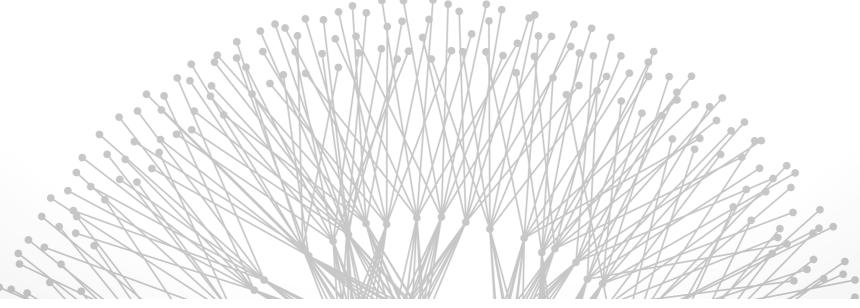
Secure Routing

Brighten Godfrey CS 538 March 28 2018



Network security



Where was security in the design of the original Internet protocols?

- Virtually nowhere!
- All the core protocols (IP,TCP, DNS, BGP) have trivial, glaring vulnerabilities

When security really matters, rely on end-to-end mechanisms

- Public key cryptography & certificate authorities
- But this doesn't protect from all kinds of attacks

With e2e security, what can an attack on BGP still do?

Attacks on Internet routing



Denial of service

- announce "more attractive" path (what does that mean?)
- More likely to be selected: more-specific prefix; shorter path; "cheaper" path

Eavesdropping

- like DoS, a kind of traffic attraction
- but somehow get data to destination or impersonate it

Evasion of accountability

steal or squat on a prefix; send spam; disappear!

How (much) do secure variants of BGP help?

Three approaches to BGP security



- I. Defensive filtering
- 2. Origin Authentication
- 3. Secure BGP (S-BGP)

Many others not discussed here

- Active area of research over the last decade
- Many tradeoffs, especially in deployment issues



Most commonly used class of techniques

Typical implementation

- Filter routes received from customers/peers
- Requires assumptions about what they should be advertising
- Imperfect, requires human maintenance

Filtering of Route Announcements from Peers

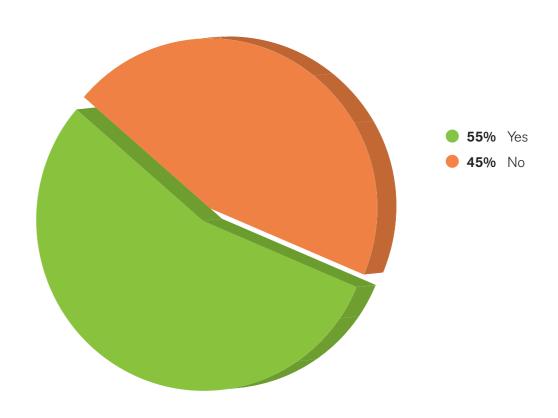


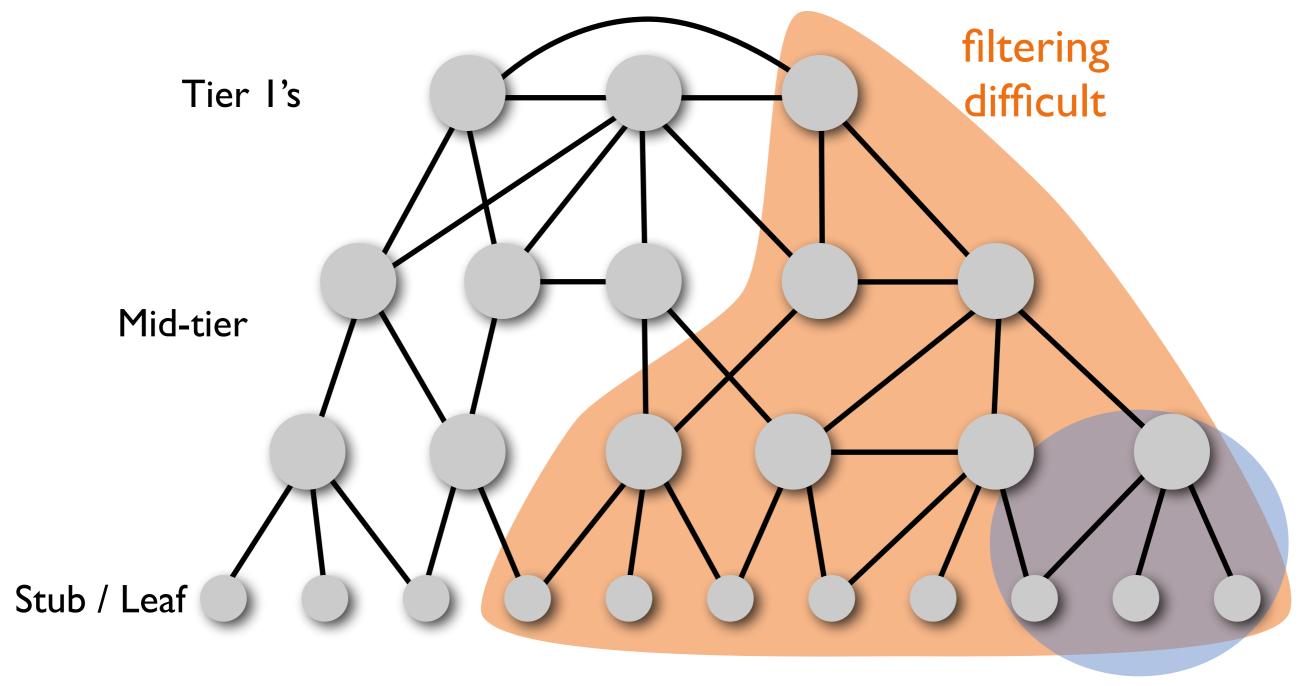
Figure 93 Source: Arbor Networks, Inc.

Arbor Networks survey 2012: 76% filter from customers 55% filter from peers 57% monitor for hijacks

Monito

Source: Arb





filtering feasible



Pretty Good BGP: Improving BGP

by Cautiously Adopting Routes

Josh Karlin, Stephanie Forrest, Jennifer Rexford

ICNP 2006

Pretty Good BGP

- Deprioritize "novel" routes for a period (e.g. 24 hours)
- Routers prefer older (known) routes
- May still pick new route if it's the only option
- Why does this help?

Advantages

- Raises the bar for attacker: route must persist
- Gives time for response
- No protocol changes for deployment

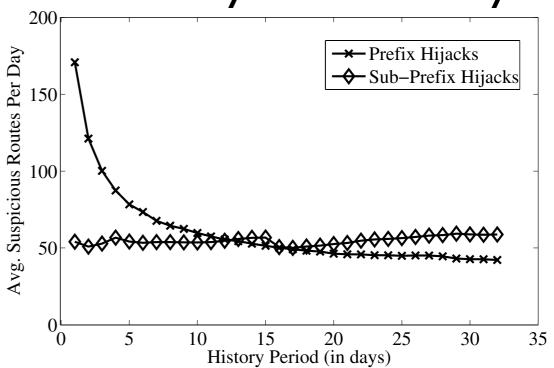
Disadvantages?



Pretty Good BGP

Take-away points

- Prioritization is important: not just good vs. bad route
- Think about human-level solutions
 - # suspicious advertisements is only about 50/day
 - vs. O(400k/day) total



Pretty Good BGP: Improving BGP

by Cautiously Adopting Routes

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ICNP 2006

Fig. 1. Average number of announcements (per day) classified as suspicious using a suspicious period of 1 day and a variety of history periods (h).

2. Origin Authentication



Idea

- Use a Routing Public Key Infrastructure (RPKI) to certify
 AS number assignment and IP address allocation
- An AS can only claim to originate a prefix it owns
- Analogous to PKI for web TLS/SSL security

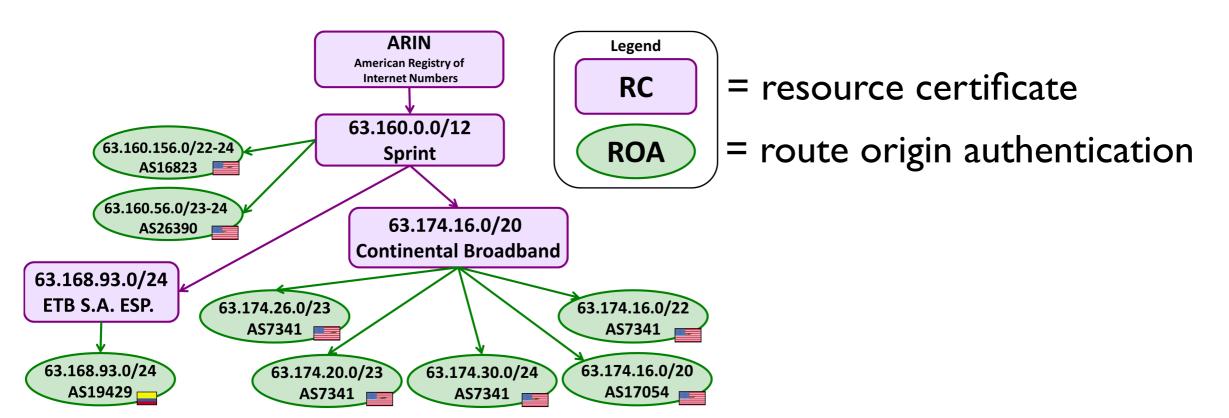


Figure 2: Excerpt of a model RPKI

2. Origin Authentication



Deployment challenges

- Needs router changes to authenticate, filter
- Needs PKI...

Status

- RPKI standardized in 2012
- Now seeing some limited regional deployment

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Internet Engineering Task Force (IETF)

Internet Engineering Task Force (IETF)

Request for Comments: 6480

Category: Informational

ISSN: 2070-1721

An Infrastructure to Support Secure Internet Routing

Abstract

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Secure Origin BGP



Scheme

- Origin Authentication + trusted database of AS-level topology
- Announced routes checked against database to see if they are "plausible" (exist in the topology)

Disadvantages

- Requires trusted database
- Route may be plausible without actually having been announced

3. S-BGP



Scheme

Origin Authentication
 + hop-by-hop
 cryptographic
 validation that path
 was announced

Deployment challenges

- Requires PKI
- Requires significant computational resources

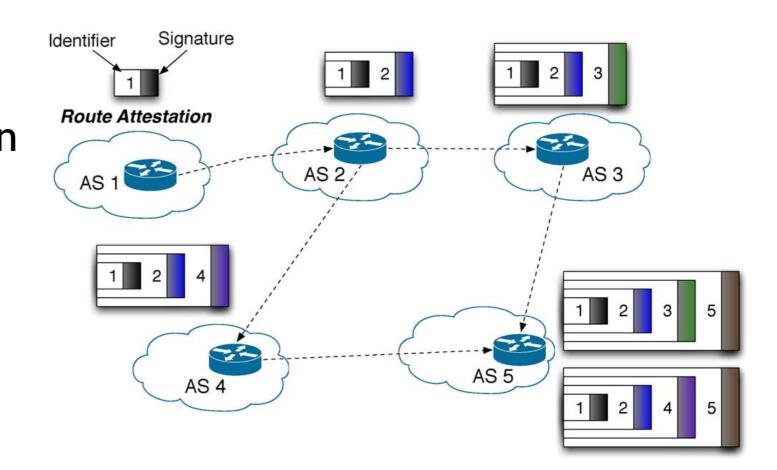


Fig. 5. Route attestations in S-BGP. As UPDATE messages are passed between peers, the receiving peer signs the received message before passing it to another neighbor. The result is an "onion-style" attestation that contains signatures from all routers along the path.

How well do they work?



How Secure are Secure

Interdomain Routing Protocols?

Goldberg, Schapira, Hummon, and

SIGCOMM 2010

Quantifying the attack

- Attacker's goal: attract traffic
- Measure fraction of ASes attacker can "steal" traffic from

How does the attacker do that?

- Basic "smart" strategy
 - Select or invent shortest route you can get away with
 - Advertise it to everyone
- Weird fact: this is not actually the attacker's best strategy; that's NP-complete to compute!

Results



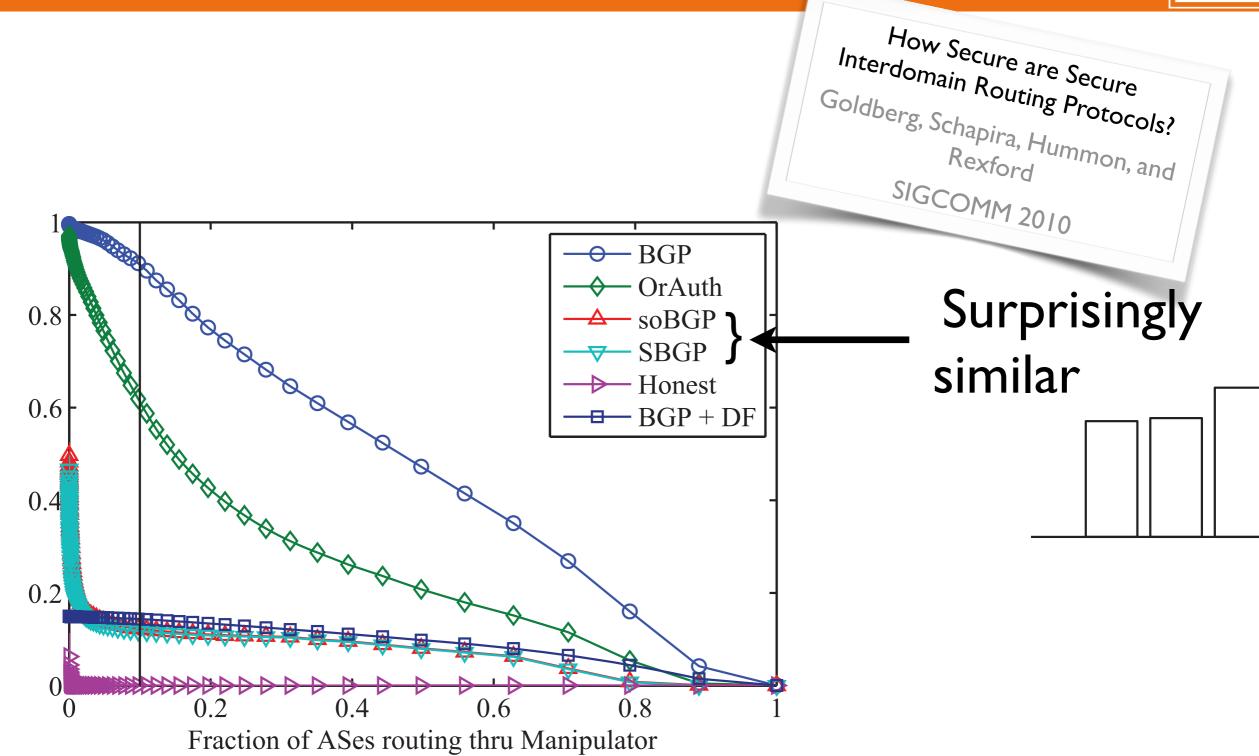


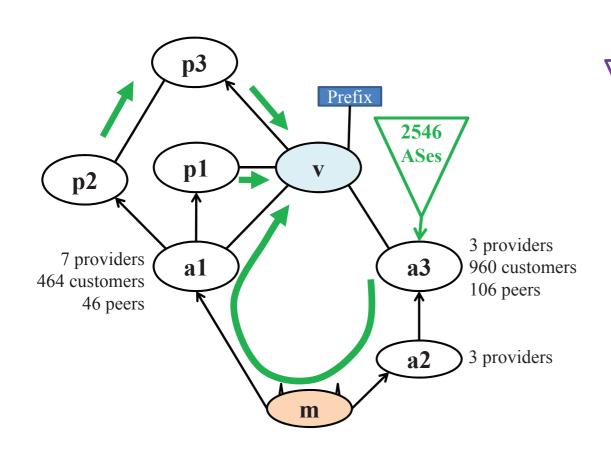
Figure 3: CCDF for the "Shortest-Path Export-All" attack strategy.

soBGP vs. S-BGP



Two components to success

- What you announce soB
- Who you announce it to –



p3 3236 **ASes ASes** Prefix **p2** 3 providers 7 providers 960 customers 464 customers 106 peers 46 peers 3 providers 1682 [Figures from peer & customer Goldberg et al.]

Announce 3-AS path, intercept 5,569 ASes

Announce 4-AS path, intercept 18,664 ASes

Results



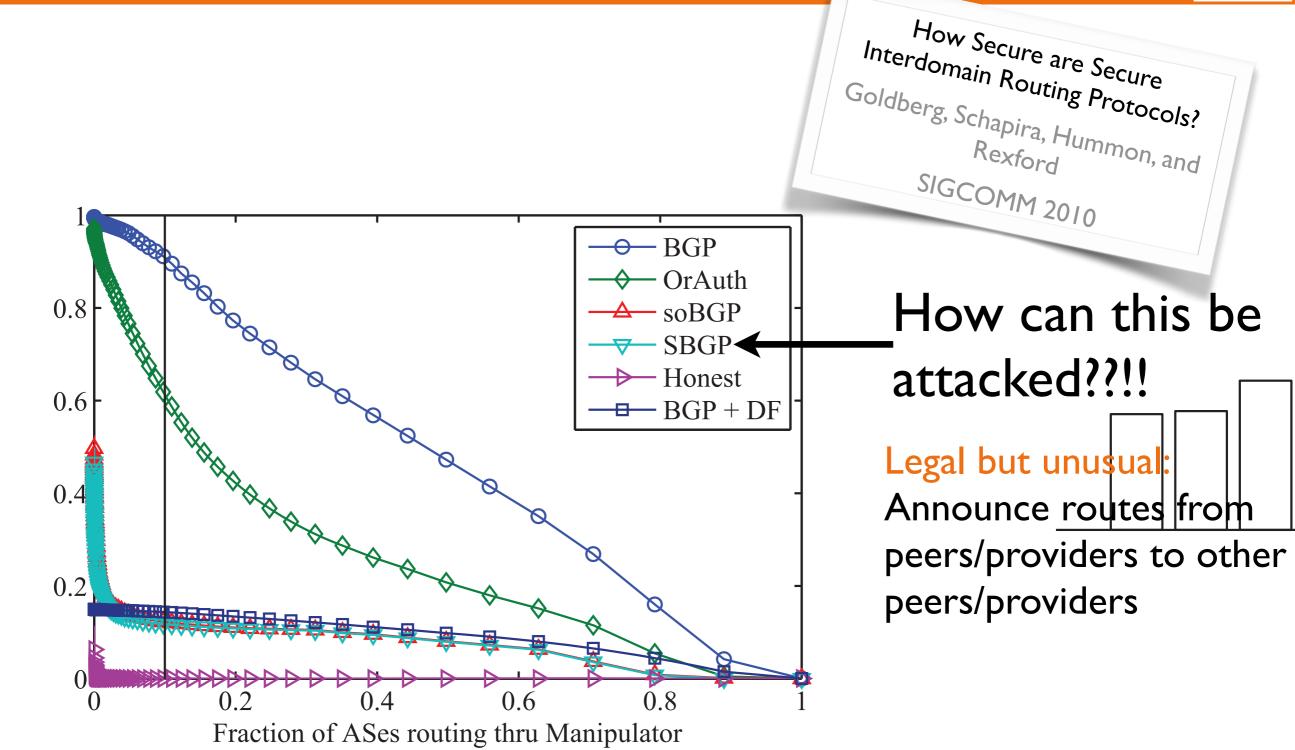


Figure 3: CCDF for the "Shortest-Path Export-All" attack strategy.

Discussion



Is the attack on S-BGP really an attack?

- No, not technically in the protocol
 - ASes are allowed to export whatever routes they like
- Yes, effectively
 - Key point I: unusual export can grab nearly as much traffic as prefix hijack!
 - Key point 2: Want protection against accidents well as attackers

Not just malicious attackers



Many or most high-profile outages and hijacks are likely just configuration errors

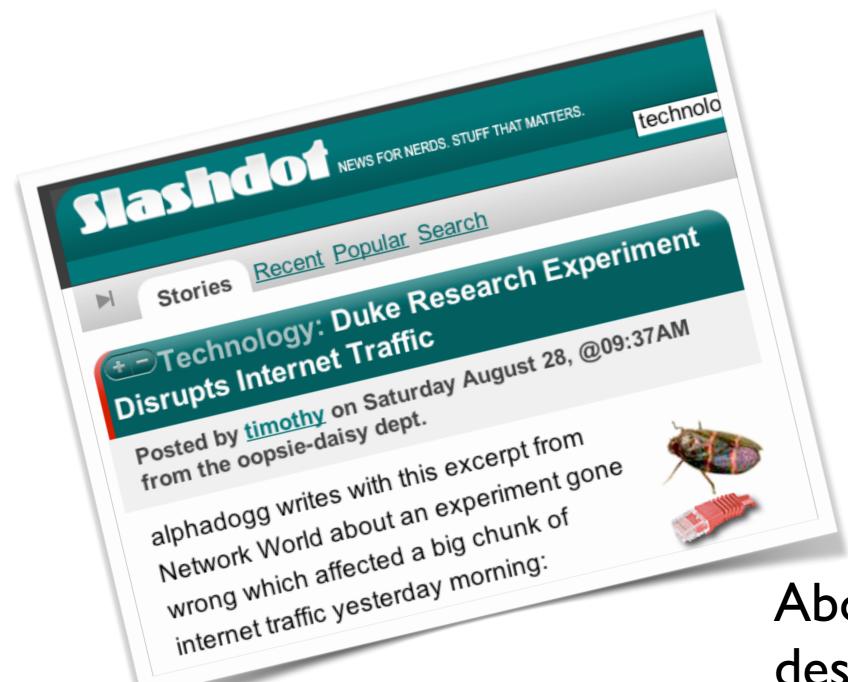
Natural correspondence between attackers and bugs

- behavior unknown ahead of time
- defense is to limit and contain worst-case effects

What about a bug in the protocol?

- worst-case scenario: zero-day exploit on large fraction of routers across the entire Internet
- many are running the same software!

A (bad) day in the life of the Internet

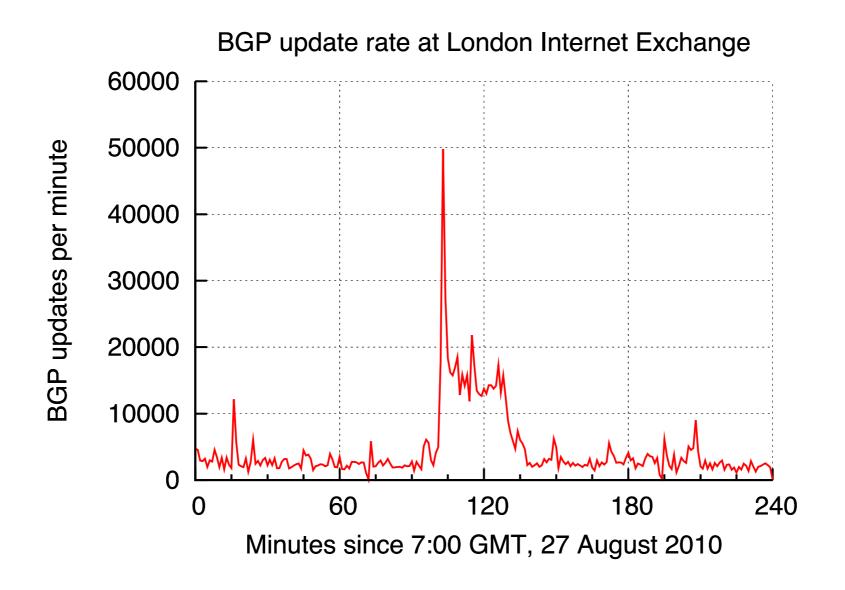


About 1% of Internet destinations disrupted for about 30 minutes

How did this happen?

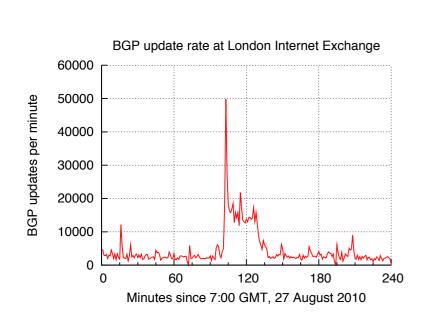
Internet had a bad Friday

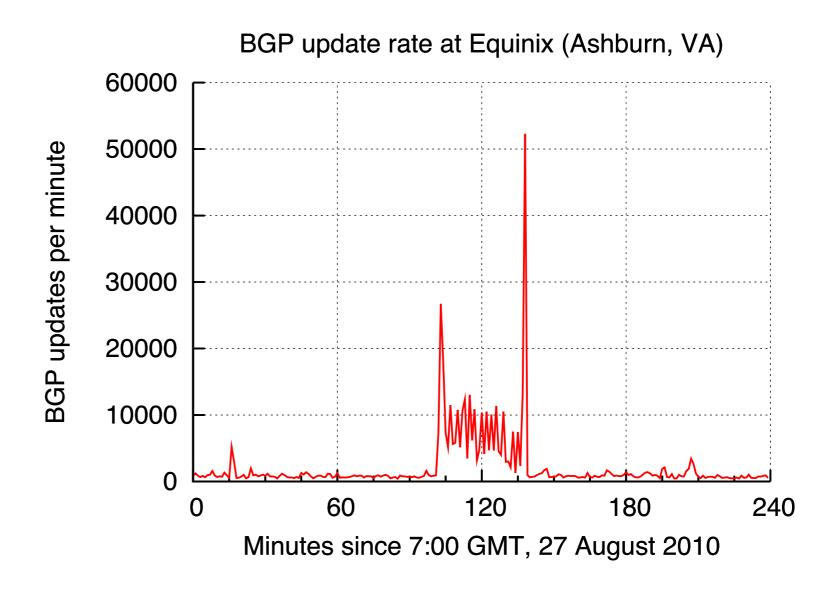




Internet had a bad Friday



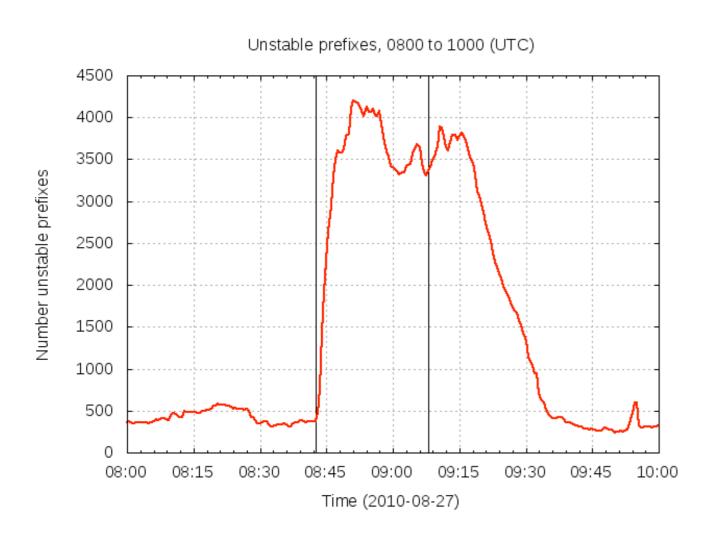


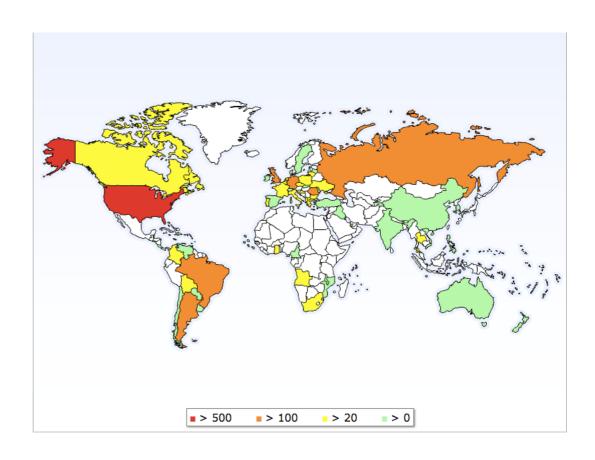


~1% of prefixes affected



[Earl Zmijewski, Renesys]





Brewing a storm



- An unusual announcement
- 2. Propagation from router to router
- 3. Buggy software mangles announcement
- 4.while(true)
 - I. Buggy router propagates announcement to neighbor
 - 2. BGP session dropped upon receipt of mangled message
 - 3. BGP session reestablished

Lessons



Many unsavory BGP announcements can be contained, but this one wasn't

- Spread geographically because it was an entirely valid announcement
- Spread to many prefixes because BGP spec lets one bad announcement from a router affect all traffic to that router

Widespread correlated failures from similar software

Bugs and attacks can have similar effects and solutions

 Lucky in this case: bug triggered by researchers, not attackers!

Discussion



Partial deployment crucial. Issues?

Given all this, why does the Internet work so well?

Announcements



Next time: Data Center Network Architecture