

Lecture 8 Scribble
09/17/2020

- Topics for today:
- First CFG with closure
 - Intro to Turing machines
 - Formal def TMs
 - TM examples
 - Universal Turing Machines

CFLs/CFGs $L = \{a^n b^n \mid n \geq 0\}$ $\Sigma = \{a, b\}$
 $P = \{S \rightarrow aSb \mid \epsilon\}$ $\Gamma = \{S\}$ $S \Rightarrow S$

Closure CFGs: A set is closed under an operation iff application of that operation on members of the set produces more members of that set.

$L = \{a^n b^n c^n \mid n > 0\}$ is not CF Pumping Lemma

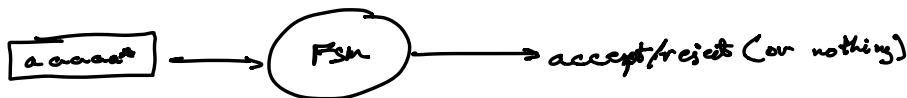
$L_1 \cap L_2$ is not CF $L_1 = a^n b^n c^n$ $L_2 = a^n b^n c^n$

Closed for CFL: $L_1 \cup L_2$ $S \rightarrow S_1 \mid S_2$
 $L_1 \cap L_2$ $S \rightarrow S_1 \cdot S_2$
 L_1^* $S \rightarrow S_1 S_1^*$

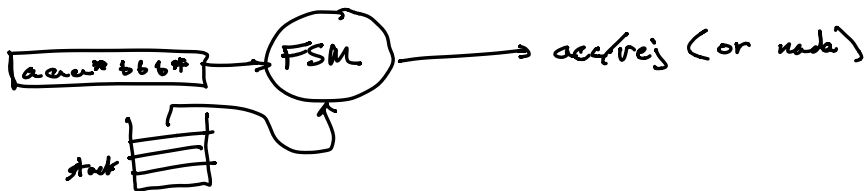
Turing Machine

Radiolab - The Turing Problem
(Jim) completely loopy

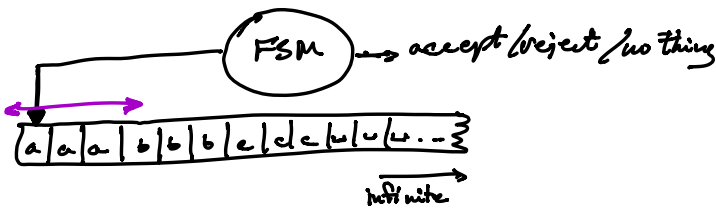
Finite State -
(DFA / NFA)



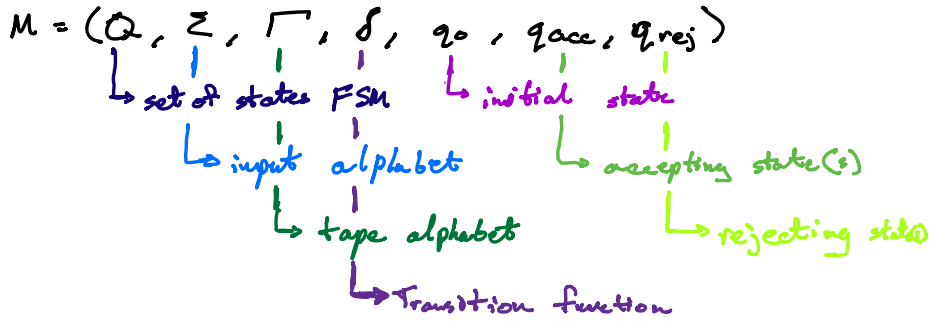
Pushdown Automata



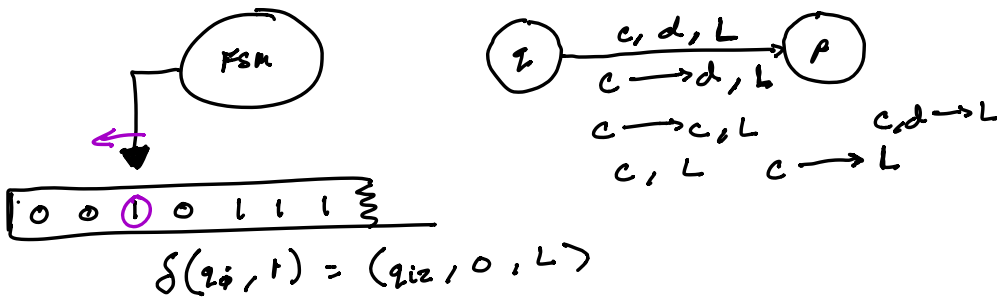
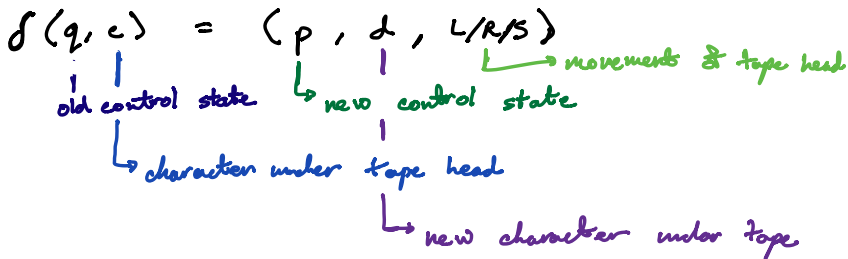
Turing Machine



Formal Definition of TM:



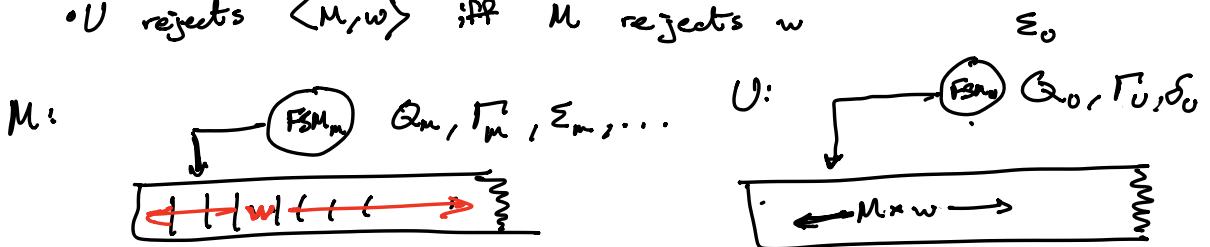
$\delta: Q \times \Gamma \rightarrow Q \times \Gamma \times \{L, R, \emptyset\}$



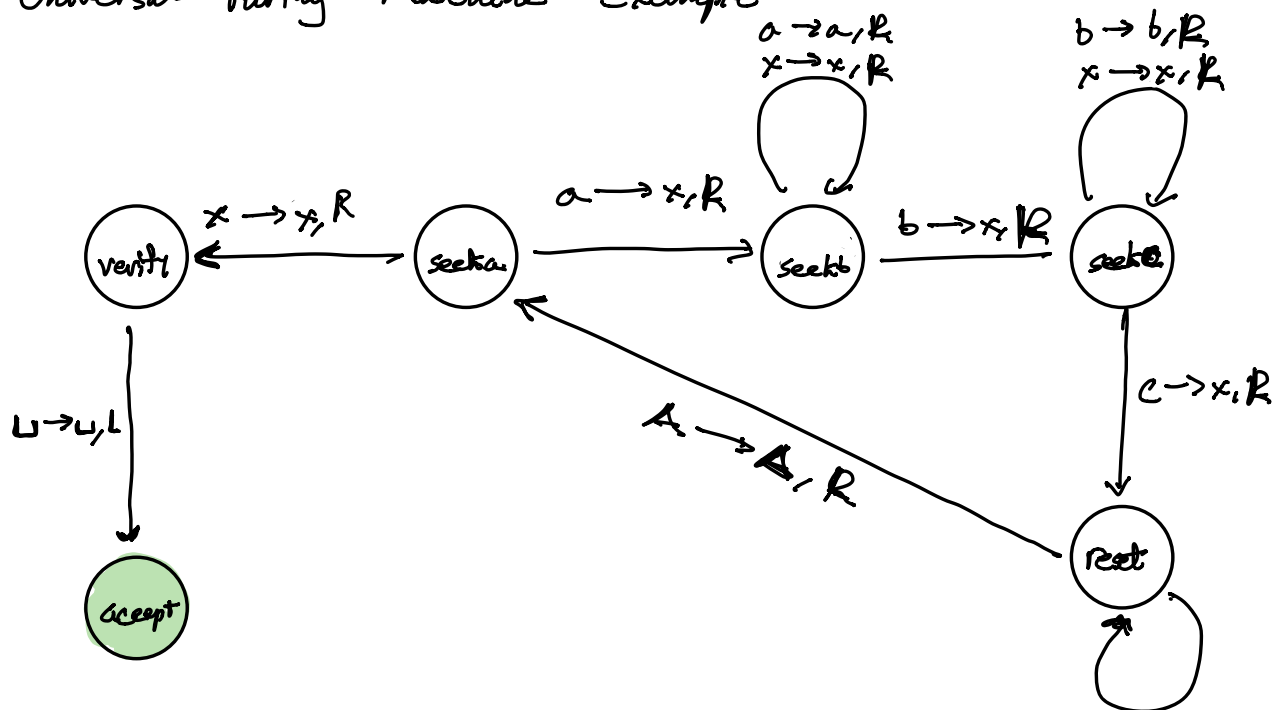
Universal Turing Machine:

Turing Machine that simulates Turing machines

- U accepts $\langle M, w \rangle$ iff M accepts w
- U rejects $\langle M, w \rangle$ iff M rejects w



Universal Turing Machine Example



$Q = \langle \text{start}, \text{seeka}, \text{seekb}, \text{seekc}, \text{verify}, \text{accept} \rangle$

start = 000

seeka = 001 seekb = 001 seekc = 010

$\Gamma = (a, b, c, \sqcup)$

0	0	0	0
0	1	1	1

$\delta(\text{seekc}, c) = (\text{seeka}, b, b)$

[010, 010, 1001, 001, 1]

Input String = aabba and get to

How do we indicate where the head is?

$w = a b c$

