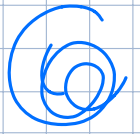


Backtracking \rightarrow Exponential



2nd c. BCE Sanskrit poetry Pingali

0 - short
1 - long

- Fix # syllables
vary short + long

- Fix # beats

short = 1 beat
long = 2 beats

① 0 beats |

① 1 beat | S

② 2 beats | S S
| L

③ 3 beats | S S S
| S L
| L S

⑤ 4 beats | S S S S
| S S L
| S L S
| L S S
| L L

7th century India

$M(n) = \#$ meters n beats long

$$M(n) = \begin{cases} 1 & \text{if } n=0 \\ & \text{or } n=1 \\ M(n-1) + M(n-2) & \text{o/w} \end{cases}$$

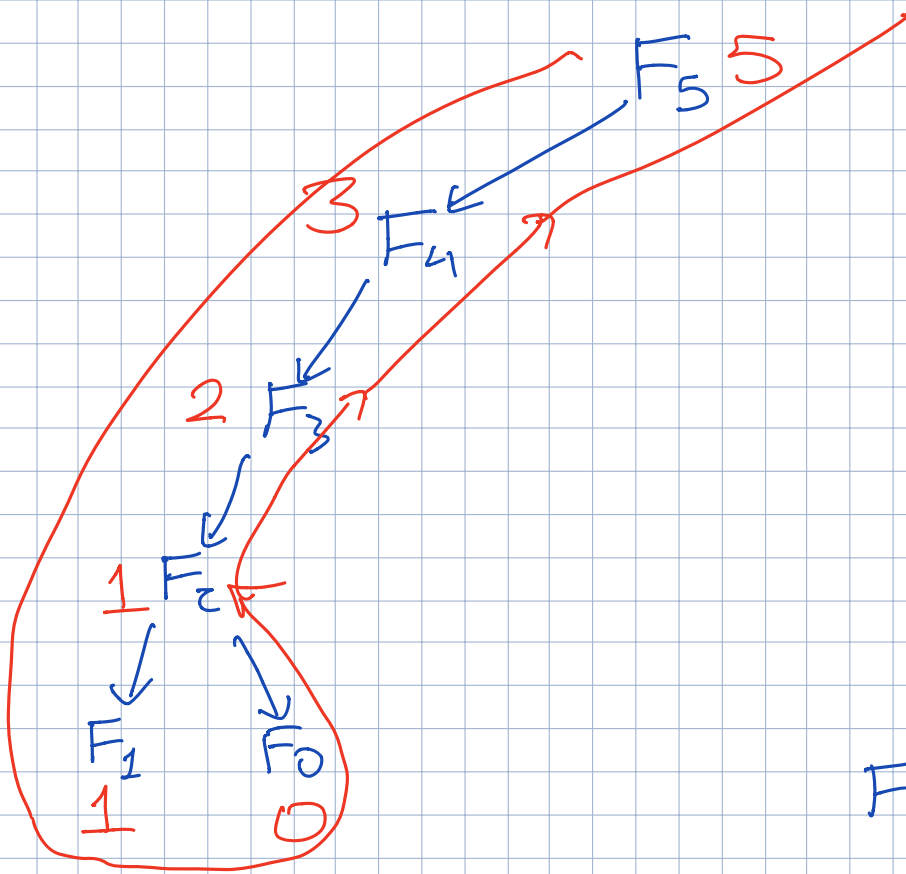
Fibonacci #s

$$F_n = \begin{cases} 0 & \text{if } n=0 \\ 1 & \text{if } n=1 \\ F_{n-1} + F_{n-2} & n > 1 \end{cases} = M_{(n-1)}$$

RecFibo(n):

if $n \leq 1$
return n // assume $n \geq 0$

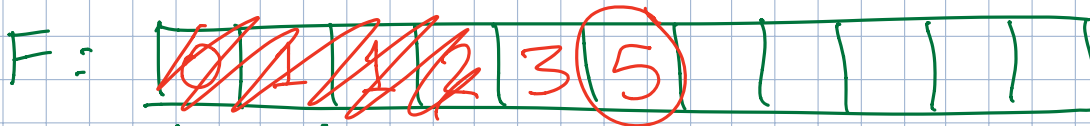
else
return RecFibo($n-1$) + RecFibo($n-2$)



$$F_n = \Theta(\phi^n)$$

$$\phi = \frac{1+\sqrt{5}}{2} \approx 1.618$$

Memoization \rightarrow $O(n)$ time



$F[k]$? \rightarrow

Von Neumann — Two-player games 1940s
Zobrist 1960s

IterFib(n):

```

F(n) =
  F[0] ← 0
  F[1] ← 1
  for i ← 2 to n
    F[i] ← F[i-1] + F[i-2]
  return F[n]

```

$O(n)$ "time"

FastFib(n):

prev ← 1
cur ← 0

for i ← 1 to n

next ← prev + cur

prev ← cur

cur ← next

return cur

$$\begin{bmatrix} 0 & 1 \\ 1 & 1 \end{bmatrix} \begin{bmatrix} x \\ y \end{bmatrix} = \begin{bmatrix} y \\ x+y \end{bmatrix}$$

$$\begin{bmatrix} 0 & 1 \\ 1 & 1 \end{bmatrix}^n \begin{bmatrix} 1 \\ 0 \end{bmatrix} = \begin{bmatrix} F_{n-1} \\ F_n \end{bmatrix}$$

$$\begin{bmatrix} 0 & 1 \\ 1 & 1 \end{bmatrix}^n = \begin{cases} \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix} & \text{if } n=0 \\ \begin{bmatrix} 0 & 1 \\ 1 & 1 \end{bmatrix} & \text{if } n=1 \\ \begin{bmatrix} 0 & 1 \\ 1 & 1 \end{bmatrix}^{\lfloor n/2 \rfloor} \cdot \begin{bmatrix} 0 & 1 \\ 1 & 1 \end{bmatrix}^{\lfloor n/2 \rfloor} \cdot \underbrace{\begin{bmatrix} 0 & 1 \\ 1 & 1 \end{bmatrix}}_{\text{if odd}} & \text{if } n \geq 2 \end{cases}$$

F_n in $O(\log n)$ ~~xxx~~ multiplications
integer

1. Develop recurrence (backtracking)

2. Find memoization structure

→ one entry per subproblem

USE INDICES

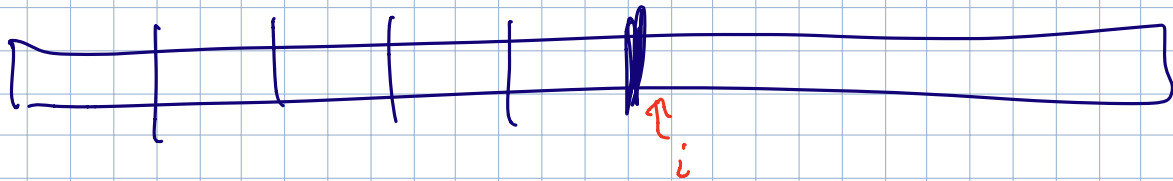
3. Find evaluation order

Text segmentation

Given $A[1..n]$ of letters

Can it be split into English words?

Given $IsWord(i,j) = \text{True}$ iff $A[i..j]$ is an English word



$Splittable(i) = \text{True}$ iff $A[i..n]$ can be split into words

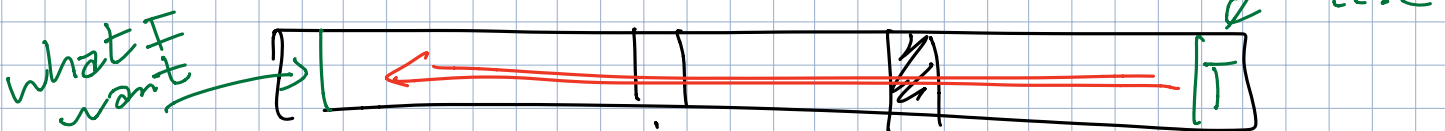
$Splittable(i) = \begin{cases} \text{True} & \text{if } i > n \\ \text{True} & \text{if } IsWord(i,j) \text{ and } Splittable(j+1) \\ & \text{for some } i \leq j \leq n \\ \text{FALSE} & \text{otherwise} \end{cases}$

One param $\Rightarrow O(n)$ space

two params $\Rightarrow O(n^2)$ time

We need $Splittable(1)$.

Array $Splittable(1..n+1)$



Each $S-e[i]$ depends on $S-e[j]$ with $j > i$

So fill in decreasing index order.

$O(n^2)$ time

Splittable?(A[1..n]):

Sp[n+1] ← TRUE

for i ← n down to 1

Sp[i] ← FALSE

for j ← i to n

if IsWord(i, j) and Sp[j+1]

Sp[i] ← TRUE

return Sp[1]