## Programming Languages and Compilers (CS 421)

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Based in part on slides by Mattox Beckman, as updated by Vikram Adve and Gul Agha

## Functions Over Lists

\# let rec double_up list = match list with [ ] -> [ ] (* pattern before ->, expression after *)
| (x :: xs) -> (x :: x :: double_up xs);;
val double_up : 'a list -> 'a list = <fun> \# let fib5_2 = double_up fib5;;
val fib5_2 : int list = [8; $8 ; 5 ; 5 ; 3 ; 3 ; 2 ; 2 ; 1$; 1; 1; 1]

## Functions Over Lists

\# let silly = double_up ["hi"; "there"];;
val silly : string list = ["hi"; "hi"; "there"; "there"]
\# let rec poor_rev list =
match list
with [] -> []
| (x::xs) -> poor_rev xs @ [x];;
val poor_rev : 'a list -> 'a list = <fun>
\# poor_rev silly;;

- : string list = ["there"; "there"; "hi"; "hi"]


## Structural Recursion

- Functions on recursive datatypes (eg lists) tend to be recursive
- Recursion over recursive datatypes generally by structural recursion
- Recursive calls made to components of structure of the same recursive type
- Base cases of recursive types stop the recursion of the function


## Structural Recursion : List Example

\# let rec length list = match list with [ ] -> 0 (* Nil case *)
| x :: xs -> 1 + length xs;; (* Cons case *)
val length : 'a list -> int = <fun>
\# length [5; 4; 3; 2];;

- : int = 4
- Nil case [ ] is base case
- Cons case recurses on component list xs


## Forward Recursion

- In Structural Recursion, split input into components and (eventually) recurse on components
- Forward Recursion form of Structural Recursion
- In forward recursion, first call the function recursively on all recursive components, and then build final result from partial results
- Wait until whole structure has been traversed to start building answer


## Forward Recursion: Examples

\# let rec double_up list = match list with [ ] -> [ ]
| (x :: xs) -> (x :: x :: double_up xs);;
val double_up : 'a list -> 'a list = <fun>
\# let rec poor_rev list =
match list
with [] -> []
| (x::xs) -> let pr = poor_rev xs in pr @ [x] $;$;
val poor_rev : 'a list -> 'a list = <fun>

## Question

- How do you write length with forward recursion?
let rec length I =


## Question

- How do you write length with forward recursion?
let rec length I = match I with [] -> | (a :: bs) ->


## Question

- How do you write length with forward recursion?
let rec length I = match I with [] ->
| (a :: bs) -> length bs


## Question

- How do you write length with forward recursion?
let rec length I =
match I with [] -> 0
| (a :: bs) -> 1 + length bs


## Your turn now

## Try Problem 2 on ML2

## An Important Optimization

- When a function call is made,

Normal call
 the return address needs to be saved to the stack so we know to where to return when the call is finished

- What if $f$ calls $g$ and $g$ calls $h$, but calling $h$ is the last thing $g$ does (a tail cal)?


## An Important Optimization

- When a function call is made, the return address needs to be saved to the stack so we know to where to return when the call is finished
- What if $f$ calls $g$ and $g$ calls $h$, but calling $h$ is the last thing $g$ does (a tail call)?
- Then $h$ can return directly to $f$ instead of $g$


## Tail Recursion

- A recursive program is tail recursive if all recursive calls are tail calls
- Tail recursive programs may be optimized to be implemented as loops, thus removing the function call overhead for the recursive calls
- Tail recursion generally requires extra "accumulator" arguments to pass partial results
- May require an auxiliary function


## Example of Tail Recursion

- Forward recursive:
\# let rec prod I =
match I with [] -> 1
| (x :: rem) -> x * prod rem;;
val prod : int list -> int = <fun>
- Tail recursive:
\# let prod list =
let rec prod_aux I acc =
match I with [] -> acc
| (y :: rest) -> prod_aux rest (acc * y)
(* Uses associativity of multiplication *) in prod_aux list 1;;
val prod : int list -> int = <fun>


## Question

- How do you write length with tail recursion? let length I =


## Question

- How do you write length with tail recursion? let length I =
let rec length_aux list $\mathrm{n}=$
in


## Question

- How do you write length with tail recursion? let length I =
let rec length_aux list $\mathrm{n}=$ match list with [] ->
| (a :: bs) ->
in


## Question

- How do you write length with tail recursion? let length I =
let rec length_aux list $\mathrm{n}=$ match list with [] -> n
| (a :: bs) ->
in


## Question

- How do you write length with tail recursion? let length I =
let rec length_aux list $\mathrm{n}=$ match list with [] -> n
| (a :: bs) -> length_aux
in


## Question

- How do you write length with tail recursion? let length I =
let rec length_aux list $\mathrm{n}=$ match list with [] -> n
| (a :: bs) -> length_aux bs
in


## Question

- How do you write length with tail recursion? let length I =
let rec length_aux list $\mathrm{n}=$ match list with [] -> n
| (a :: bs) -> length_aux bs ( $\mathrm{n}+1$ )
in


## Question

- How do you write length with tail recursion? let length I =
let rec length_aux list $\mathrm{n}=$ match list with [] -> n
| (a :: bs) -> length_aux bs (n + 1) in length_aux I 0


## Your turn now

## Try Problem 4 on MP2

## Mapping Recursion

One common form of structural recursion applies a function to each element in the structure
\# let rec doubleList list = match list with [ ] -> [ ]
| x::xs -> 2 * x :: doubleList xs;;
val doubleList : int list -> int list = <fun>
\# doubleList [2;3;4];;

- : int list = $[4 ; 6 ; 8]$


## Mapping Functions Over Lists

\# let rec map f list =
match list
with [] -> []
| (h::t) -> (f h) :: (map ft);;
val map : ('a -> 'b) -> 'a list -> 'b list = <fun> \# map plus_two fib5;;

- : int list = [10; 7; 5; 4; 3; 3]
\# map (fun x-> x-1) fib6;;
: int list = $12 ; 7 ; 4 ; 2 ; 1 ; 0 ; 0]$


## Mapping Recursion

Can use the higher-order recursive map function instead of direct recursion
\# let doubleList list =
List.map (fun x -> 2 * x) list;;
val doubleList : int list -> int list = <fun>
\# doubleList [2;3;4];;

- : int list = $[4 ; 6 ; 8]$
- Same function, but no rec


## Your turn now

Write a function
make_app : (('a -> 'b) * 'a) list -> 'b list that takes a list of function - input pairs and gives the result of applying each function to its argument. Use map, no explicit recursion.
let make_app I =

## Folding Recursion

- Another common form "folds" an operation over the elements of the structure
\# let rec multList list = match list
with [ ] -> 1
| x::xs -> x * multList xs;;
val multList : int list -> int = <fun>
\# multList [2;4;6];;
- : int = 48
- Computes (2 * (4 * (6 * 1)))


## Folding Functions over Lists

How are the following functions similar?
\# let rec sumlist list = match list with
[ ] -> 0 | x::xs -> x + sumlist xs;;
val sumlist : int list -> int = <fun>
\# sumlist [2;3;4];;

- : int = 9
\# let rec prodlist list = match list with
[ ] -> 1 | x::xs -> x * prodlist xs;;;
val prodlist : int list -> int = <fun>
\# prodlist [2;3;4];;
- : int = 24


## Folding Functions over Lists

How are the following functions similar?
\# let rec sumList list = match list with
[ ] -> 0|| x::xs -> x + sumList xs;;
val sumList : in list -> int = <fun>
\# sumList [2;3;4];;

- : int = 9


## Base Case

\# let rec multlist list $=$ match list with
[ ] -> 1 x: :xs -> x * multList xs;;
val multList : int list -> int $=$ <fun>
\# multList [2;3;4];;

- : int = 24


## Folding Functions over Lists

How are the following functions similar?
\# let rec sumList list = match list with
[ ] -> 0| x::xs -> x + sumList xs;
val sumList : int list $->$ int $=$ <fun>
\# sumList [2;3;4];;

- : int = 9
\# let rec multList list = match list with
[ ] -> 1| x: :xs -> x * multList xs; ;
val multList : int list -> int = <fun>
\# multList [2;3;4];;
- : int = 24


## Folding Functions over Lists

How are the following functions similar?
\# let rec sumList list = match list with
[ ] -> $0 \mid \mathrm{x}:: \mathrm{xs}->\mathrm{X}+$ SumList xs; ;
val sumList : int list -> int $=<$ fun $>$
\# sumList [2;3;4];;

- : int = 9
\# let rec multList list = match list with
[ ] -> 1| x: :xs -> x * multList xs;
val multList : int list -> int = <fun>
\# multList [2;3;4];;
- : int = 24


## Folding Functions over Lists

How are the following functions similar?
\# let rec sumList list = match list with
[ ] -> $0 \mid \mathrm{x}:: \mathrm{xs}->\mathrm{x}+$ SumList xs; ;
val sumList : int list -> int $=<$ fun >
\# sumList [2;3;4];;

- : int = 9
\# let rec multList list $=$ match list with
[ ] -> $1 \mid x:: x s ~->区$ multList $x$; ;
val multList : int list -> int = <fun>
\# multList [2;3;4];;
- : int = 24


## Folding Functions over Lists

How are the following functions similar?
\# let rec sumList list = match list with
[ ] ->0| x::xs -> $\mathrm{x}+$ Rec value
val sumList : int list -> int =<tun>
\# sumList [2;3;4];;

- : int = 9
\# let rec multList list $=$ match list with
[ ] -> $1 \mid x:: x s ~->区 *$ Rec value;
val multList : int list -> int = <fun>
\# multList [2;3;4];;
- : int = 24R


## Recursing over lists

\# let rec fold_right f list $\mathrm{b}=$
match list
with [] -> b
The Primitive
| (x : : xs) -> fx (fold_right f xs b);, Recursion Fairy val fold_right : ('a -> 'b -> 'b) -> 'a list -> 'b -> 'b = <fun>
\# fold_right
(fun s -> fun () -> print_string s)
["hi"; "there"]
();;
therehi- : unit $=()$

## Folding Recursion

- multList folds to the right
- Same as:
\# let multList list =
List.fold_right
(fun x -> fun $\mathrm{p}->\mathrm{x}$ * p )
list 1;;
val multList : int list -> int = <fun>
\# multList [2;4;6];;
- : int = 48


## Encoding Recursion with Fold

\# let rec append list1 list2 = match list1 with
[ ] -> list2 | x::xs -> x :: append xs list2;;
val append : 'a list -> 'a list -> 'a list = <fun>
Base Case
Operation Recursive Call
\# let append list1 list2 =
fold_right (fun x y -> x :: y) list1 list2;;
val append : 'a list -> 'a list -> 'a list = <fun>
\# append $[1 ; 2 ; 3][4 ; 5 ; 6] ;$;

- : int list = [1; 2; 3; 4; 5; 6]


## Question

let rec length I = match I with [] -> 0 | (a :: bs) -> 1 + length bs

- How do you write length with fold_right, but no explicit recursion?


## Question

let rec length I =

## match I with [] -> 0

| (a :: bs) -> 1 + length bs

- How do you write length with fold_right, but no explicit recursion?
let length list =
List.fold_right (fun x -> fun n -> $\mathrm{n}+1$ ) list 0


## Map from Fold

\# let map f list = fold_right (fun $x$-> fun $y->f x:: y$ ) list [];
val map : ('a -> 'b) -> 'a list -> 'b list = <fun>
\# map ((+)1) [1;2;3];;

- : int list = [2; 3; 4]
- Can you write fold_right (or fold_left) with just map? How, or why not?


## Iterating over lists

\# let rec fold_left falist =
match list
with [] -> a
| (x :: xs) -> fold_left f (f a x) xs;;
val fold_left : ('a -> 'b -> 'a) -> 'a -> 'b list -> 'a = <fun>
\# fold_left
(fun () -> print_string)
()
["hi"; "there"];;
hithere- : unit $=()$

## Encoding Tail Recursion with fold_left

\# let prod list = let rec prod_aux I acc = match I with [] -> acc
| (y :: rest) -> prod_aux rest (acc * y)
in prod_aux list.1;;
val prod : int list $->$ int $=$ <fun>
Init Acc Value Recursive Call Operation
\# let prod list =
List.fold_left (fun acc y -> acc * y) 1 list;;
val prod: int list -> int $=$ <fun>
\# prod [4;5;6];;

- : int =120


## Question

let length I =
let rec length_aux list $\mathrm{n}=$ match list with [] -> n | (a :: bs) -> length_aux bs ( $\mathrm{n}+1$ ) in length_aux I 0

- How do you write length with fold_left, but no explicit recursion?


## Question

let length I =
let rec length_aux list $\mathrm{n}=$ match list with [] -> n
| (a :: bs) -> length_aux bs ( $\mathrm{n}+1$ )
in length_aux I 0

- How do you write length with fold_left, but no explicit recursion?
let length list =
List.fold_left (fun n -> fun x -> n + 1) 0 list
9/11/17


## Folding

\# let rec fold_left falist = match list
with [] -> a | (x :: xs) -> fold_left f (f a x) xs;;
val fold_left : ('a -> 'b -> 'a) -> 'a -> 'b list -> 'a = <fun>
fold_left fa $\left[x_{1} ; x_{2} ; \ldots ; x_{n}\right]=f\left(\ldots\left(f\left(f\right.\right.\right.$ a $\left.\left.\left.x_{1}\right) x_{2}\right) \ldots\right) x_{n}$
\# let rec fold_right f list $b=$ match list
with [ ] -> b | (x :: xs) -> fx (fold_right f xs b);;
val fold_right : ('a -> 'b -> 'b) -> 'a list -> 'b -> 'b = <fun>
fold_right $f\left[x_{1} ; x_{2} ; \ldots ; x_{n}\right] b=f x_{1}\left(f x_{2}\left(\ldots\left(f x_{n} b\right) \ldots\right)\right)$

## Recall

\# let rec poor_rev list = match list with [] -> []
| (x::xs) -> poor_rev xs @ [x];;
val poor_rev : 'a list -> 'a list = <fun>

What is its running time?

## Quadratic Time

- Each step of the recursion takes time proportional to input
- Each step of the recursion makes only one recursive call.
- List example:
\# let rec poor_rev list = match list with [] -> []
| (x::xs) -> poor_rev xs @ [x];;
val poor_rev : 'a list -> 'a list = <fun>


## Tail Recursion - Example

\# let rec rev_aux list revlist = match list with [ ] -> revlist
| x :: xs -> rev_aux xs (x::revlist);;
val rev_aux : 'a list -> 'a list -> 'a list = <fun>
\# let rev list = rev_aux list [ ];;
val rev : 'a list -> 'a list = <fun>

- What is its running time?


## Comparison

- poor_rev [1,2,3] =
- (poor_rev [2,3]) @ [1] =
- ((poor_rev [3]) @ [2]) @ [1] =
- (((poor_rev [ ]) @ [3]) @ [2]) @ [1] =
- (([ ] @ [3]) @ [2]) @ [1]) =
- ([3] @ [2]) @ [1] =
- (3:: ([ ] @ [2])) @ [1] =
- [3,2] @ [1] =
- $3::([2]$ @ [1]) =
- 3 :: (2:: ([ ] @ [1])) = [3, 2, 1]


## Comparison

- $\operatorname{rev}[1,2,3]=$
- rev_aux [1,2,3] [ ] =
- rev_aux [2,3] [1] =
- rev_aux [3] [2,1] =
- rev_aux [ ] [3,2,1] = [3,2,1]


## Folding - Tail Recursion

- \# let rev list =
fold_left

> (fun I -> fun x -> x :: I) //comb op
[] //accumulator cell
list

## Folding

- Can replace recursion by fold_right in any forward primitive recursive definition
- Primitive recursive means it only recurses on immediate subcomponents of recursive data structure
- Can replace recursion by fold_left in any tail primitive recursive definition


## Continuation Passing Style

- A programming technique for all forms of "non-local" control flow:
- non-local jumps
- exceptions
- general conversion of non-tail calls to tail calls
- Essentially it's a higher-order function version of GOTO


## Continuations

- Idea: Use functions to represent the control flow of a program
- Method: Each procedure takes a function as an argument to which to pass its result; outer procedure "returns" no result
- Function receiving the result called a continuation
- Continuation acts as "accumulator" for work still to be done


## Example of Tail Recursion

\# let rec app fl $\mathrm{x}=$
match fl with [] -> $x$
| (f :: rem_fs) -> f (app rem_fs x)ii
val app: ('a -> 'a) list ->'a -> 'a = <fun>
\# let app fs x =
let rec app_aux fl acc=
match fl with [] -> acc
| (f :: rem_fs) -> app_aux rem_fs (fun z -> acc (f z))
in app_aux fs (fun y -> y) x;;
val app : ('a -> 'a) list -> 'a -> 'a = <fun>

## Continuation Passing Style

- Writing procedures so that they take a continuation to which to give (pass) the result, and return no result, is called continuation passing style (CPS)


## Example of Tail Recursion \& CSP

\# let app fs $x=$ let rec app_aux flacc= match fl with [] -> acc
| (f :: rem_fs) -> app_aux rem_fs
(fun z -> acc (f z))
in app_aux fs (fun y -> y) x;;
val app : ('a -> 'a) list -> 'a -> 'a = <fun>
\# let rec appk fl x k =
match fl with [] -> kx
| (f :: rem_fs) -> appk rem_fs x (fun z -> k (f z));;
val appk : ('a -> 'a) list -> 'a -> ('a -> 'b) -> 'b

## Continuation Passing Style

- A compilation technique to implement nonlocal control flow, especially useful in interpreters.
- A formalization of non-local control flow in denotational semantics
- Possible intermediate state in compiling functional code


## Terms

- A function is in Direct Style when it returns its result back to the caller.
- A Tail Call occurs when a function returns the result of another function call without any more computations (eg tail recursion)
- A function is in Continuation Passing Style when it passes its result to another function.
- Instead of returning the result to the caller, we pass it forward to another function.


## Continuation Passing Style

- A compilation technique to implement nonlocal control flow, especially useful in interpreters.
- A formalization of non-local control flow in denotational semantics
- Possible intermediate state in compiling functional code


## Example

- Simple reporting continuation:
\# let report x = (print_int x; print_newline( ) );;
val report : int -> unit = <fun>
- Simple function using a continuation:
\# let plusk a b k = k (a + b)
val plusk : int -> int -> (int -> 'a) -> 'a = <fun> \# plusk 2022 report;;
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- : unit = ()


## Simple Functions Taking Continuations

- Given a primitive operation, can convert it to pass its result forward to a continuation
- Examples:
\# let subk x y k = k(x + y) ;;
val subk : int -> int -> (int -> 'a) -> 'a = <fun>
\# let eqk $x$ y $k=k(x=y)$;;
val eqk : 'a -> 'a -> (bool -> 'b) -> 'b = <fun>
\# let timesk x y k = k( $x^{*}$ y) ; ;
val timesk : int -> int -> (int -> 'a) -> 'a = <fun>


## Nesting Continuations

\# let add_three x y z = x + y + z; ;
val add_three : int -> int -> int -> int = <fun> \# let add_three $x y z=$ let $p=x+y$ in $p+z ;$; val add_three : int -> int -> int -> int = <fun> \# let add_three_k x y z k =

val add_three_k : int -> int -> int -> (int -> 'a)
-> 'a = <fun>

