CS 525 Advanced Distributed Systems Spring 2018

Indranil Gupta (Indy) Lecture 3 Cloud Computing (Contd.) January 24, 2018 All slides © IG

What is MapReduce?

- Terms are borrowed from Functional Language (e.g., Lisp) Sum of squares:
- (map square '(1 2 3 4))
 - Output: (1 4 9 16)

[processes each record sequentially and independently]

- (reduce + '(1 4 9 16))
 - (+16(+9(+41)))
 - Output: 30

[processes set of all records in batches]

- Let's consider a sample application: Wordcount
 - You are given a <u>huge</u> dataset (e.g., Wikipedia dump or all of Shakespeare's works) and asked to list the count for each of the words in each of the documents therein

Map

• Process individual records to generate intermediate key/value pairs.



Input <filename, file text>

3

Map

• Parallelly Process individual records to generate intermediate key/value pairs.



Map

- Parallelly Process a large number of individual records to generate intermediate
 - key/value pairs.



Reduce

• Reduce processes and merges all intermediate values associated <u>per key</u>



Reduce

- Each key assigned to one Reduce
- Parallelly Processes and merges all intermediate values <u>by partitioning</u> <u>keys</u>



• Popular: *Hash partitioning, i.e.,* key is assigned to reduce # = hash(key)%number of reduce servers

Hadoop Code - Map

```
public static class MapClass extends MapReduceBase
                                                      implements
Mapper<LongWritable, Text, Text, IntWritable> {
 private final static IntWritable one =
    new IntWritable(1);
 private Text word = new Text();
 public void map(LongWritable key, Text value, /* value is a line */
     OutputCollector<Text, IntWritable> output, Reporter reporter)
    throws IOException {
    String line = value.toString();
    StringTokenizer itr = new StringTokenizer(line);
    while (itr.hasMoreTokens()) {
    word.set(itr.nextToken());
    output.collect(word, one);
```

Hadoop Code - Reduce

```
public static class ReduceClass extends MapReduceBase implements
Reducer<Text, IntWritable, Text, IntWritable> {
    public void reduce( /* called once per key */
        Text key,
        Iterator<IntWritable> values,
        OutputCollector<Text, IntWritable> output,
        Reporter reporter)
        throws IOException {
        int sum = 0;
        while (values.hasNext()) {
            sum += values.next().get();
        }
    }
}
```

output.collect(key, new IntWritable(sum));

}

Hadoop Code - Driver

// Tells Hadoop how to run your Map-Reduce job
public void **run** (String inputPath, String outputPath)
 throws Exception {

// The job. WordCount contains MapClass and Reduce.

JobConf conf = new JobConf(WordCount.class);

conf.setJobName("mywordcount");

// The keys are words

(strings) conf.setOutputKeyClass(Text.class);

// The values are counts (ints)

conf.setOutputValueClass(IntWritable.class);

conf.setMapperClass(MapClass.class);

conf.setReducerClass(ReduceClass.class);

FileInputFormat.addInputPath(

conf, newPath(inputPath));

FileOutputFormat.setOutputPath(

```
conf, new Path(outputPath));
```

```
JobClient.runJob(conf);
```

Some Applications of MapReduce

Distributed Grep:

- Input: large set of files
- Output: lines that match pattern
- Map Emits a line if it matches the supplied pattern
- Reduce Copies the intermediate data to output

Some Applications of MapReduce (2) Reverse Web-Link Graph

- Input: Web graph: tuples (a, b) where (page a \rightarrow page b)
- Output: For each page, list of pages that link to it

- Map process web log and for each input <source, target>, it outputs <target, source>
- Reduce emits <target, list(source)>

Some Applications of MapReduce

Count of URL access frequency

- Input: Log of accessed URLs, e.g., from proxy server
- Output: For each URL, % of total accesses for that URL
- Map Process web log and outputs <URL, 1>
- Multiple Reducers *Emits* < *URL*, *URL_count*>
- (So far, like Wordcount. But still need %)
- Chain another MapReduce job after above one
- Map Processes < URL, URL_count> and outputs <1, (< URL, URL_count>)>
- 1 Reducer Does two passes over input. First sums up URL_count's to calculate overall_count.

Emits multiple <URL, URL_count/overall_count>

Some Applications of MapReduce

Map task's output is sorted (e.g., quicksort) Reduce task's input is sorted (e.g., mergesort)

Sort

- Input: Series of (key, value) pairs
- Output: Sorted <value>s
- $Map \langle key, value \rangle \rightarrow \langle value, _ \rangle$ (identity)
- Reducer <key, value> \rightarrow <key, value> (identity)
- Partitioning function partition keys across reducers based on ranges (can't use hashing!)
 - Take data distribution into account to balance reducer tasks

Programming MapReduce

Externally: For user

- 1. Write a Map program (short), write a Reduce program (short)
- 2. Specify number of Maps and Reduces (parallelism level)
- 3. Submit job; wait for result
- 4. Need to know very little about parallel/distributed programming!

Internally: For the Paradigm and Scheduler

- 1. Parallelize Map
- 2. Transfer data from Map to Reduce
- 3. Parallelize Reduce

4. Implement Storage for Map input, Map output, Reduce input, and Reduce output
(Ensure that no Reduce starts before all Maps are finished. That is, ensure the *barrier* between the Map phase and Reduce phase)

Inside MapReduce

For the cloud:

- 1. Parallelize Map: easy! each map task is independent of the other!
 - All Map output records with same key assigned to same Reduce
- 2. Transfer data from Map to Reduce ("Shuffle" phase):
 - All Map output records with same key assigned to same Reduce task
 - use partitioning function, e.g., hash(key)%number of reducers
- 3. Parallelize Reduce: easy! each reduce task is independent of the other!
- 4. Implement Storage for Map input, Map output, Reduce input, and Reduce output
 - Map input: from distributed file system
 - Map output: to local disk (at Map node); uses local file system
 - Reduce input: from (multiple) remote disks; uses local file systems
 - Reduce output: to distributed file system

local file system = Linux FS, etc.

distributed file system = GFS (Google File System), HDFS (Hadoop Distributed File System)



Resource Manager (assigns maps and reduces to servers)

The YARN Scheduler

- Used in Hadoop 2.x +
- YARN = Yet Another Resource Negotiator
- Treats each server as a collection of *containers*
 - Container = fixed CPU + fixed memory
- Has 3 main components
 - Global Resource Manager (RM)
 - Scheduling
 - Per-server Node Manager (NM)
 - Daemon and server-specific functions
 - Per-application (job) Application Master (AM)
 - Container negotiation with RM and NMs
 - Detecting task failures of that job

YARN: How a job gets a container



Fault Tolerance

- Server Failure
 - NM heartbeats to RM
 - If server fails, RM lets all affected AMs know, and AMs take action
 - NM keeps track of each task running at its server
 - If task fails while in-progress, mark the task as idle and restart it
 - AM heartbeats to RM
 - On failure, RM restarts AM, which then syncs up with its running tasks
- RM Failure
 - Use old checkpoints and bring up secondary RM
- Heartbeats also used to piggyback container requests
 - Avoids extra messages

Slow Servers

Slow tasks are called **Stragglers**

- •The slowest task slows the entire job down (why?)
- •Due to Bad Disk, Network Bandwidth, CPU, or Memory
- •Keep track of "progress" of each task (% done)
- •Perform proactive backup (replicated) execution of straggler task: task considered done when first replica complete. Called Speculative Execution.

Locality

- Locality
 - Since cloud has hierarchical topology (e.g., racks)
 - GFS/HDFS stores 3 replicas of each of chunks (e.g., 64 MB in size)
 - Maybe on different racks, e.g., 2 on a rack, 1 on a different rack
 - Mapreduce attempts to schedule a map task on
 - a machine that contains a replica of corresponding input data, or failing that,
 - on the same rack as a machine containing the input, or failing that,
 - Anywhere

Mapreduce: Summary

• Mapreduce uses parallelization + aggregation to schedule applications across clusters

• Need to deal with failure

• Plenty of ongoing research work in scheduling and fault-tolerance for Mapreduce and Hadoop

10 Challenges [Above the Clouds]

(Index: Performance Data-related Scalability Logistical)

- Availability of Service: Use Multiple Cloud Providers; Use Elasticity; Prevent DDOS
- Data Lock-In: Standardize APIs; Enable Surge Computing
- Data Confidentiality and Auditability: Deploy Encryption, VLANs, Firewalls, Geographical Data Storage
- Data Transfer Bottlenecks: Data Backup/Archival; Higher BW Switches; New Cloud Topologies; FedExing Disks
- Performance Unpredictability: QoS; Improved VM Support; Flash Memory; Schedule VMs
- Scalable Storage: Invent Scalable Store
- Bugs in Large Distributed Systems: Invent Debuggers; Real-time debugging; predictable prerun-time debugging
- Scaling Quickly: Invent Good Auto-Scalers; Snapshots for Conservation
- Reputation Fate Sharing
- Software Licensing: Pay-for-use licenses; Bulk use sales

A more Bottom-Up View of Open Myriad interesting problems that acknowledge the characteristics that make today's cloud computing unique: massive scale + on-demand + data-intensive + new programmability + and infrastructure- and application-

- specific details.
- Monitoring: of systems&applications; single site and multi-site
- Storage: massive scale; global storage; for specific apps or classes
- Failures: what is their effect, what is their frequency, how do we achieve fault-tolerance?
- Scheduling: Moving tasks to data, dealing with federation
- Communication bottleneck: within applications, within a site
- Locality: within clouds, or across them
- Cloud Topologies: non-hierarchical, other hierarchical
- Security: of data, of users, of applications, confidentiality, integrity
- Availability of Data
- Seamless Scalability: of applications, of clouds, of data, of everything
- Geo-distributed clouds: Inter-cloud/multi-cloud computations
- Second Generation of Other Programming Models? Beyond MapReduce! Storm, GraphLab, Hama
- Pricing Models, SLAs, Fairness
- Green cloud computing
- Stream processing

Example: Rapid Atmospheric Modeling System, ColoState U

- Hurricane Georges, 17 days in Sept 1998
 - "RAMS modeled the mesoscale convective complex that dropped so much rain, in good agreement with recorded data"
 - Used 5 km spacing instead of the usual 10 km
 - Ran on 256+ processors
- Computation-intenstive computing (or HPC = High Performance Computing)
- Can one run such a program without access to a supercomputer?

Distributed Computing Resources



An Application Coded by a Physicist/Biologist/Meterologist



An Application Coded by a Physicist/Biologist/Meterologist



Next: Scheduling Problem



2-level Scheduling Infrastructure



Intra-site Protocol



Internal Allocation & Scheduling Monitoring Distribution and Publishing of Files

Condor (now HTCondor)

- High-throughput computing system from U. Wisconsin Madison
- Belongs to a class of Cycle-scavenging systems

Such systems

- Run on a lot of workstations
- When workstation is free, ask site's central server (or Globus) for tasks
- If user hits a keystroke or mouse click, stop task
 - Either kill task or ask server to reschedule task
- Can also run on dedicated machines

Inter-site Protocol



Globus

- Globus Alliance involves universities, national US research labs, and some companies
- Standardized several things, especially software tools
- Separately, but related: Open Grid Forum
- Globus Alliance has developed the Globus Toolkit

http://toolkit.globus.org/toolkit/

Globus Toolkit

- Open-source
- Consists of several components
 - GridFTP: Wide-area transfer of bulk data
 - GRAM5 (Grid Resource Allocation Manager): submit, locate, cancel, and manage jobs
 - Not a scheduler
 - Globus communicates with the schedulers in intra-site protocols like HTCondor or Portable Batch System (PBS)
 - RLS (Replica Location Service): Naming service that translates from a file/dir name to a target location (or another file/dir name)
 - Libraries like XIO to provide a standard API for all Grid IO functionalities
 - Grid Security Infrastructure (GSI)

Security Issues

- Important in Grids because they are *federated*, i.e., no single entity controls the entire infrastructure
- Single sign-on: collective job set should require once-only user authentication
- Mapping to local security mechanisms: some sites use Kerberos, others using Unix
- Delegation: credentials to access resources inherited by subcomputations, e.g., job 0 to job 1
- Community authorization: e.g., third-party authentication
- These are also important in clouds, but less so because clouds are typically run under a central control
- In clouds the focus is on failures, scale, on-demand nature

Points to Ponder

- Cloud computing vs. Grid computing: what are the differences?
- What has happened to the Grid Computing Community?
 - See Open Cloud Consortium
 - See CCGrid conference
 - See Globus



Summary

- Grid computing focuses on computation-intensive computing (HPC)
- Though often federated, architecture and key concepts have a lot in common with that of clouds
- Are Grids/HPC converging towards clouds?
 - E.g., Compare OpenStack and Globus

Projects:

Where to get your ideas from

- Read through papers. Read ahead! Read both main and optional papers.
- Leverage area overlaps: x was done for problem area 1, but not for problem area 2
- Look at hot areas:
 - Stream processing
 - Machine Learning
 - IoT
 - RDMA
- Look at the JIRAs of these projects
 - Lots of issues listed but not being worked on

Announcements

- Please sign up for a presentation/scriber slot by next Wednesday office hours
- Next up: Peer to peer systems