

CS 473: Algorithms

Ruta Mehta

University of Illinois, Urbana-Champaign

Spring 2018

Administrivia, Introduction

Lecture 1

Jan 16, 2018

Some of the slides are courtesy Prof. Chekuri

Part I

Administrivia

Instructional Staff

- ① **Instructors:** Ruta Mehta
- ② **Teaching Assistants:** Shant Boodaghian and Vasilis Livanos
- ③ **Graders:** TBD
- ④ **Office hours:** See course webpage
- ⑤ **Email:** Use private notes on Piazza to reach course staff.

Online resources

- 1 **Webpage:** General information, lecture schedule/slides/notes, homeworks, course policies
courses.engr.illinois.edu/cs473
- 2 **Gradescope:** HW submission, grading, regrade requests
- 3 **Moodle:** HW solutions, grades
- 4 **Piazza:** Announcements, online questions and discussion, contacting course staff (via private notes)

See course webpage for links

Important: check Piazza/course web page at least once each day

Prerequisites

- 1 **Prerequisites:** CS 173 (discrete math), CS 225 (data structures), CS 374 (algorithms and models of computation) or sufficient mathematical maturity
- 2 Concretely:
 - 1 Good ability to write formal proofs of correctness
 - 2 Comfort with recursive thinking/algorithms, reductions
 - 3 Comfort with basic data structures: balanced binary search trees, priority queues, heaps, etc.
 - 4 Basic graph algorithms: reachability (DFS/BFS), undirected vs directed, strong connected components, shortest paths and Dijkstra's algorithm, minimum spanning trees
 - 5 Probability: random variables, expectation, variance
 - 6 Exposure to models of computation and NP-Completeness (optional but will help)

Textbooks and Resources

No one specific textbook for the course.

- 1 **Recommended books:** (not required)
 - 1 Algorithms by Dasgupta, Papadimitriou & Vazirani.
Available online for free!
 - 2 Algorithm Design by Kleinberg & Tardos
- 2 **Lecture notes/slides/pointers:** available on course web-page
- 3 **Additional References**
 - 1 Lecture notes of Jeff Erickson, Timothy Chan, Sarel HarPeled, and others
 - 2 Computers and Intractability: Garey and Johnson.

Grading Policy: Overview

- ① **Homeworks:** 25%
- ② **Midterms:** 45% (**2 × 22.5%**)
- ③ **Finals:** 30% (covers the full course content)

Midterms dates:

- ① Midterm 1: Mon, Feb 26, 7–9pm, 1320 DCL
- ② Midterm 2: Mon, Apr 9, 7–9pm, 1320 DCL
- ③ Final Exam: Fri, May 11

No conflict exam offered unless you have a valid reason (see course webpage).

Homeworks

- ① One homework every week: Due on Wednesday at 8pm. To be submitted electronically in pdf form in *Gradescope*. Assigned at least a week in advance.
- ② Homeworks can be worked on in groups of up to 3 and each group submits *one* written solution (**except Homework 0**).
- ③ **Important:** academic integrity policies. See course web page.

More on Homeworks

- ① No extensions or late homeworks accepted.
- ② To compensate, five problems will be dropped. Homeworks typically have three problems each.
- ③ **Important:** Read homework faq/instructions on website.

Advice

- ① Attend lectures, please ask plenty of questions.
- ② Don't skip homework and don't copy homework solutions.
- ③ Study regularly and keep up with the course.
- ④ This is a course on problem solving. Solve as many as you can!
Books/notes have plenty.
- ⑤ Ask for help promptly. Make use of office hours/Piazza.
- ⑥ This is an optional mixed undergrad/grad course.
(Mathematical) maturity and independence are expected.

Homework 0

- ① HW 0 is posted on the class website.
- ② HW 0 due on Wednesday Jan 24 at 8pm
- ③ HW 0 to be done and submitted individually.

Miscellaneous

Please contact instructors if you need special accommodations.

Lectures are being taped. A link to the videos will be put up on course webpage.

Emergencies: see information at link <http://police.illinois.edu/dpsapp/wp-content/uploads/2017/12/CEOP-2017.pdf>

Part II

Course Goals and Overview

Course Structure

Course divided into four parts:

- 1 Recursion, dynamic programming.
- 2 Randomization in algorithms
- 3 Combinatorial and Discrete Optimization: flows/cuts, matchings, introduction to linear and convex programming
- 4 Intractability and heuristics

Course Goals

Mostly algorithms:

- 1 Some fundamental problems and algorithms
 - 1 FFT, Hashing, Flows/Cuts, Matchings, LP, approximation, ...
- 2 Broadly applicable techniques in algorithm design
 - 1 Recursion, Divide and Conquer, Dynamic Programming
 - 2 Randomization in algorithms and data structures
 - 3 Optimization via convexity and duality
 - 4 Approximation and heuristics
 - 5 Role of mathematics in algorithm design: graph theory, (linear) algebra, geometry, convexity ...

Complexity and Algorithms

- *P*: class of problems that can be solved in polynomial time
- *EXP*: class of problems that can be solved in exponential time
- *NP*: non-deterministic polynomial-time.
- *DECIDABLE*: class of problems that have an algorithm

Complexity and Algorithms

- ***P***: class of problems that can be solved in polynomial time
- ***EXP***: class of problems that can be solved in exponential time
- ***NP***: non-deterministic polynomial-time.
- ***DECIDABLE***: class of problems that have an algorithm

Theorem

There exist (many) undecidable problems.

Complexity and Algorithms

- P : class of problems that can be solved in polynomial time
- EXP : class of problems that can be solved in exponential time
- NP : non-deterministic polynomial-time.
- $DECIDABLE$: class of problems that have an algorithm

Theorem

There exist (many) undecidable problems.

Theorem

P is a strict subset of EXP .

Complexity and Algorithms

- P : class of problems that can be solved in polynomial time
- EXP : class of problems that can be solved in exponential time
- NP : non-deterministic polynomial-time.
- $DECIDABLE$: class of problems that have an algorithm

Theorem

There exist (many) undecidable problems.

Theorem

P is a strict subset of EXP .

$P \subseteq NP \subseteq EXP$. Major open problem: Is $P = NP$?

Complexity and Algorithms

- P : class of problems that can be solved in polynomial time
- EXP : class of problems that can be solved in exponential time
- NP : non-deterministic polynomial-time.
- $DECIDABLE$: class of problems that have an algorithm

Theorem

There exist (many) undecidable problems.

Theorem

P is a strict subset of EXP .

$P \subseteq NP \subseteq EXP$. Major open problem: Is $P = NP$?

Many useful and important problems are *intractable*: $NPComplete$, $EXPCOMplete$, $UNDECIDABLE$.

Complexity and Algorithms

Goals of algorithm design.

- find the “best” possible algorithm for some specific problems of interest
- develop broadly applicable techniques for algorithm design

Goals of complexity:

- prove lower bounds for specific problems
- develop broadly applicable techniques for proving lower bounds
- develop complexity classes to characterize many problems

Rich interplay between the two areas.

Important Ingredients in Algorithm Design

- 1 What is the problem (really)?
 - 1 What is the input? How is it represented?
 - 2 What is the output?
- 2 What is the model of computation? What basic operations are allowed?
- 3 Algorithm design
 - Understand the structure of the problem
 - Relate it to standard and known problems via reductions
 - Try algorithmic paradigms: recursion, divide and conquer, dynamic programming, greedy, convex optimization, ...
 - Failing, try to prove a lower bound via reduction to existing hard problems or settle for approximation, heuristics.
- 4 Proving correctness of algorithm
- 5 Analysis of time and space complexity
- 6 Algorithmic engineering

Recall: Time (Space) Complexity and Notations

Representing running time of an algorithm on an n sized input:

- Upper bound $O(f(n))$: Takes at most $c \cdot f(n)$ time for some constant $c \in \mathbb{R}_+$.
- Lower bound $\Omega(f(n))$: Takes at least $c \cdot f(n)$ time for some constant $c \in \mathbb{R}_+$.
- Tight bound $\Theta(f(n))$: Takes at least $c \cdot f(n)$ and at most $c' \cdot f(n)$ time for some $c, c' \in \mathbb{R}_+$.