# Data Center Networks

#### Brighten Godfrey CS 538 April 4 2018

Thanks to Ankit Singla for some slides in this lecture

Introduction: The Driving Trends

#### Computing as a utility

- Purchase however much you need, whenever you need it
- Service ranges from access to raw (virtual) machines, to higher level: distributed storage, web services

#### Implications

- Reduces barrier to entry to building large service
  - No need for up-front capital investment
  - No need to plan ahead
- May reduce cost
- Compute and storage becomes more centralized

### The physical cloud: Data centers



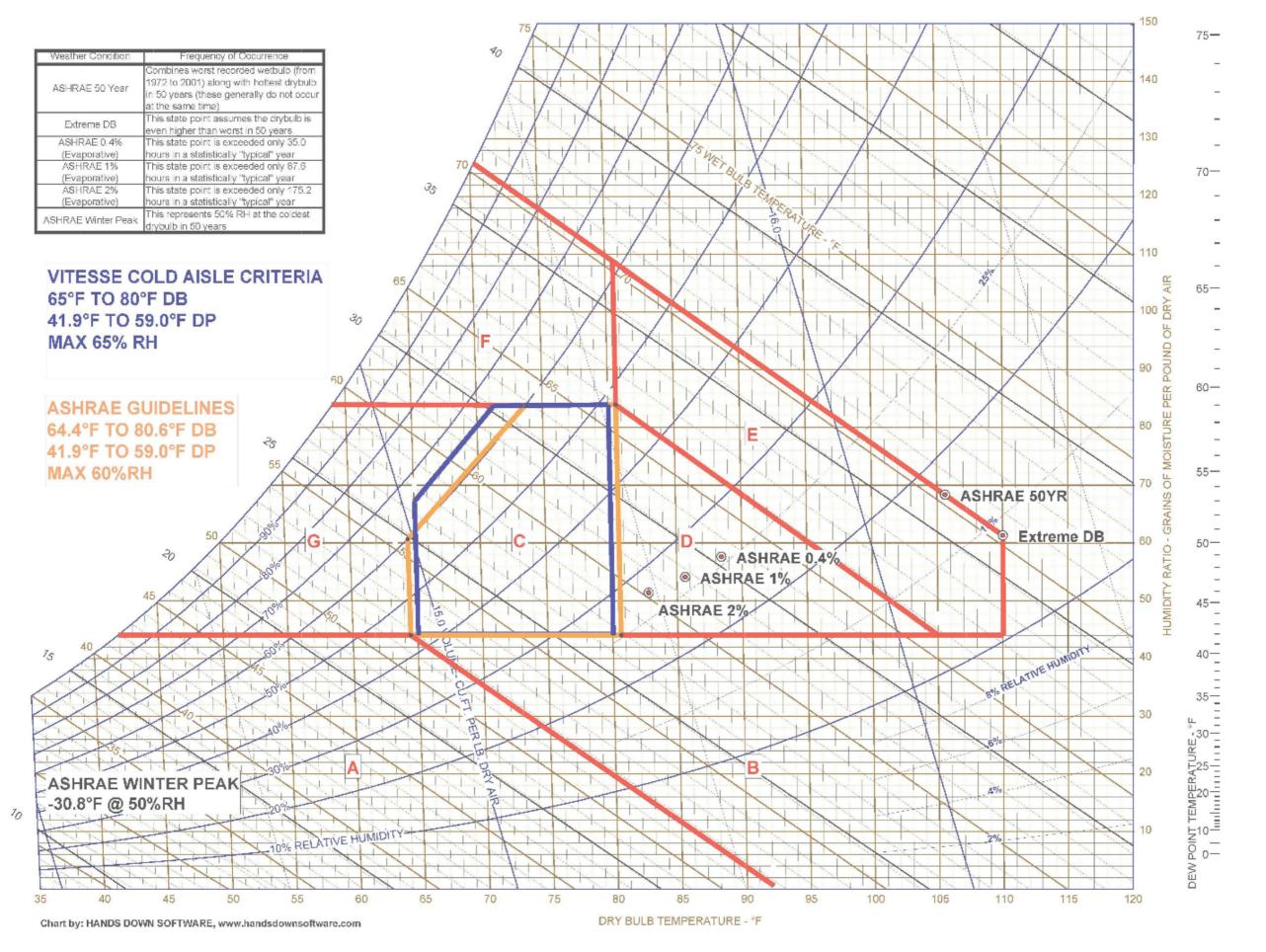
Facebook data center, North Carolina

#### National Petascale Computing Facility, UIUC









[Data Center vI.0, Open Compute Project

One technician for each 15,000 servers [Facebook]

Facility / power infrastructure operated in bulk

- Power usage efficiency (PuE) ~ 1.8 in average DCs
- Pushed down to ~ I.I in large cloud DCs

Ability to custom-design equipment

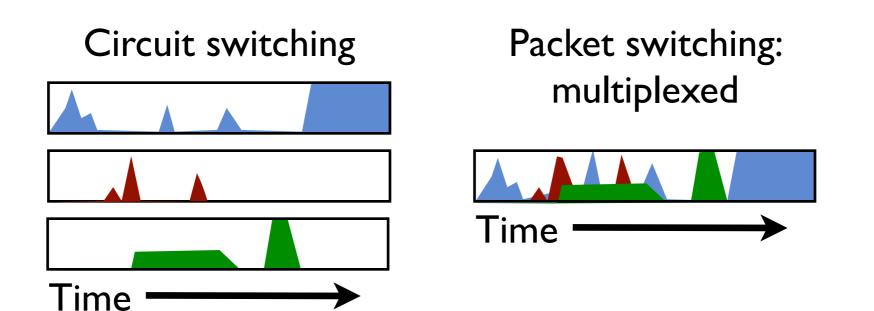
• Servers, switches, NICs...

Statistical multiplexing

- Must provision for peak load
- Many users sharing a resource are unlikely to have their peaks all at the same time

#### Statistical multiplexing

- Must provision for peak load
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- Just as in packet switching



#### Challenges

- Confidentiality of data and computation
- Isolation of resources
- Integration with existing systems
- Robustness
- Latency
- Bandwidth
- Programmability
- ...

#### Opportunities

- New systems and architectures
- Optimizations matter

#### Servers are expensive!

Amortized Cost	Component	Sub-Components
$\sim 45\%$	Servers	CPU, memory, storage systems
$\sim 25\%$	Infrastructure	Power distribution and cooling
~15%	Power draw	Electrical utility costs
~15%	Network	Links, transit, equipment

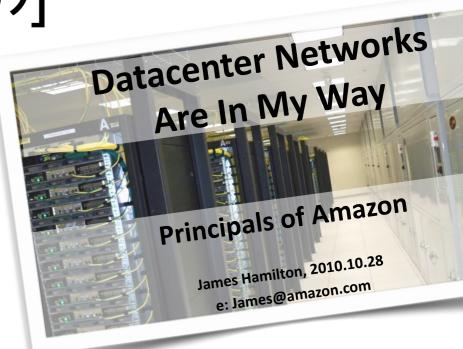
[Greenberg, CCR Jan. 2009]

Agility: Use any server for any service at any time

- Increase utilization of servers
- Reduce costs, increase reliability

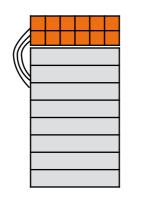
#### What we need [Greenberg, ICDCS'09]

- Rapid installation of service's code
  - Solution: virtual machines
- Access to data from anywhere
  - Solution: distributed filesystems



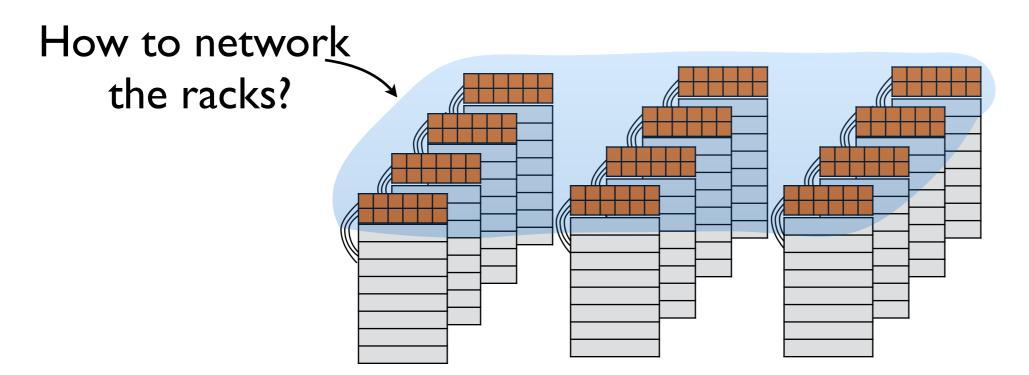
• Ability to communicate between servers quickly, regardless of where they are in the data center





A top-of-rack switch

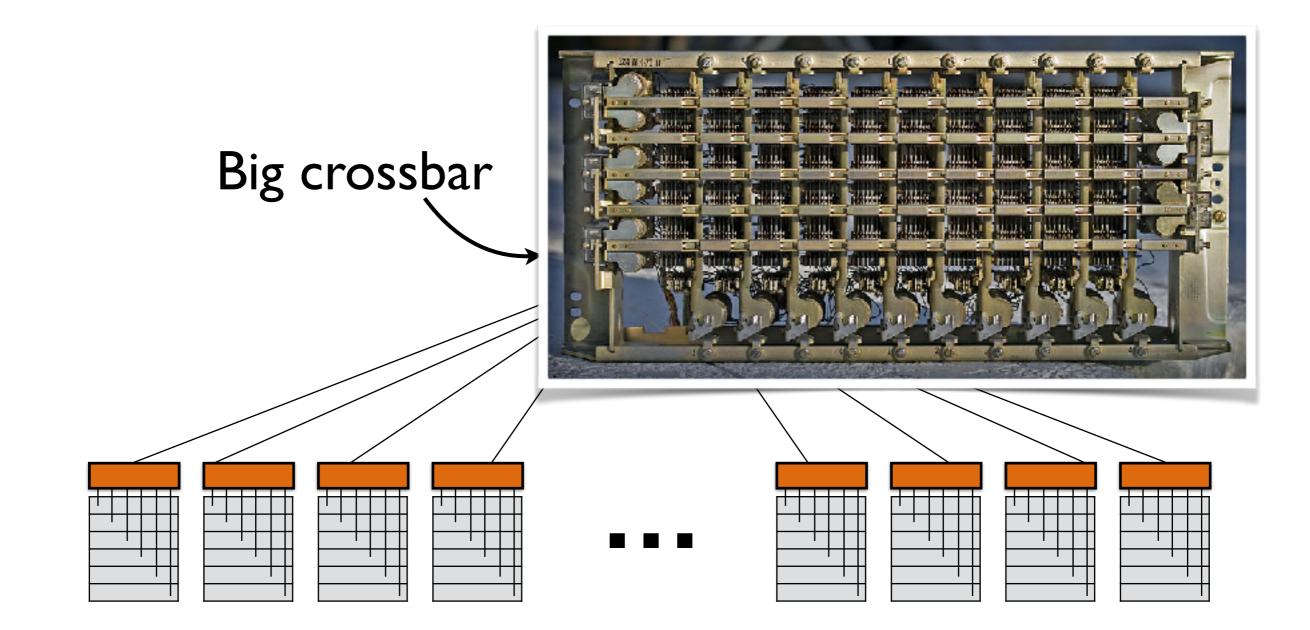
A rack of servers



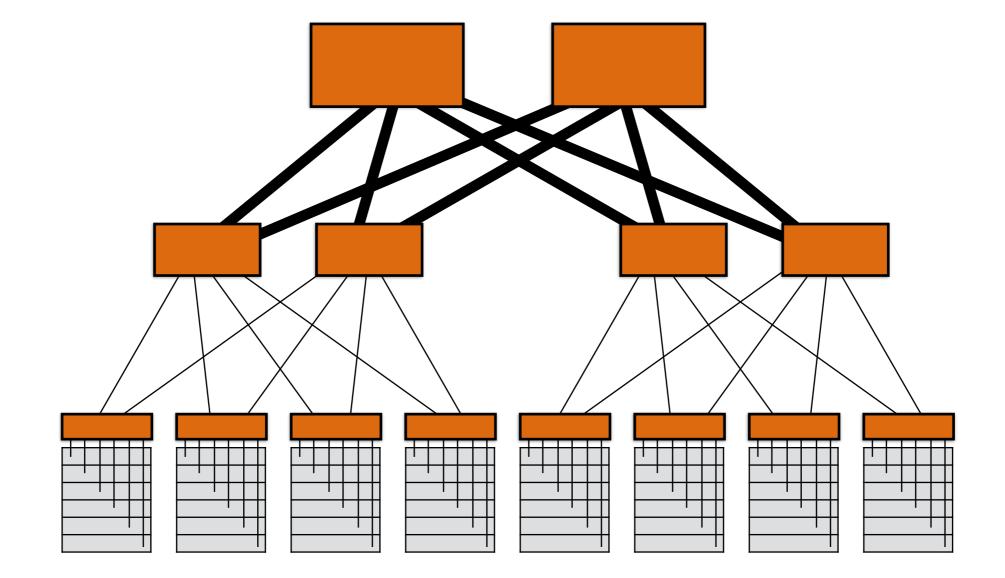
Facebook: machine-machine traffic "doubling at an interval of less than a year"

#### "Big switch" approach

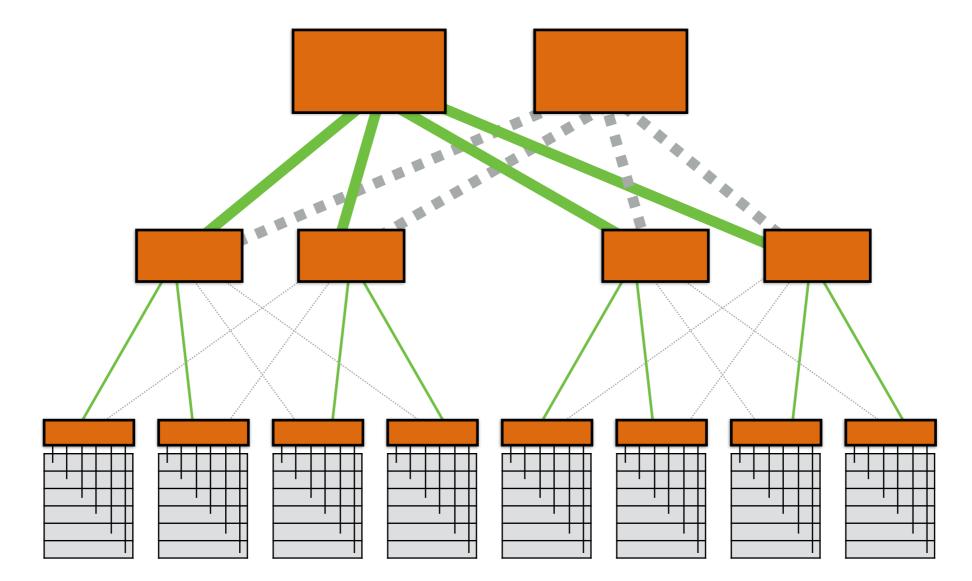




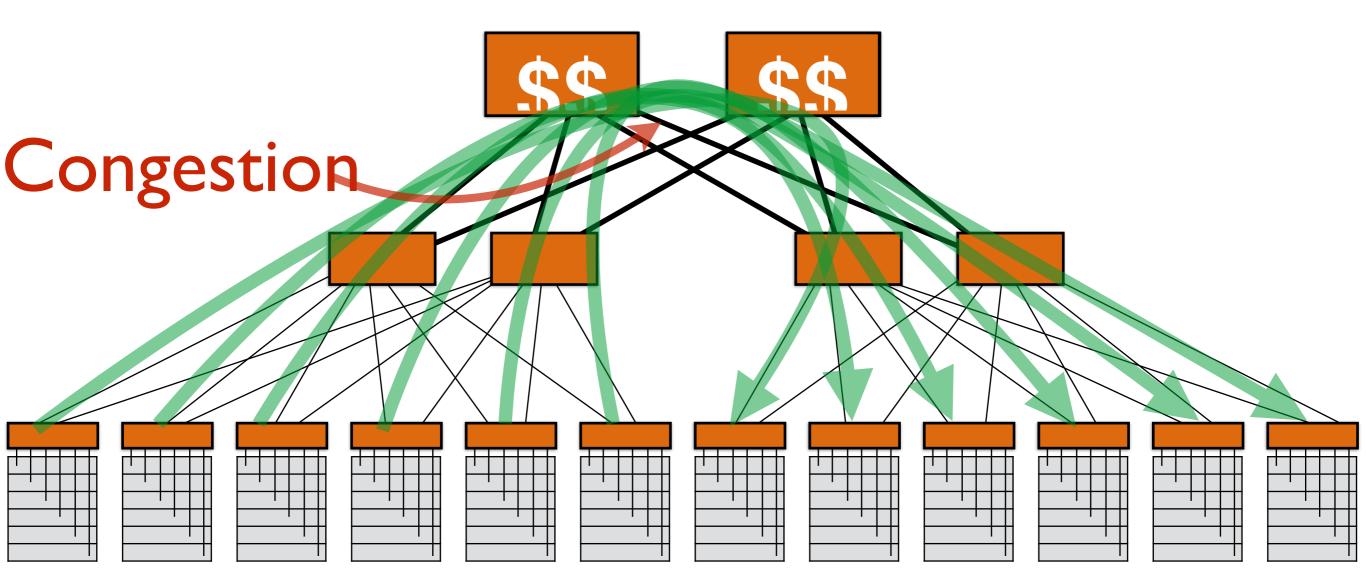
### Alternative: tree network



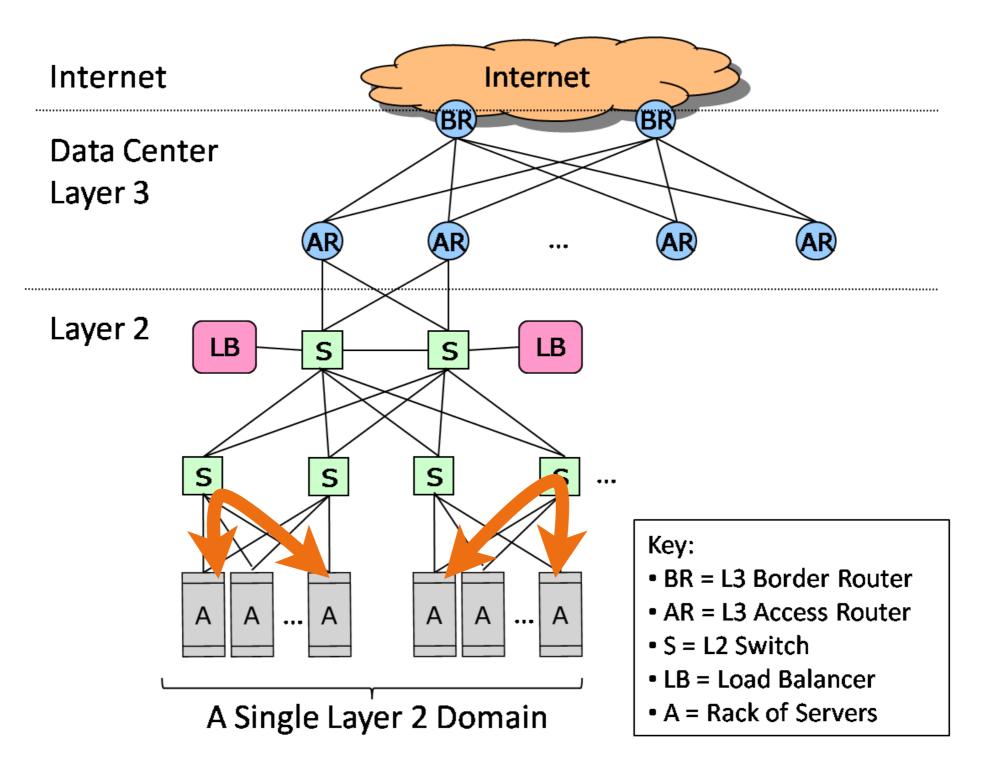
# Alternative: tree network



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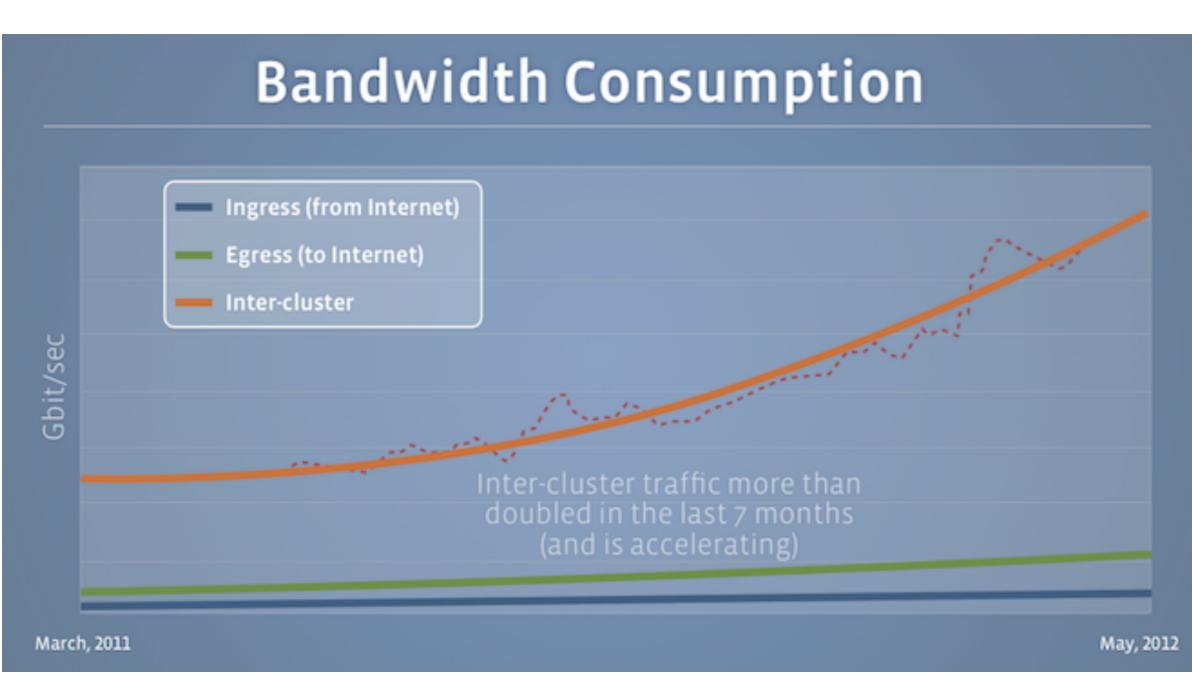


### Traditional data center network



[Greenberg et al, CCR Jan. 2009]

### The need for performance



March 2011

[Facebook, via Wired]

May

2012

#### Modern Data Center Networks

### Fat trees in data centers

#### Argued for nonblocking bandwidth

- Servers limited only by their network card's speed, regardless of communication pattern between servers
- Also known as full throughput in the "hose model"
  - Maximum rate input from each "hose" (host)
  - Maximum rate output to each "hose"
  - Subject to those constraints, any traffic pattern is OK

A Scalable, Commodity Data

Center Network Architecture

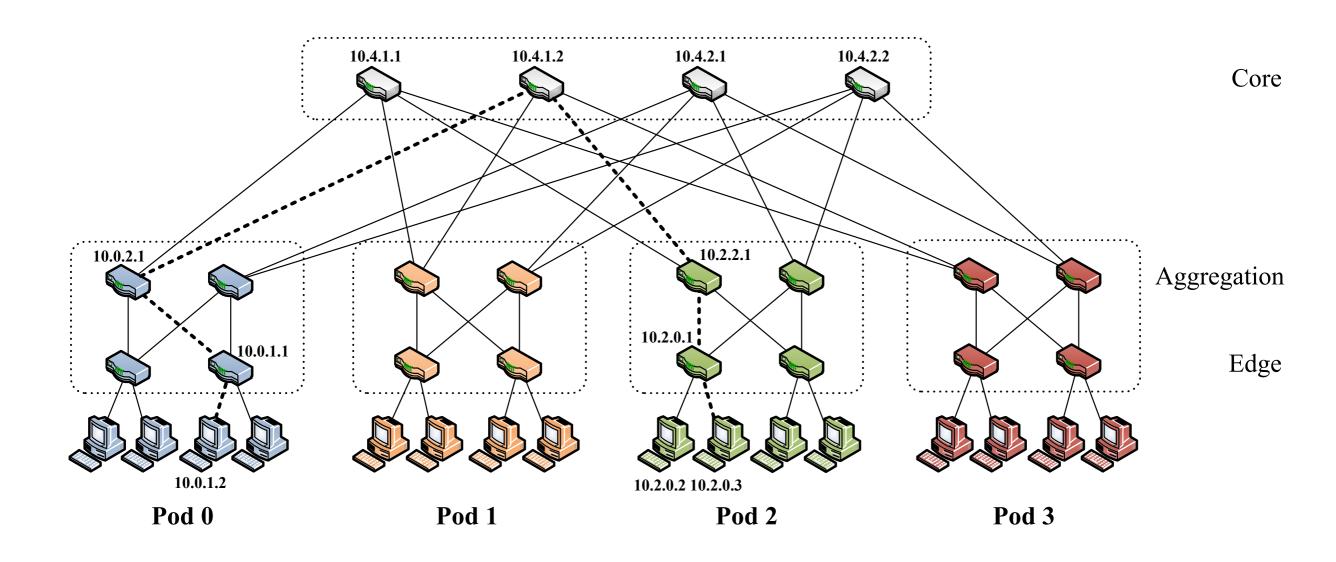
Mohammad Al-Fares, Alexander

Loukissas, Amin Vahdat

SIGCOMM 2008

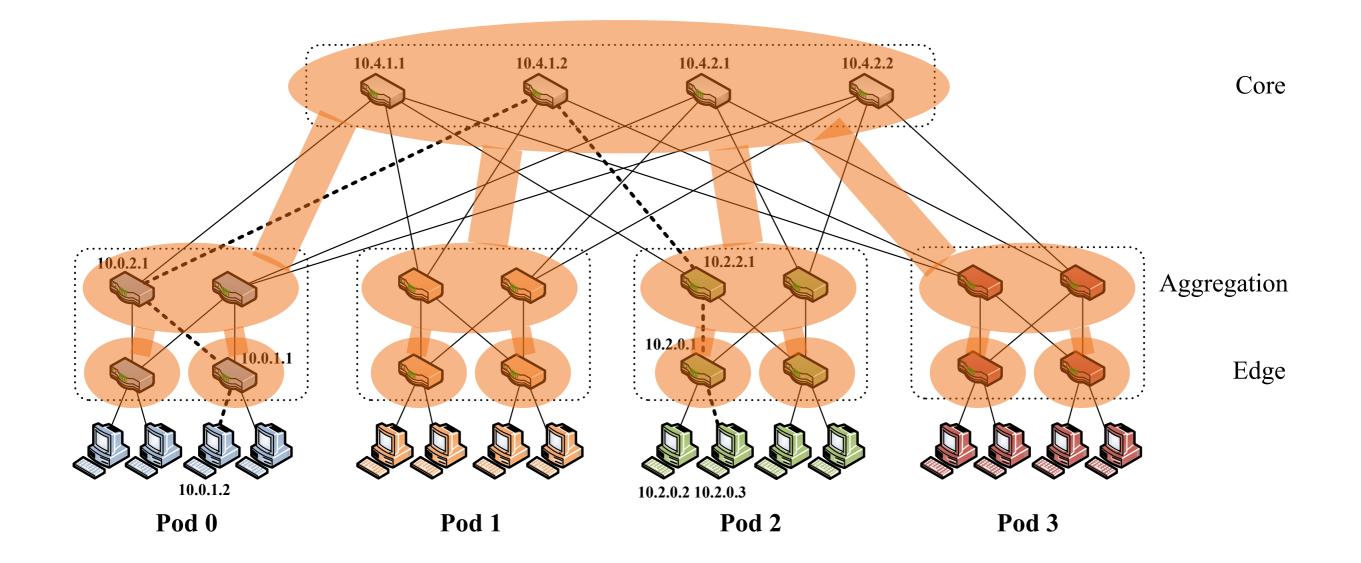
#### Design

- Employed large number of commodity switches rather than "big iron"
- Arranged in Clos topology, and specifically a "fat tree"



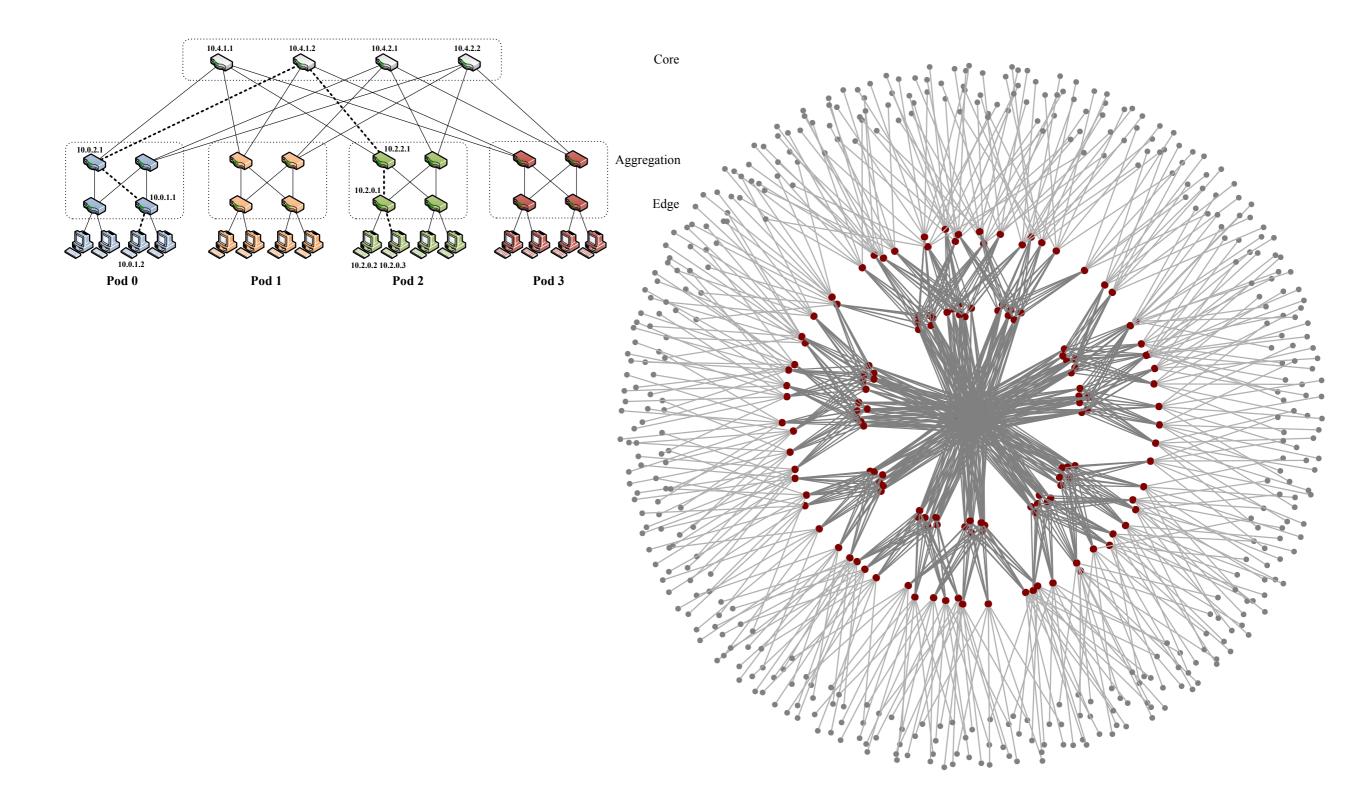
[Al-Fares, Loukissas,Vahdat, SIGCOMM '08]



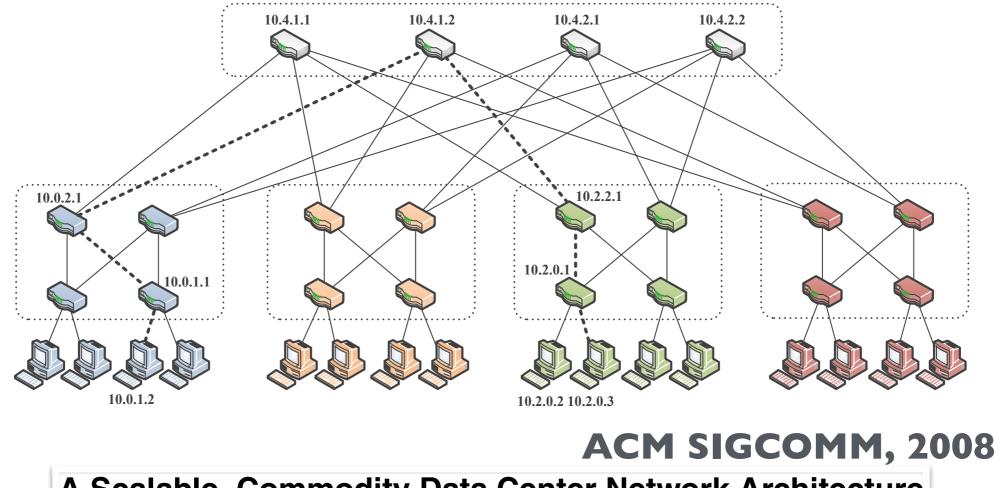


[Al-Fares, Loukissas,Vahdat, SIGCOMM '08]



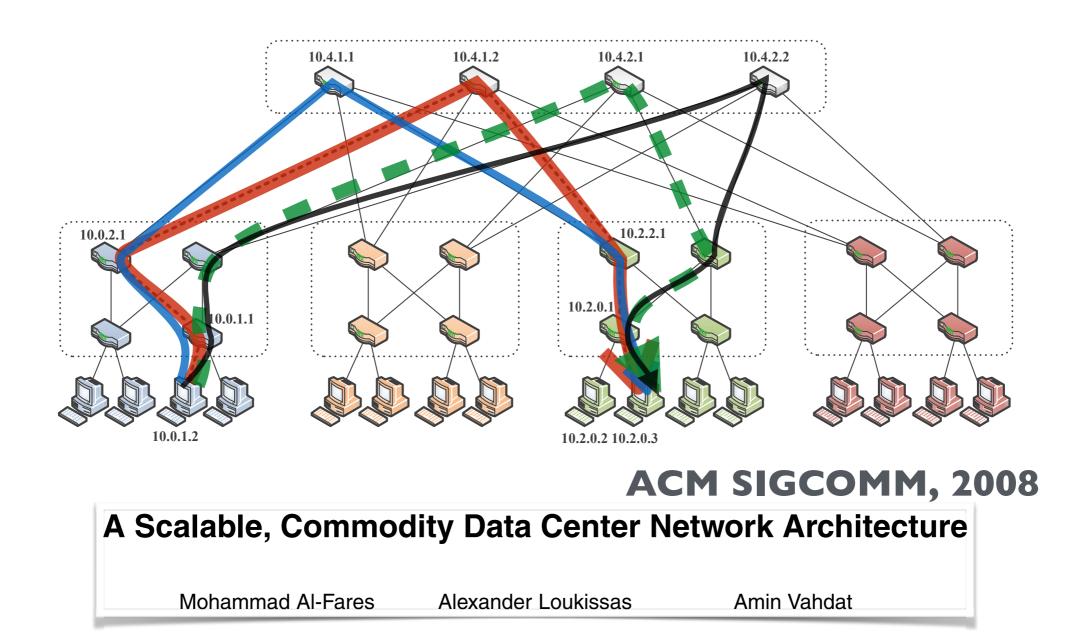




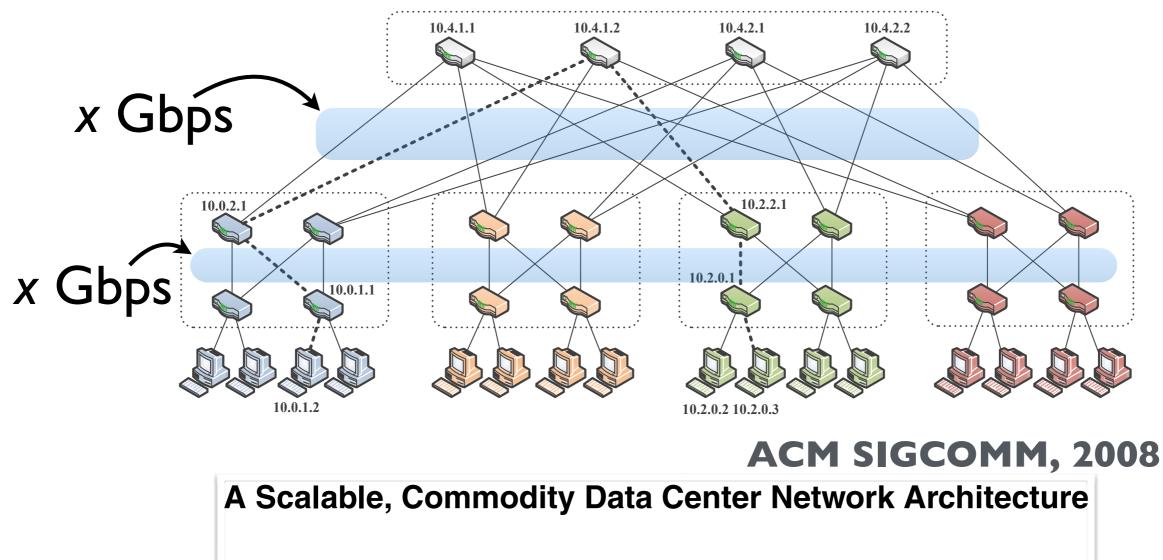








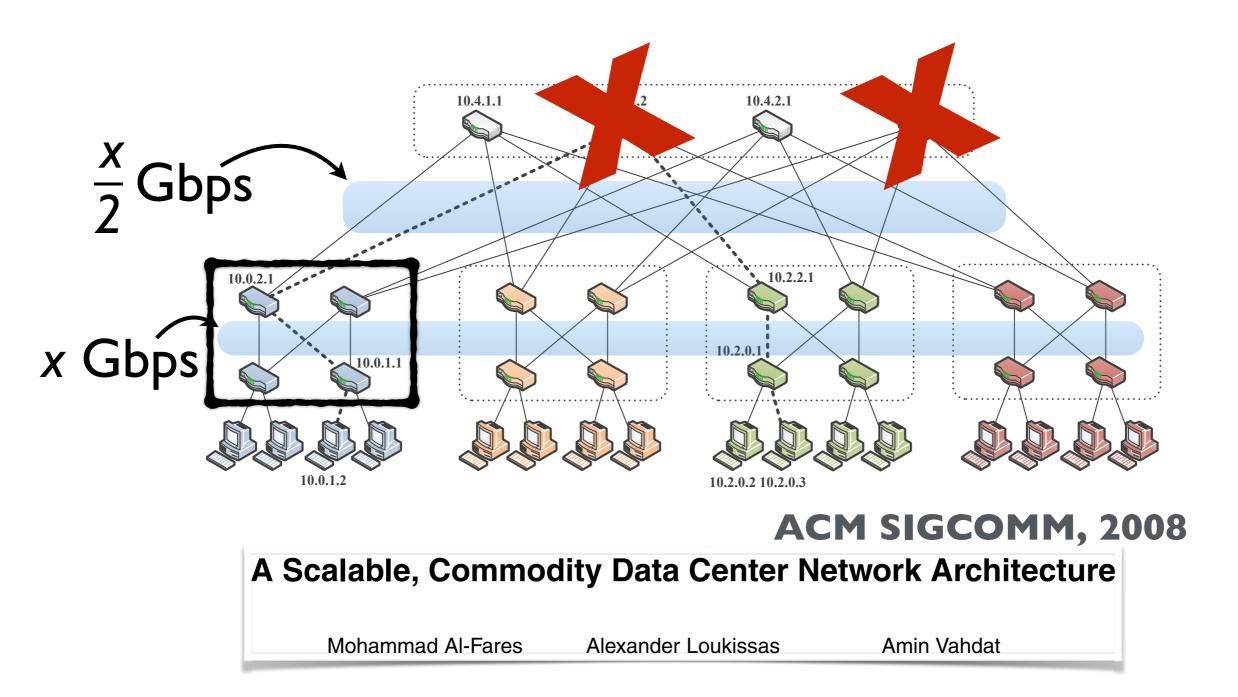




Mohammad Al-Fares Alexander Loukissas

Amin Vahdat

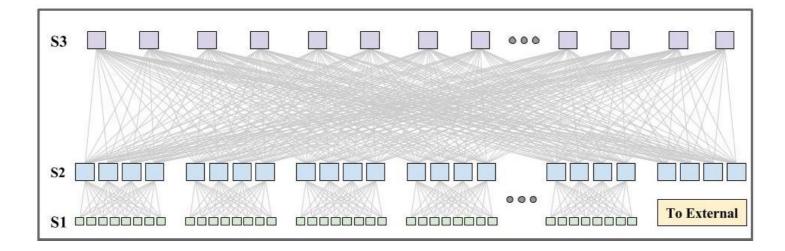
### Oversubscribed fat tree



#### ACM SIGCOMM, 2015

#### Jupiter Rising: A Decade of Clos Topologies and Centralized Control in Google's Datacenter Network

Arjun Singh, Joon Ong, Amit Agarwal, Glen Anderson, Ashby Armistead, Roy Bannon, Seb Boving, Gaurav Desai, Bob Felderman, Paulie Germano, Anand Kanagala, Jeff Provost, Jason Simmons, Eiichi Tanda, Jim Wanderer, Urs Hölzle, Stephen Stuart, and Amin Vahdat Google, Inc.



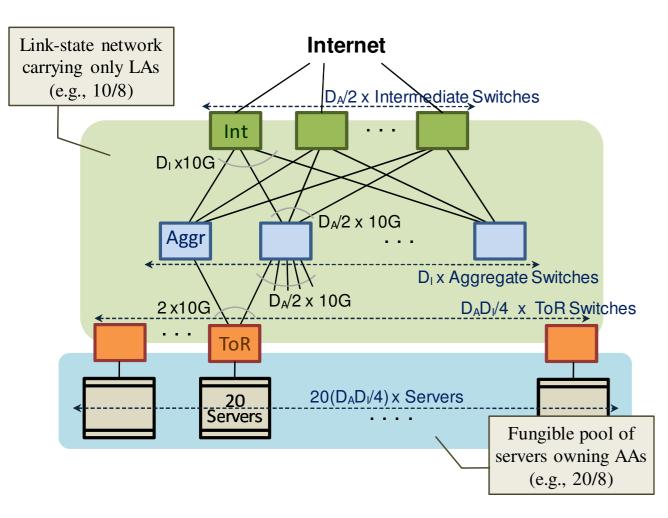
#### age: Robert Harker]

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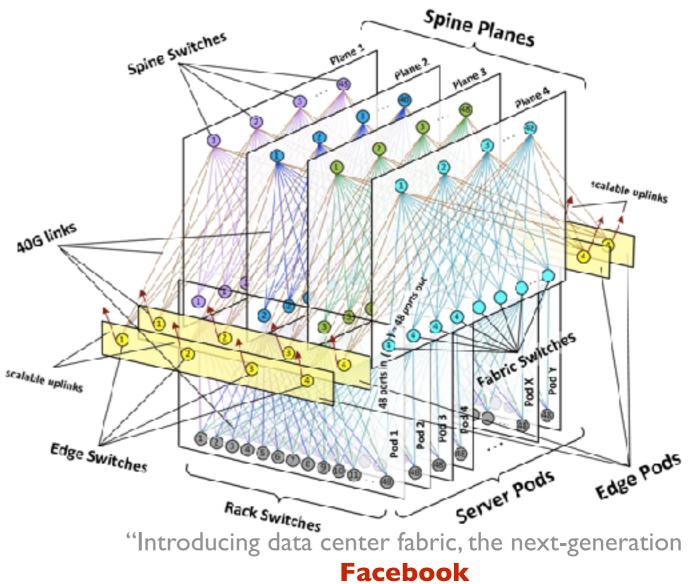
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## Variants of this design are common



VL2 @ Microsoft, ACM SIGCOMM'09 Greenburg, Hamilton, Jain, Kandula, Kim, Lahiri, Maltz, Patel, Sengupta



data center network", Alexey Andreyev, 2015

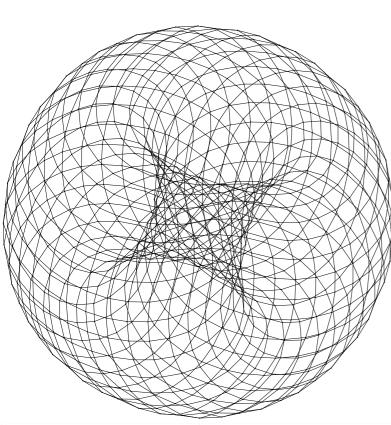


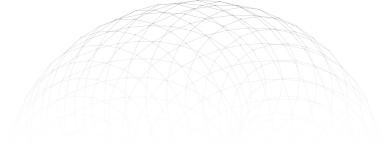
[Greenberg, Hamilton, Jain, Kandula, Kim, Lahiri, Maltz, Patel, Sengupta, SIGCOMM 2009]

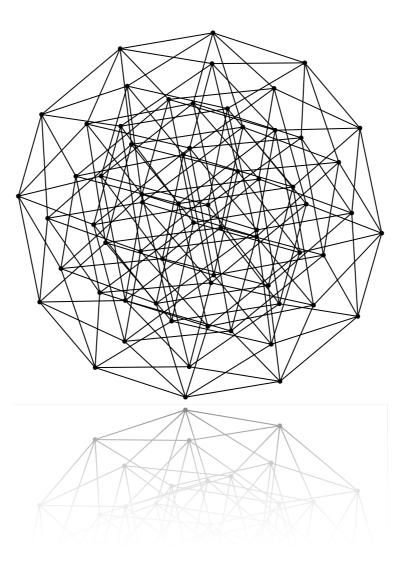
Key features

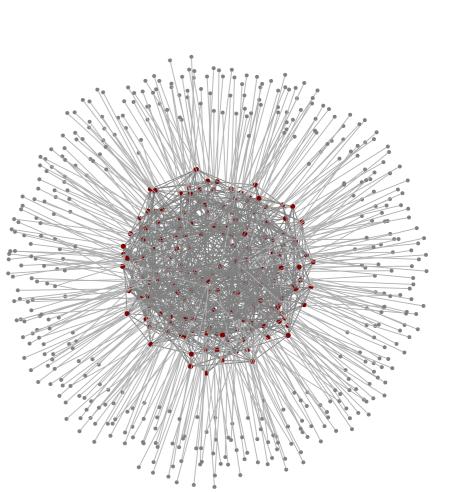
- High bandwidth network
  - Another folded Clos network
  - Slightly different than fat tree (e.g., uses faster 10 Gbps links at higher layers)
- Randomized (Valiant) load balancing
  - Makes better use of network resources
- Flat addressing
  - Ethernet-style (layer 2) addresses to forward data, rather than IP addresses
  - Separates names from locations

## Many other proposed topologies

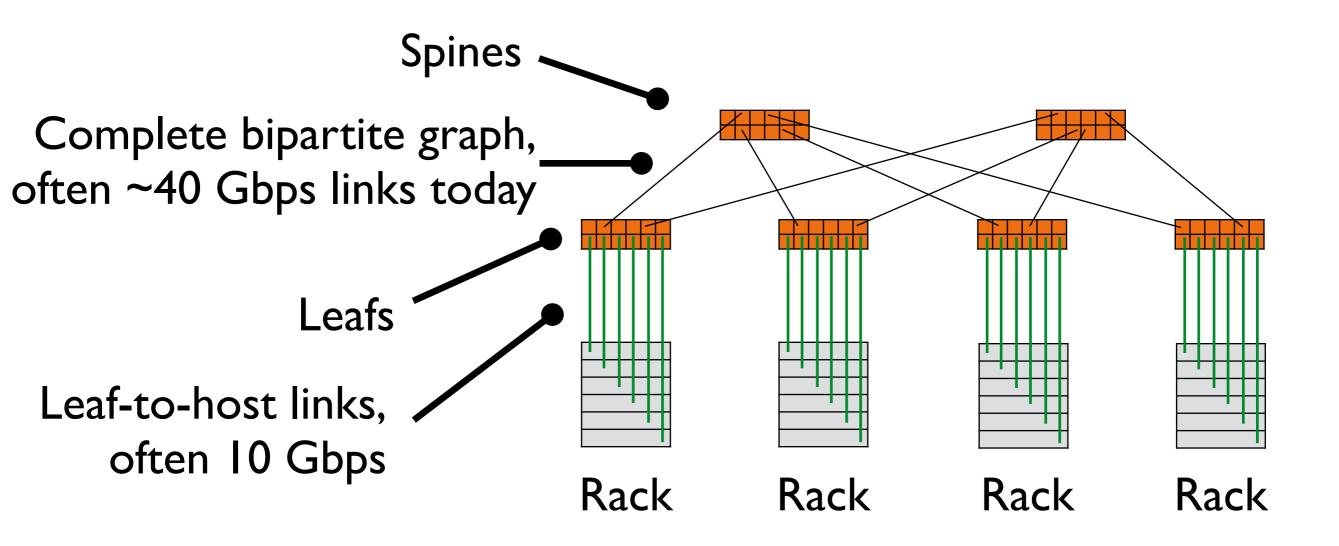








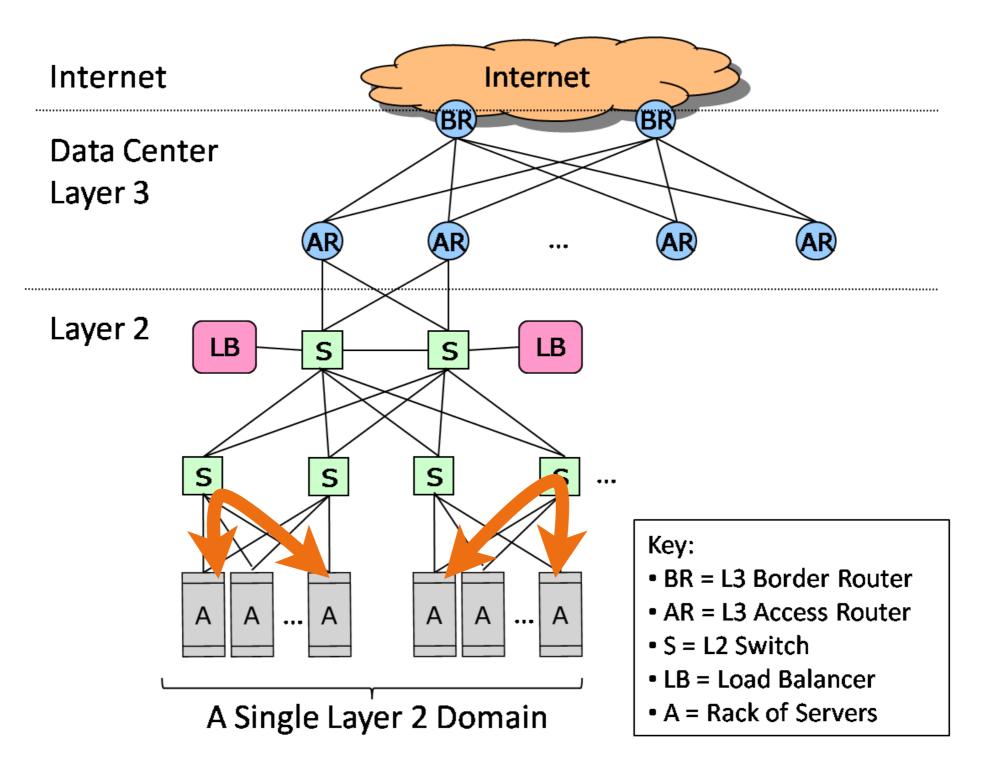
### Leaf-spine for smaller networks



Outside of the hyper scale cloud providers, this 2-tier design is typically scalable enough and is now common.

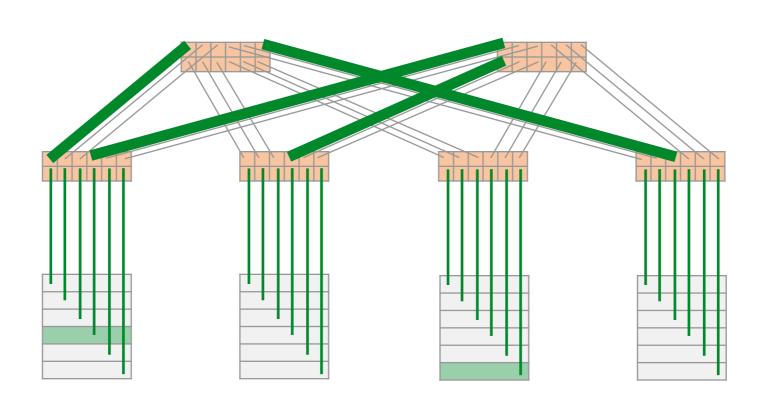
#### **Routing in Data Center Networks**

### Traditional data center network



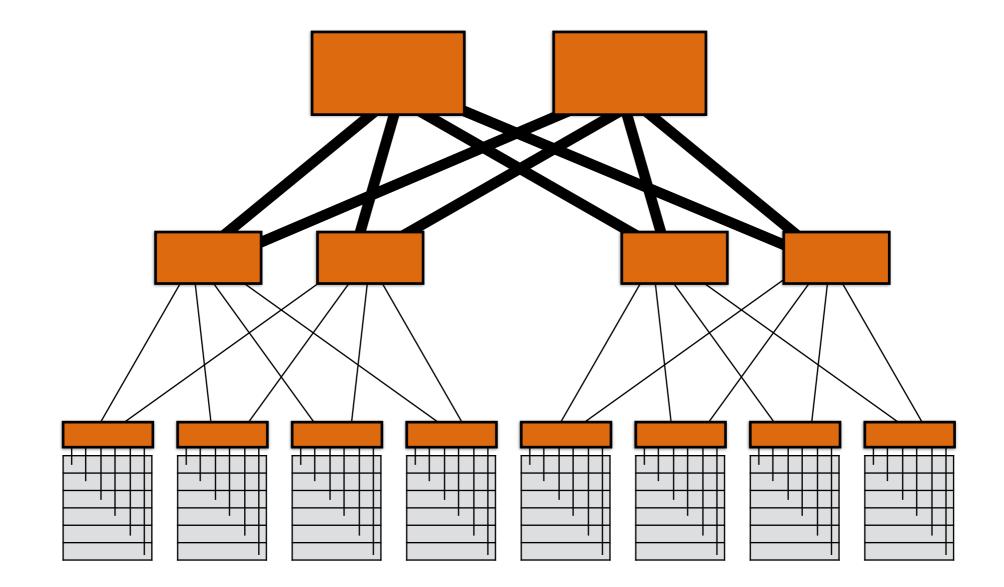
[Greenberg et al, CCR Jan. 2009]

# Spanning tree

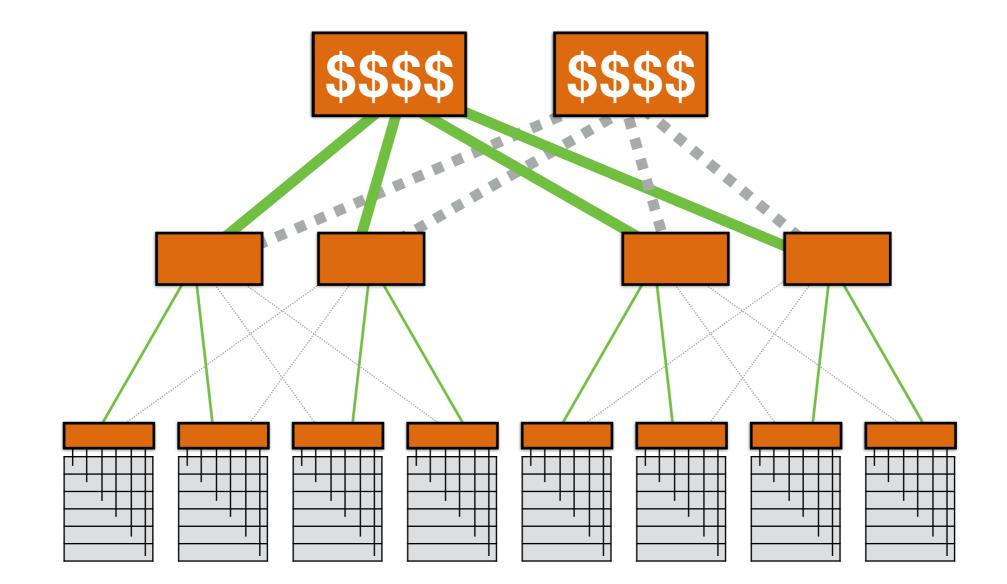


#### Disable all other links!

#### STP works for tree-like networks

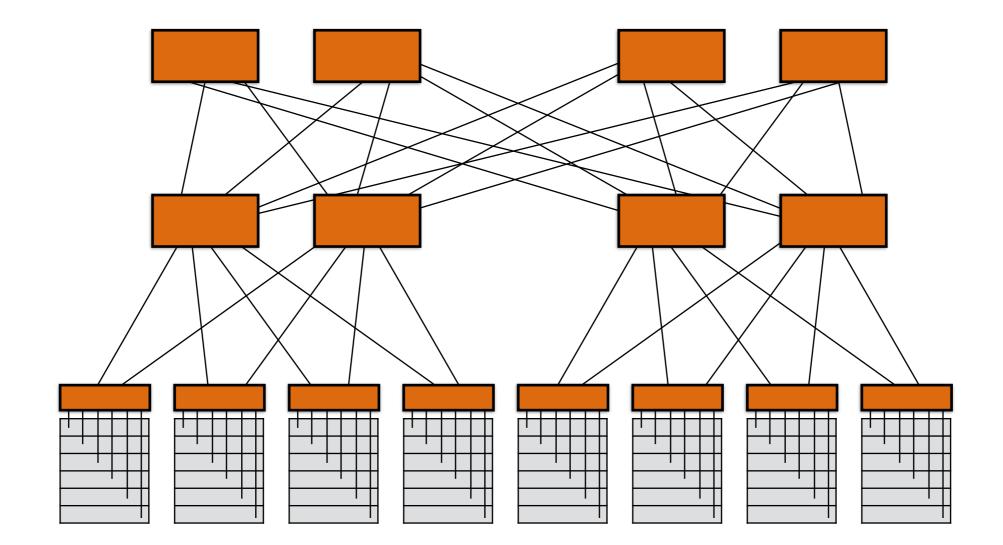


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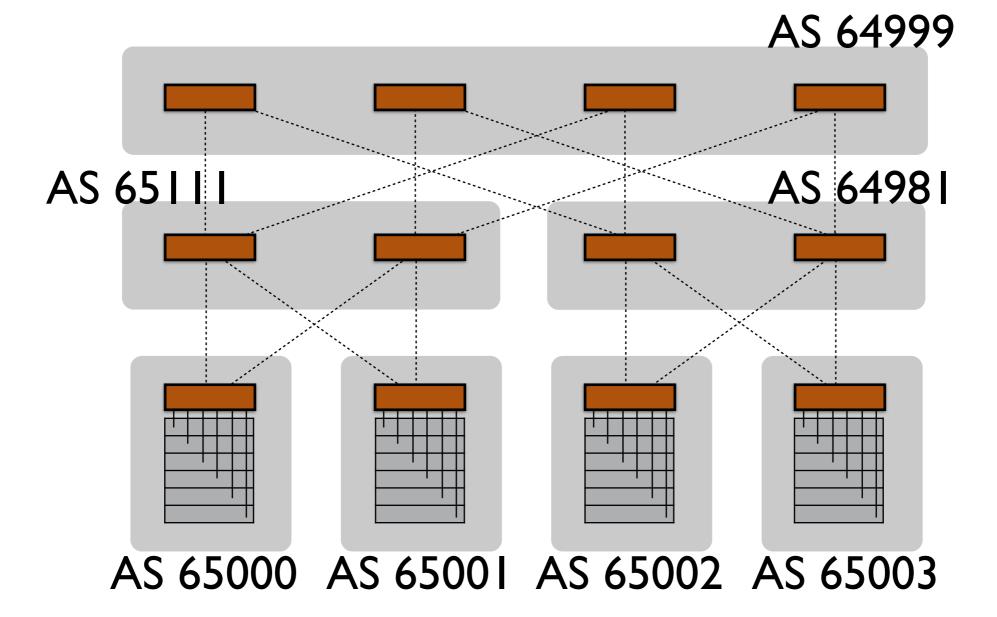


#### STP works for tree-like networks

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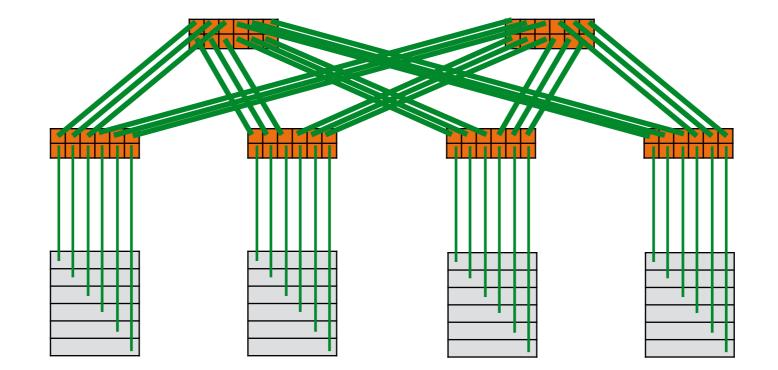


#### BGP in the data center

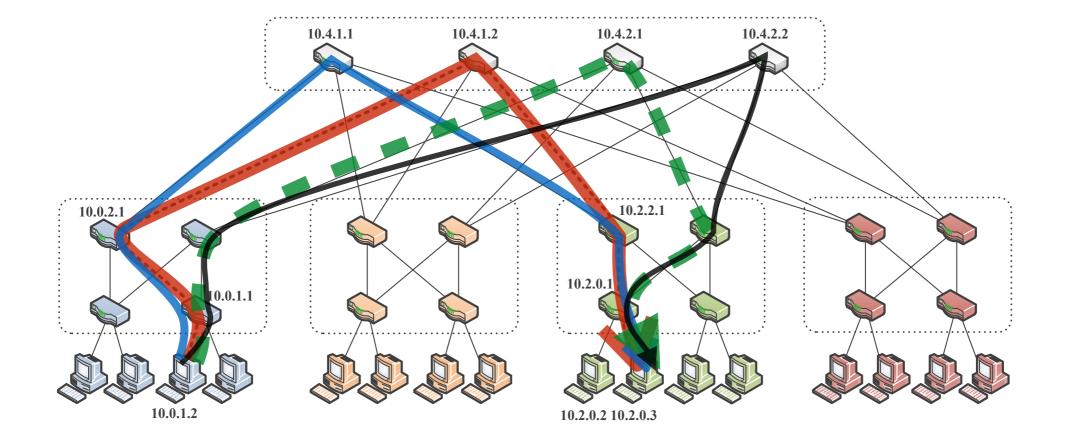


## Routing

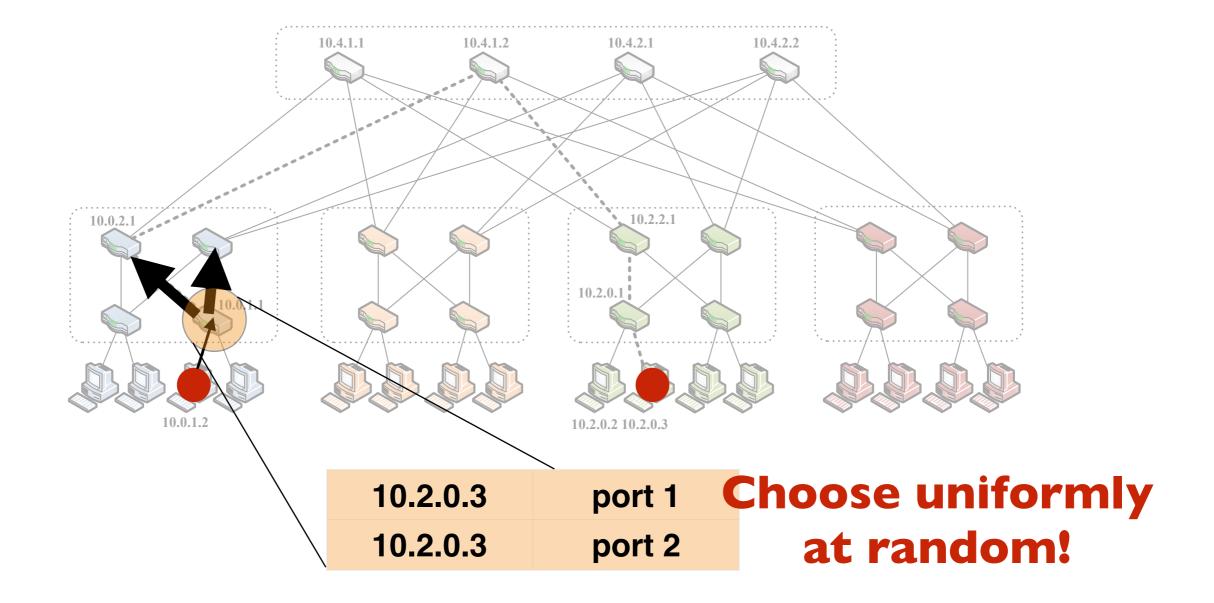


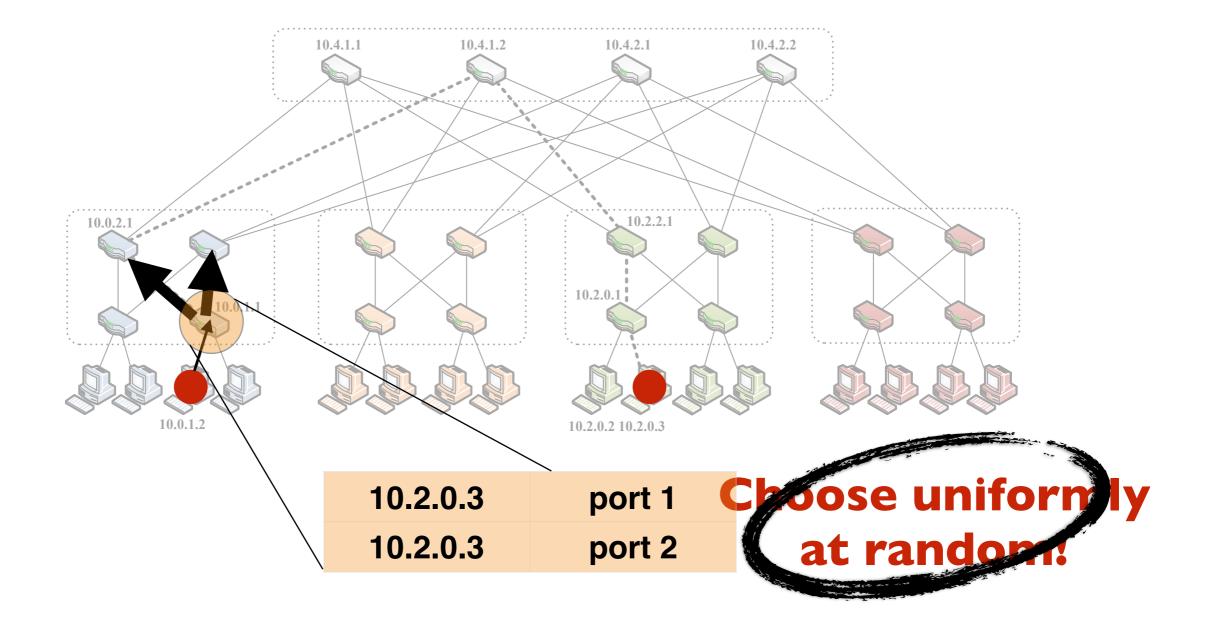


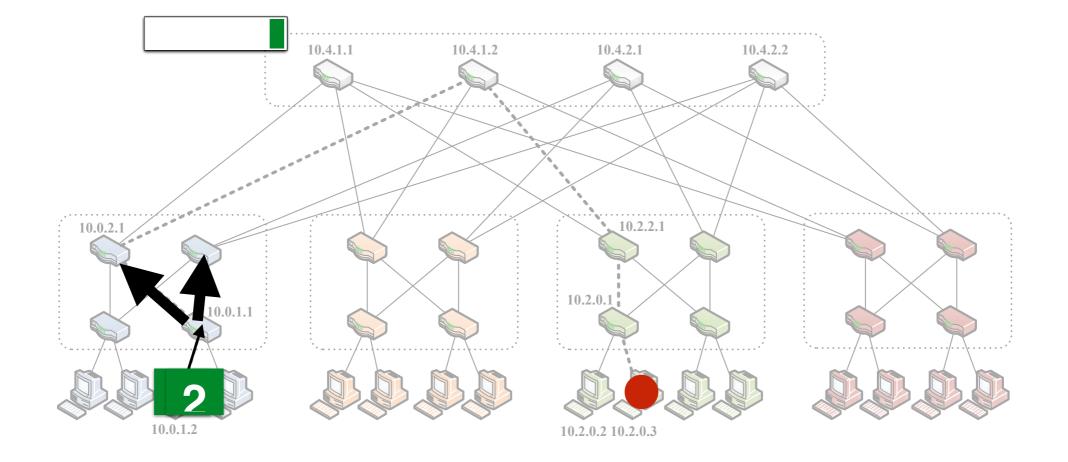
### Multipath routing



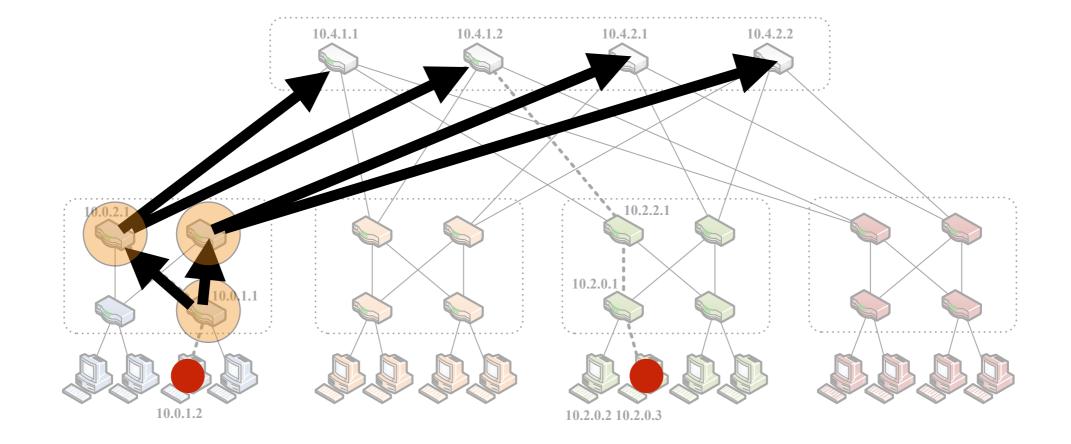
A Scalable, Commodity Data Center Network Architecture					
Mohammad Al-Fares	Alexander Loukissas	Amin Vahdat			
malfares@cs.ucsd.edu	aloukiss@cs.ucsd.edu	vahdat@cs.ucsd.edu			

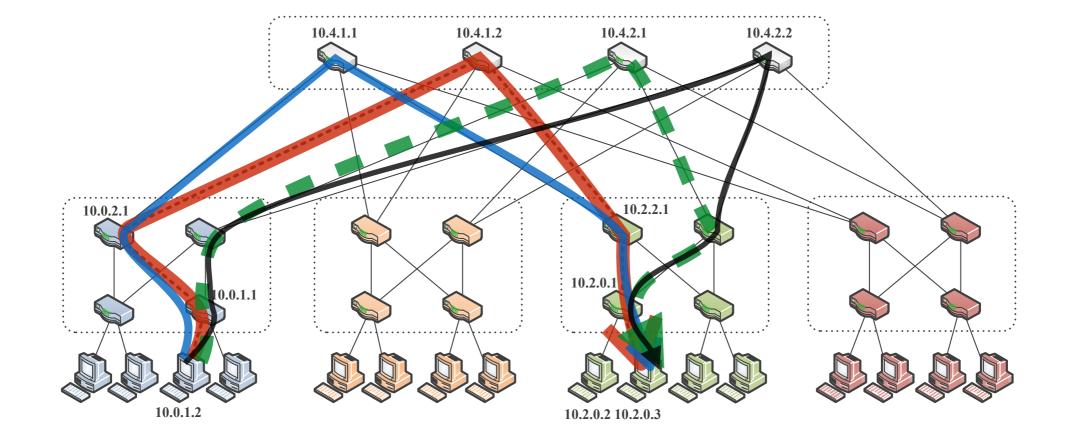




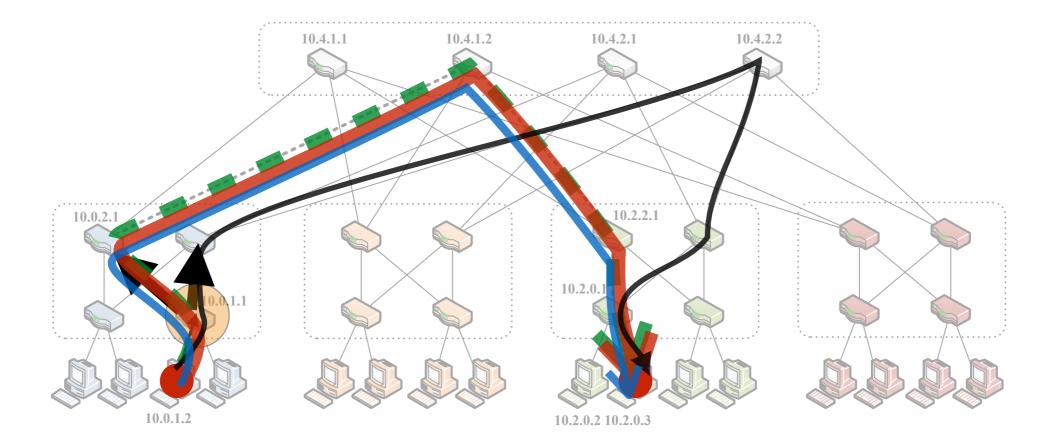


#### Output port = hash (packet header)





#### ECMP: traffic imbalance



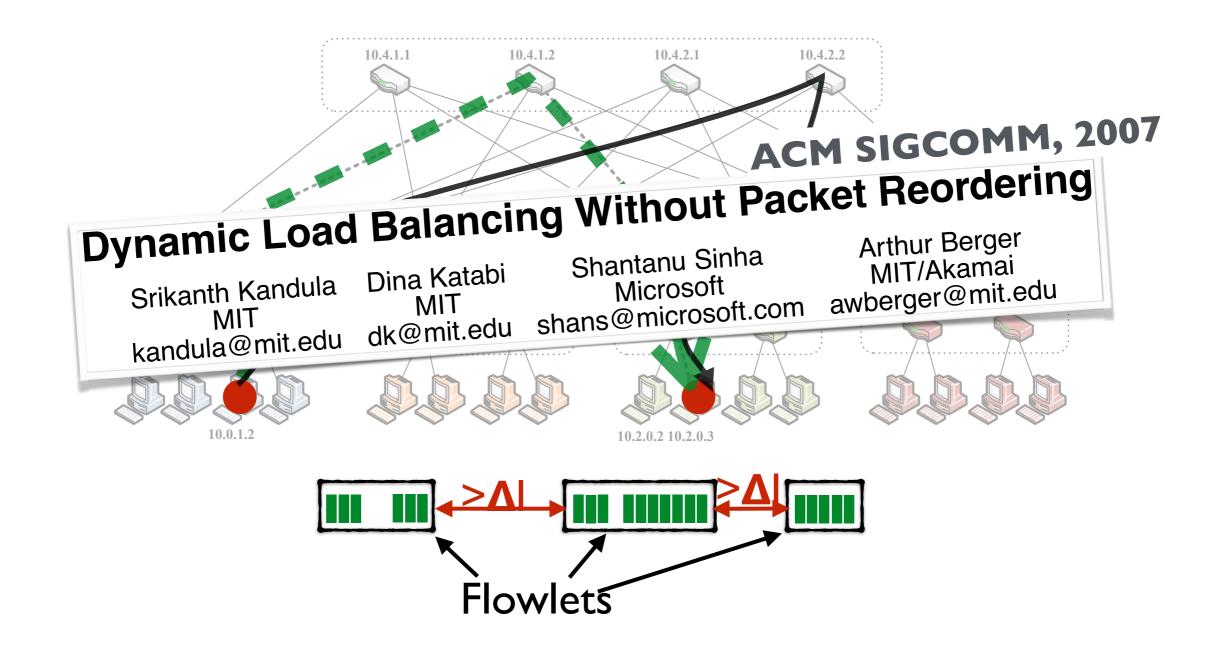
#### Output port = hash (packet header "5-tuple")

IP src & dst. Protocol number Protocol port src & dst.

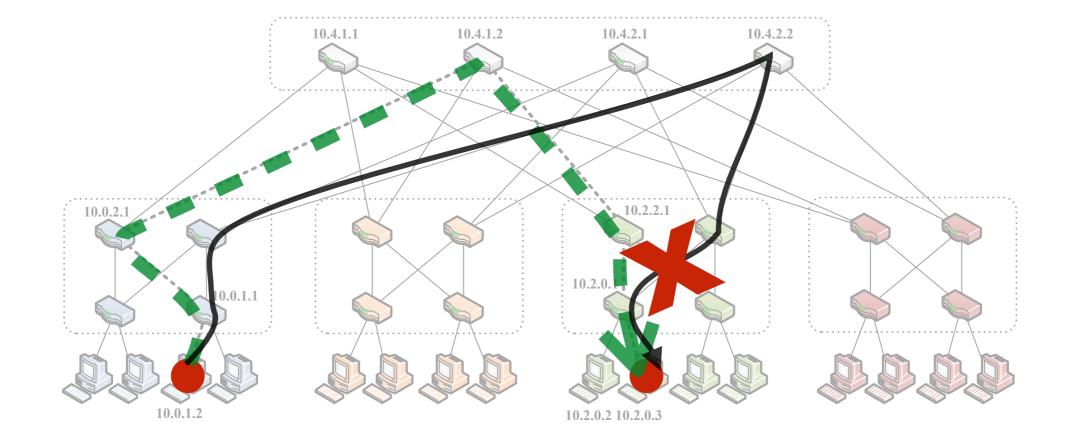
Result: each TCP connection (a "flow") stays on one path

#### Flowlets





#### ECMP: local, oblivious choices

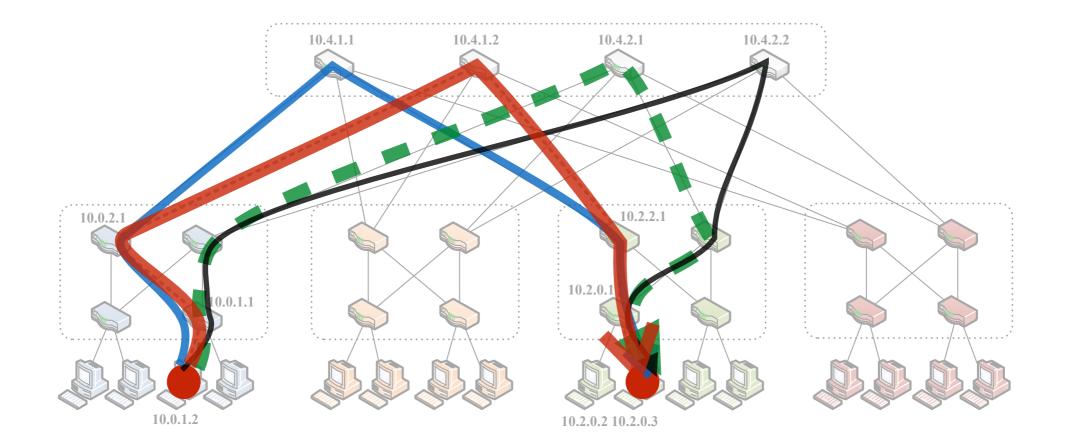


#### Cannot avoid distant failures!

### CONGA: edge based monitoring

#### ACM SIGCOM, 2014 CONGA: Distributed Congestion-Aware Load Balancing for Datacenters Mohammad Alizadeh, Tom Edsall, Sarang Dharmapurikar, Ramanan Vaidyanathan, Kevin Chu, Andy Fingerhut, Vinh The Lam (Google), Francis Matus, Rong Pan, Navindra Yaday, George Varghese (Microsoft) Lisco Systems

### CONGA: edge based monitoring



#### Leaf switch monitors path performance in real time:

	Path 1	Path 2	Path 3	Path 4
Dest. switch	3	7	2	1

Physical layer

Topology

Load balancing

Transport

Control & SDN