd



M_2 accepts #1 = odd



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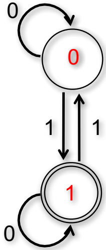
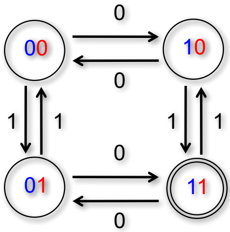


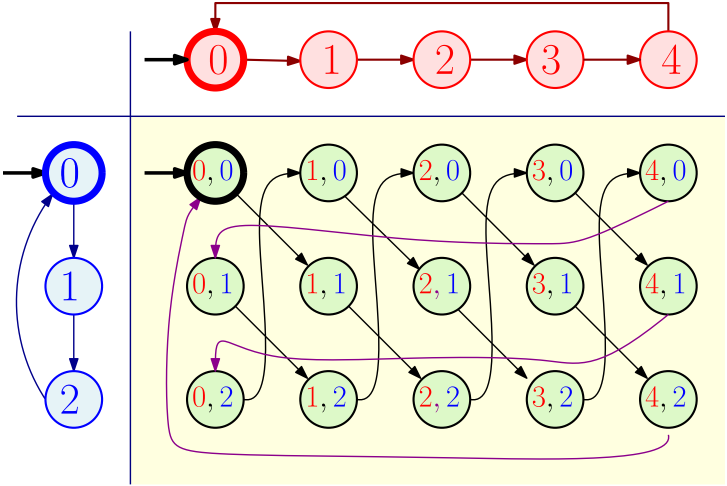
M_2 accepts #1 = odd



Cross-product machine

Example II

Accept all binary strings of length divisible by **3** and **5**



Assume all edges are labeled by $\mathbf{0, 1}$.

Product construction for intersection

$M_1 = (Q_1, \Sigma, \delta_1, s_1, A_1)$ and $M_2 = (Q_1, \Sigma, \delta_2, s_2, A_2)$

Create $M = (Q, \Sigma, \delta, s, A)$ where

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$L(M) = L(M_1) \cap L(M_2)$.

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Set Difference

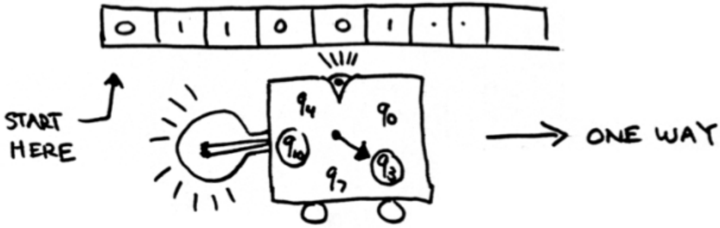
Theorem

M_1, M_2 DFAs. There is a DFA M such that
$L(M) = L(M_1) \setminus L(M_2)$.

Exercise: Prove the above using two methods.

- Using a direct product construction
- Using closure under complement and intersection and union

Things to know: 2-way DFA



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- Can show that any language recognized by a 2-way DFA can be recognized by a regular (1-way) DFA
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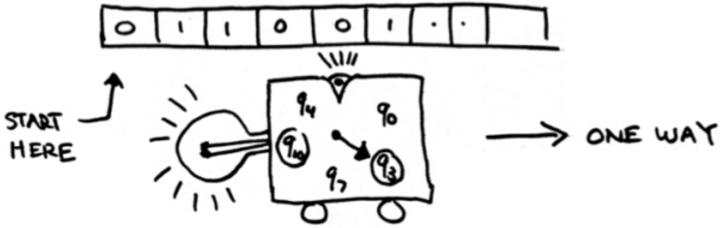

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